# 111 JANUARY 1986 \$2.95 (3.95 CANADA)

## Dr. Dobb's Journal of

# Software Tools

FOR THE PROFESSIONAL PROGRAMMER

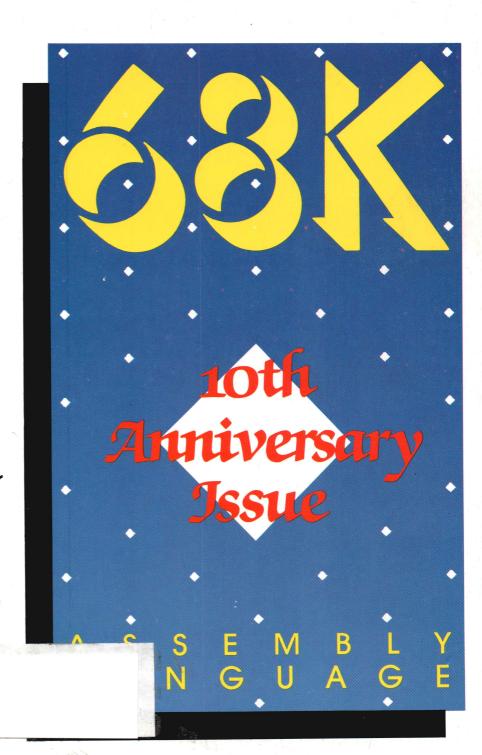
## Launching Our Second Decade

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An 8080 Simulator for the 68K

DOS Shell Notes on 80286 Big DOS







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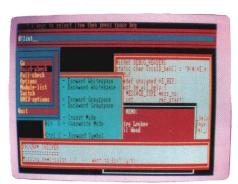
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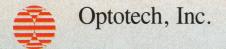
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#### Dr. Dobb's Journal of

## ware Too

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Our anniversary cover was designed by Shelley Rae Doeden, who also is responsible for the new look of the magazine. Shelley is DDJ's Art Director.

Mysterious and dangerous programs >

C CHEST: A Unix-like Shell for MS DOS 18 by Allen Holub

Allen describes how the shell works on a high level and includes examples of some often-used functions.

118 16-BIT TOOLBOX: Trojan Horse Programs by Ray Duncan

These mysterious and destructive programs are appearing on bulletin boards. What can be done to combat them?

#### This Issue:

The Motorola 68000 chip is, some say, a programmer's processor. The instruction set is rich and logically constructed, and memory access is simple and capacious. What it lacks, others say, are a standard operating system and development tools. This month we focus on programming tools for the 68K.

**EDITORIAL: The New** Look by Michael Swaine

FORUM

DDJ is now on ▶ CompuServe.

programmers

don't know can

What C

hurt them. >

**LETTERS: Comment** by our readers **CARTOON: A Word on** Formats by Rand Renfroe VIEWPOINT: Inefficient C by Hal Hardenberg DDJ ON LINE: The Electronic Edition by Frank DeRose

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PROGRAMMERS'

**SERVICES** 

New products of interest to programmers ADVERTISER INDEX:

Where to find that ad

#### **Next Issue:**

Next month we'll look at structured programming and at some languages that have a reputation for encouraging structured design. We'll publish programming tools in Pascal, Modula-2, and Ada. We'll tell you where to get public-domain Ada utilities to speed Ada software development, and we'll look at a program that ports between dialects of Pascal.

3

# 了一个一个一个一个

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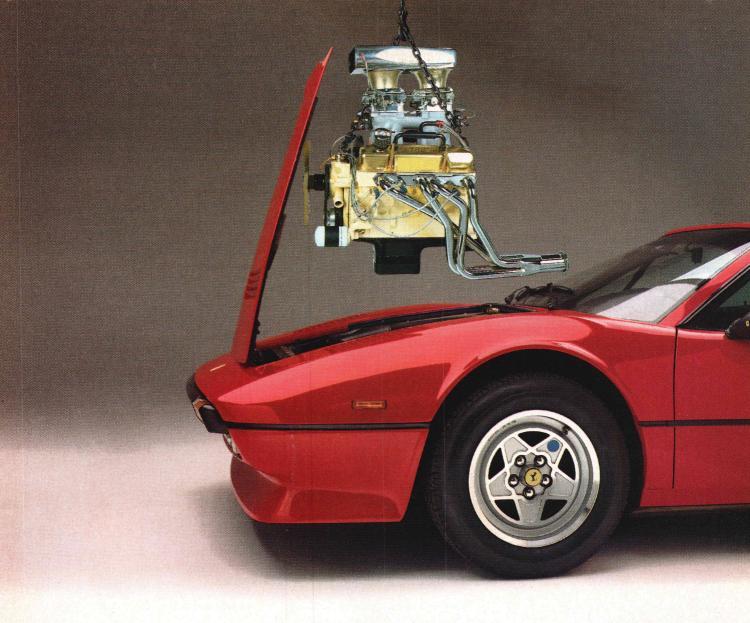
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Programmers' Pfantasies<sup>TM</sup>

XT and AT are trademarks of International Business Machines Corporation. Pfaster286 and Programmers' Pfantasies are trademarks of Phoenix Computer Products Corporation For the Ferrari aficionado: yes, we know this is a rear engine car. We are showing the addition of a second engine to symbolize how Pfaster can be added to your PC or XT to increase performance

ate 1975 and early 1976 saw a number of new magazines for a new group of readers: the pioneering users of the first microcomputers. Most were heavy on hardware because there was essentially no software to write

about back then. In January 1976, a remedy for this situation appeared in the form of a magazine providing the tools needed to develop software: Dr. Dobb's Journal of Tiny BASIC Calisthenics & Orthodontia.

Ten years later, DDJ is still providing software tools for serious programmers. We've replaced "calisthenics and orthodontia" with other metaphors and have moved upscale from a black-and-white newsletter to a four-color magazine, but we still publish the only pages in which you can find compilers, assemblers, and programming tools reviewed, designed, analyzed, and source-listed, with the whole process watched over by the most knowledgable readership in the industry.

Was that last paragraph excessively self-congratulatory? Pardon it, but there is occasion for it: This month is DDJ's tenth birthday. This birthday issue features articles on 68000 programming, leading off with Edward Ream's assembler that thinks it's a compiler. In a more overtly festive act, this issue also unveils the new design for DDJ.

The new design should more directly indicate what the magazine provides: programming tools in the form of articles and listings, reviews, columns, a forum for discussion of programming issues, and services such as product listings. The title treatment, while harking back to the original cover design, makes it clear that what this magazine contains are software tools. Inside, we've gathered all the contents into five sec-



tions: Articles, Reviews, Columns, Forum, and Programmers' Services. Forum contains the editorial, letters, an editorial cartoon, and an invited viewpoint-this month by 68K authority Hal Hardenberg. Programmers' Services

will grow; this month it contains a section of information on the professional side of programming. No room for reviews this month.

We are fortunate to have some excellent columnists who provide a personal view along with useful tips and utilities. Unfortunately, two of them are retiring from their columnist duties: Dave Cortesi and Bob Blum. Bob is just too busy to keep up a monthly column and to do as good a job as he wants; he'll continue to contribute articles and to maintain his CP/M bulletin board. He has written a sign-off letter that appears on page 8. Dave is leaving computer writing for the writing of fiction. Both were longtime contributors and will be missed. We are actively, if grudgingly, seeking replacements. Perhaps you know a candidate?

Finally, there's the electronic component we've added to the magazine. By the time you read this, DDJ page 17 for details.

Michael Swans should be up on CompuServe. See

Michael Swaine



#### Dr. Dobb's Journal of

#### **Software Tools**

#### **Editorial**

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## The C for Microcomputers

PC-DOS, MS-DOS, CP/M-86, Macintosh, Amiga, Apple II, CP/M-80, Radio Shack, Commodore, XENIX, ROM, and Cross Development systems

#### MS-DOS, PC-DOS, CP/M-86, XENIX, 8086/80x86 ROM

#### Manx Aztec C86

"A compiler that has many strengths ... quite valuable for serious work"

Computer Language review, February 1985

Great Code: Manx Aztec C86 generates fast executing compact code. The benchmark results below are from a study conducted by Manx. The Dhrystone benchmark (CACM 10/84 27:10 p1018) measures performance for a systems software instruction mix. The results are without register variables. With register variables, Manx, Microsoft, and Mark Williams run proportionately faster, Lattice and Computer Innovations show no improvement.

	Execution Time	Code Size	Compile/ Link Time
Dhrystone Benchmark	(		
Manx Aztec C86 3.3	34 secs	5,760	93 secs
Microsoft C 3.0	34 secs	7,146	119 secs
Optimized C86 2.201	53 secs	11,009	172 secs
Mark Williams 2.0	56 secs	12,980	113 secs
Lattice 2.14	89 secs	20,404	117 secs

Great Features: Manx Aztec C86 is bundled with a powerful array of well documented productivity tools, library routines and features.

and teatures.
Optimized C compiler
AS86 Macro Assembler
80186/80286 Support
8087/80287 Sensing Lib
Extensive UNIX Library
Large Memory Model
Z (vi) Source Editor -c
ROM Support Package -c
Library Source Code -c
MAKE, DIFF, and GREP -c
One year of updates -c
CP/M-8
COMPM-8
COMPM

Symbolic Debugger
LN86 Overlay Linker
Librarian
Profiler
DOS, Screen, & Graphics Lib
Intel Object Option
CP/M-86 Library -c
INTEL HEX Utility -c
Mixed memory models -c
Source Debugger -c
CP/M-86 Library -c

Manx offers two commercial development systems, Aztec C86-c and Aztec C86-d. Items marked -c are special features of the Aztec C86-c system.

Aztec C86-c Commercial System	\$499
Aztec C86-d Developer's System	\$299
Aztec C86-p Personal System	\$199
Aztec C86-a Apprentice System	\$49

All systems are upgradable by paying the difference in price plus \$10.

Third Party Software: There are a number of high quality support packages for Manx Aztec C86 for screen management, graphics, database management, and software development.

C-tree \$395	Greenleaf \$185
PHACT \$250	PC-lint \$98
HALO \$250	<b>Amber Windows \$59</b>
PRE-C \$395	Windows for C \$195
WindScreen \$149	FirsTime \$295
SunScreen \$99	C Util Lib \$185
PANEL \$295	Plink-86 \$395

#### MACINTOSH, AMIGA, XENIX, CP/M-68K, 68k ROM

#### Manx Aztec C68k

"Library handling is very flexible ... documentation is excellent ... the shell a pleasure to work in ... blows away the competition for pure compile speed ... an excellent effort."

Computer Language review, April 1985

Aztec C68k is the most widely used commercial C compiler for the Macintosh. Its quality, performance, and completeness place Manx Aztec C68k in a position beyond comparison. It is available in several upgradable versions.

Optimized C
Macro Assembler
Overlay Linker
Resource Compiler
Debuggers
Librarian
Source Editor
MacRam Disk -c

Library Source -c

Creates Clickable Applications Mouse Enhanced SHELL Easy Access to Mac Toolbox UNIX Library Functions Terminal Emulator (Source) Clear Detailed Documentation C-Stuff Library UniTools (vi,make,diff,grep) -c One Year of Updates -c

Items marked -c are available only in the Manx Aztec C86-c system. Other features are in both the Aztec C86-d and Aztec C86-c systems.

Aztec C68k-c Commercial System	\$499
Aztec C68d-d Developer's System	\$299
Aztec C68k-p Personal System	\$199
C-tree database (source)	\$399
AMIGA, CP/M-68k, 68k UNIX	call

#### Apple II, Commodore, 65xx, 65C02 ROM

#### Manx Aztec C65

"The AZTEC C system is one of the finest software packages I have seen"

NIBBLE review, July 1984

A vast amount of business, consumer, and educational software is implemented in Manx Aztec C65. The quality and comprehensiveness of this system is competitive with 16 bit C systems. The system includes a full optimized C compiler, 6502 assembler, linkage editor, UNIX library, screen and graphics libraries, shell, and much more. The Apple II version runs under DOS 3.3, and ProDOS, Cross versions are available.

The Aztec C65-c/128 Commodore system runs under the C128 CP/M environment and generates programs for the C64, C128, and CP/M environments. Call for prices and availability of Apprentice, Personal and Developer versions for the Commodore 64 and 128 machines.

Aztec C65-c ProDOS & DOS 3.3 \$399
Aztec C65-d Apple DOS 3.3 \$199
Aztec C65-p Apple Personal system \$99
Aztec C65-a for learning C \$49
Aztec C65-c/128 C64, C128, CP/M \$399

#### Distribution of Manx Aztec C

In the USA, Manx Software Systems is the sole and exclusive distributor of Aztec C. Any telephone or mail order sales other than through Manx are unauthorized.

#### **Manx Cross Development Systems**

Cross developed programs are edited, compiled, assembled, and linked on one machine (the HOST) and transferred to another machine (the TARGET) for execution. This method is useful where the target machine is slower or more limited than the HOST, Manx cross compilers are used heavily to develop software for business, consumer, scientific, industrial, research, and educational applications.

HOSTS: VAX UNIX (\$3000), PDP-11 UNIX (\$2000), MS-DOS (\$750), CP/M (\$750), MACINTOSH (\$750), CP/M-68k (\$750), XENIX (\$750).

TARGETS: MS-DOS, CP/M-86, Macintosh, CP/M-68k, CP/M-80, TRS-80 3 & 4, Apple II, Commodore C64, 8086/80x86 ROM, 68xxx ROM, 8080/8085/Z80 ROM, 65xx ROM.

The first TARGET is included in the price of the HOST system. Additional TARGETS are \$300 to \$500 (non VAX) or \$1000 (VAX).

Call Manx for information on cross development to the  $68000,\ 65816,\ Amiga,\ C128,\ CP/M-68K,\ VRTX,\ and others.$ 

#### CP/M, Radio Shack, 8080/8085/Z80 ROM

#### Manx Aztec CII

"Tve had a lot of experience with different C compilers, but the Aztec C80 Compiler and Professional Development System is the best I've seen."

80-Micro, December, 1984, John B. Harrell III

Aztec C II-c (CP/M & ROM)	\$349
Aztec C II-d (CP/M)	\$199
C-tree database (source)	\$399
Aztec C80-c (TRS-80 3 & 4)	\$299
Aztec C80-d (TRS-80 3 & 4)	\$199

#### How To Become an Aztec C User

To become an Aztec C user call 1-800-221-0440 or call 1-800-832-9273 (800-TEC WARE). In NJ or outside the USA call 201-530-7997. Orders can also be telexed to 4995812.

Payment can be by check, COD, American Express, VISA, Master Card, or Net 30 to qualified customers.

Orders can also be mailed to Manx Software Systems, Box 55, Shrewsbury, NJ 07701.

#### **How To Get More Information**

To get more information on Manx Aztec C and related products, call 1-800-221-0440, or 201-530-7997, or write to Manx Software Systems.

#### 30 Day Guarantee

Any Manx Aztec C development system can be returned within 30 days for a refund if it fails to meet your needs. The only restrictions are that the original purchase must be directly from Manx, shipped within the USA, and the package must be in resalable condition. Returned items must be received by Manx within 30 days. A small restocking fee may be required.

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#### **LETTERS**



#### **Blum Signs Off**

Dear DDJ,

Being a part of *DDJ* for practically three years has been an incredible experience. Even though I won't be writing a regular column anymore, I don't feel that I am losing a thing. The many new friendships that I now enjoy and the experiences that I shared will last a lifetime.

Of particular importance to me has been your feedback. The many letters I received, even when critical, were a delight to read and helped guide me in preparing material meaningful to you.

Even though I won't be a regular participant in future issues of *DDJ*, I have a number of CP/M projects under way that I hope will be accepted for publication in the future.

In closing for the last time, allow me to invite you to continue to visit my bulletin board system and to again thank each of you for looking in on the CP/M Exchange each month.

Bob Blum 5536 Colbert Trail Norcross, GA 30092

Bob Blum's RCP/M system is available for your use 24 hours a day, 7 days a week. Reach it by dialing (404) 449-6588.

Dear DDJ,

I have been searching (with no success) for the address and subscription cost for *DTACK Grounded*. Would you or any of your readers have this information?

Thank you for your time.

Calvin Dodge 4490 Yukon Ct., #2A Wheatridge, CO 80033

DTACK Grounded, one of the best newsletters on programming and the industry, ceased publication with Issue 45 to allow its editor, Hal Hardenberg, to devote more time to software development. His company, Digital Acoustics, is healthy and is mailing subscription refunds to subscribers. Back issues are still available from Digital Acoustics, 1415 E. McFadden, Ste. F, Santa Ana, CA 92705.—ed.

#### **Editors**

Dear DDJ.

Mark Edwards is to be congratulated for his review of editors in the November 1985 issue of *Dr. Dobb's*. Anyone who ventures into this highly personalized field takes his life in his hands, and to attempt a survey of ten editors is

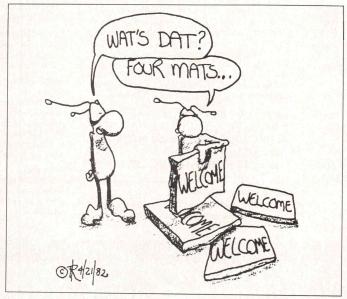
chutzpah indeed. To do it with taste, relative completeness, and fairness is no mean task.

As Mr. Edwards notes, EC is in a high degree of flux, but I would note that my version (probably later than his, given the lead time it requires for an article) can do backward searching and does not suffer the poor error handling he referred to. An important aspect of any software is that of support, and I would just like to comment that I have found the authors of EC superlative in that aspect.

A very nice feature of EC (among many others) is its maintenance of the DOS commands in a buffer. In all my editing I use APX Core, which offers multiple DOS partitions. EC did work with APX but lost the ability to maintain the DOS command buffer under it. I called Gene Brown, the principal author of EC, and informed him of this and also gave him the address of APX. Much to my surprise and delight, I received three (count them, 3) copies of the final version. One was the standard, and the other two were potential solutions to the APX compatibility problem. They all worked, and the solutions did so beautifully. How many companies will give such a response?

I have an additional comment with respect to the relevance of the compatibility of editors with such "enhancers" as Side-Kick. My experience has been that in the long run the side effects of these programs are less desirable than their offerings. Even more to the point, though, XyWrite, BRIEF, and EC (and probably others in the review also) are sufficiently internally complete that you really do not need such enhancers. BRIEF, with its powerful macro capability, allows you to create these enhancements for the most part, as does XyWrite for those it does not have; EC has most of them built-in, and the few it doesn't, you can create. Communications are accessible by any of these editors because you can exit to DOS, so you have the advantage of no overhead and no potential side effects. With BRIEF, I even wrote a perpetual calendar to mimic that aspect of SideKick.

Because people tend to use editors for short correspondence, a relatively important aspect is their "reformatting" speed as most editors do not (understandably) maintain word wrap for text insertion after a line has been entered. One of the problems I found with EDIX was its intolerably slow reformatting—in



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(Continued from page 8)

this respect VEDIT and EC are lightning fast, and BRIEF (with its free reformat macro) is only slightly faster than EDIX. XyWrite of course does it like any respectable word processor.

enjoyed reading "Wired Tales" in Dave Cortesi's column in the December issue and offer the following. I recently decided to upgrade my PC with a larger power supply (64 watts to 135 watts) and, being relatively frugal, ordered one from a mail-order house for \$89. When it arrived I (foolishly) started to install it: removed my old power supply, put the new one in, and started to make the connections. To my dismay, it turned out that the connectors were poorly constructed and I could not get a good fit. I decided to return the supply for credit, not wanting to chance a repeat performance. It took me about \$15 worth of phone calls to finally talk to someone who was authorized to issue an RA, and that same person informed me that there was a 10 percent restocking fee. Adding in the cost of the UPS shipment back, I wound up paying \$30 for some futile labor in taking out and replacing my power supply, losing a nontrivial amount of time in the process.

Finally, I offer a poser. I recently purchased an uninterruptible power supply (300 watts) from Qubie. When it was delivered, I tested it in the usual way-I pulled the power plugand it worked fine. Just to be doubly sure, I also opened the main circuit breaker in the house, and the same excellent performance repeated. A few days ago, we had our first real power outage while I was working at my PC, and much to my surprise, I lost all my work. When the power returned about one minute later, I immediately made the "pull the plug" test and it worked. I called Qubie's technical department, and no one could offer a solution, but the company did offer to send me a replacement (which I accepted). In the meantime I have been thinking about the possible causes of the loss of my work and believe I may understand why it happened, but I would love to hear from others.

Morton F. Kaplon 11 White Birch Dr. Pomona, NY 10970

Dear DDJ,

Although I greatly enjoyed Mark Edwards' editor review (DDJ, November 1985). I would like to make a correction or two and point out a problem with this kind of review. I think Mr. Edwards did a terrific job given the difficulty of trying to learn about ten text editors, much less trying to comment intelligently about them all. I use Pmate (Version 3.37) on a daily basis, however, and I am familiar enough with it to have caught a couple of errors in the review.

The first (and smallest) is that Pmate does indeed allow you to undo the deletion of a single character. What it will not allow is the retrieval of a character you have backspaced away. On a PC, this is the difference between the Delete key and the ← (backspace) key. The difference can be annoying, but it is frequently more helpful than not once you are aware that Pmate works in this way.

A more serious error is in Mr. Edwards' Pmate macro to count braces. The macro he presents does indeed work, but it is terribly slow. Trying it on a sample C program out of one of my working directories (on a Compag Plus), it ran for more than six minutes before I aborted it-unfinished. There is a much better Pmate macro for brace counting, which I present in Table 1, page 10. This brace-counting macro finished the same piece of code as above in less than six seconds (and is slowed slightly because it is constructed to work with a file

of any size as opposed to one entirely in memory). This is a huge difference from the results reported in the review. Because of differences in the test files, it is impossible to say whether this performance is better or worse than BRIEF or EMACS, but it is certainly of the same order.

It is worth noting that the macro presented in Table 1 works in much the same manner as the BRIEF and EMACS macros presented by Mr. Edwards. In particular, it allows the editor itself to search for the characters to be counted rather than stepping through the file in a clumsy characterby-character fashion. I suspect that a properly rewritten macro for VEDIT PLUS would show the same kind of dramatic speed improvement. There might even be some hope for XTC if the proper macro formulation were used.

This in turn points out a serious difficulty in benchmarking editors. A user who is familiar with an editor will undoubtedly be able to generate a better benchmark than one who isn't. I would never have even considered writing a brace-counting program for Pmate in the way Mr. Edwards did. The way Pmate works guarantees that his macro is among the slowest possible for the task.

Based on my experience, I would tend to agree with Mr. Edwards' comments about the functionality and features of each editor he reviewed. Based on my findings for Pmate, though, I would be inclined to disregard the entire benchmark table except for items A, B, F, and maybe C.

Mr. Edwards has certainly proved that it is im-

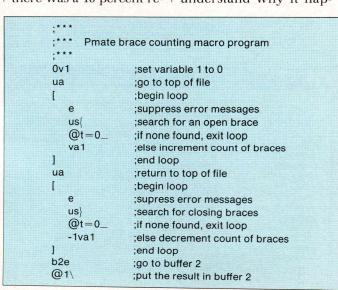


Table 1

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(Continued from page 10)

possible for a single person to learn enough about ten different editors to be able to compare each at its best.

Brad Chase P.O. Box 705 Exeter, NH 03883

#### Modula-2

Dear DDJ,

While perusing your otherwise excellent November issue, I felt a sudden stab of pain. Looking closer, I saw the source of my discomfort was "Bit Manipulation in Modula-2." I feel an overwhelming need to comment.

"I knew from the outset that it was impossible to match C's simplicity." This tells me that the author is primarily a C hacker. Modula-2 is perfectly capable of matching this "simplicity," and in exactly the same manner. (See Table 2, below.)

The other operations are left as exercises for the reader.

Lawrence C. Smith 51 Lake St. Nashua, NH 03060

#### **C** Compilers

Dear DDJ,

I love *Dr. Dobb's* and have subscribed for years. I notice a trend away from assembly, CP/M, 8-bit, and the less wealthy user, though. In particular, your issue on C language was a

disappointment. The following are my reasons, along with some suggestions. I understand your problems in appealing to a divergent group, however, and appreciate your willingness to listen to all our suggestions.

- 1. The main article on C comparisons told me more than I really cared to know about the languages. I would have trusted the writers if they had summarized their results. An article this long and detailed should have a summary at the beginning.
- 2. The article did not state which of the languages examined were available on either S-100 machines, on CP/M, or on 8-bit machines.
- 3. I did not see mention of several of the Cs I use—BDS C, Small-C, or Tiny C or any mention of Lifeline's products—and that would have been important to me. I am sure other Cs could have been found as well—so how much longer would it have taken to wait a bit and do a truly definitive review?

Frederic Schlamp 2205 Meadowview Rd. Sacramento, CA 95832

The comparative review of C compilers in our August 1985 issue was strictly dedicated to MS DOS-compatible compilers. We will review C compilers for other environments, including CP/M,

in 1986.—ed.

#### Columns

Dear DDJ,

In your July 1985 issue the 16-Bit Software Toolbox column began with the line "One of the most novel features added in Version 2 of MS DOS is the concept of 'installable device drivers.' "I would like to say that this concept may be new and novel for Microsoft and MS DOS, but it is certainly not a new and novel concept for other operating systems available for microprocessors.

The OS-9 operating system has had the concept of user-installable device drivers since its initial 6809 Level 1 release in 1978. In fact, OS-9 is totally modular in nature and allows the user to add new device descriptors, device drivers, and new file managers if required. In addition to supporting installable drivers, OS-9 has included the "novel" MS DOS concepts of hierarchical director structures and pipes. OS-9 also gives full support to I/O redirection, multiprocessing, and multitasking-concepts much more akin to Unix.

OS-9 may not be as well known as MS DOS is, but it does have a large and rapidly growing following in the 6809 and 68XXX world today. MS DOS has added nothing novel to its OS; it is adding features that are expected and required for an OS in today's world. These features have been around in OS-9 and OS-9/68000 for some time.

Tim Harris Microware Systems Corp. 1866 N.W. 114th St. Des Moines, IA 50322

Dear DDJ,
"It isn't what you don't know that hurts you, it's

what you know that ain't so." In the September Of Interest column, the author states that APL on the IBM PC requires an 8087. This is only true for IBM's own APL, not for the other five (STSC, Sharp, Portable Software, Watcom, and NIAL Systems). Most of them run on the PCjr.

Edward M. Cherlin 6611 Linville Dr. Weed, CA 96094

Dear DDJ,

I was interested to read the portion of your Dr. Dobb's Clinic (October 1985) regarding the intricasies of manipulating path names. Some time ago, we at POLY-TRON also had the pleasure of figuring out how to do exactly what you discussed. Our solution, which produced some functions that are more general purpose, might interest your readers. The primary components of our solution are:

1. A function to determine if a given path name represents a directory—that is, it is a drive ID or an actual director but not a normal file:

int is\_dir(name) char \*name;

2. A function to determine if a given file name contains wildcard characters:

int is\_wild(name) char \*name;

3. A general-purpose function for generating a new file name given an existing file name, an optional new extension, and an optional new path:

char \*
bld\_fnam(pathp,
 namep, extp)
char \*pathp;

```
(* assume ch:=CHAR(255) *)
(* clear bit 7 *);
ch:=CHAR(BOOLEAN(ch) & BOOLEAN(07FH));
(* result is ch=CHAR(127) *)
(* assume ch:=CHAR(15) *)
(* set bit 5 *)
ch:=CHAR(BOOLEAN(ch) OR BOOLEAN(10H));
(* result is ch=CHAR(63) *)
```

Table 2

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(Continued from page 12)

/\* the new path \*/ char \*namep;

/\* the original name \*/
char \*extp;

/\* the new extension \*/

This returns a pointer to allocated memory containing the newly constructed file name. The original name may or may not have a path or extension. If either pathp or extp point to a null string, the respective component of the original name is omitted from the new name. If either pathp or extp is a null pointer, the respective component of the original is used in the new name.

4. A wildcard file name expansion function with optional ability to prepend the path of the wildcard name to each matching name is shown in Table 3 (below). This returns a pointer to a linked list of names matching the wildcard name. Each node of the list is contained in sep-

arately allocated memory and is of sufficient length to contain the matching name beginning at name[0].

5. Functions to extract each of the three components of a file name—the path, the root, and the extension are shown in Table 4 (below). These functions each return a pointer to allocated memory containing the respective component of a file name.

Using these functions, you can do just about anything you wish with file names. In fact, they are used in nearly all our products and continue to be used in new development work.

These functions, along with a host of others, are available in the POLYTRON C Library 1. This package contains routines that function under DOS 1.x, DOS 2.x and later, and both/either. The functions are provided in linkable form (libraries), as well as in full source form (C and assembler) for the IBM PC/

XT/AT and compatibles.
Donald K. Kinzer
POLYTRON Corp.
P.O. Box 787
Hillsboro, OR 97123

Dear DDJ,

Here is a report on a software company for the benefit of prospective buyers. I purchased a cross-assembler from 2500AD in November 1984. What they sent me was a nonfunctioning program. Soon after receiving the program, I sent two letters, made two phone calls, and finally spoke to someone who said the company would send me a good copy when it had fixed the program. It has been a year, the firm still has my \$200, and I have not yet received a functional copy of the program. Even JRT Pascal was functional, and look at the price difference!

For the record, here are the defects that I have found (so far):

- 1. The assembler and the linker are incompatible. The linker makes the wrong assumptions about the relative order of least/most significant bytes, which makes it impossible to link object modules. The only way to get an executable file is to put all your code into one module and use absolute addressing (with an ORG statement).
- 2. The assembler gives the wrong machine code for several of the opcodes. This makes for difficult debugging.
- 3. Some of the pseudoops listed in the manual don't exist in the actual assembler.
- 4. Some of the pseudoops that do exist do not work (at least one).
- 5. The assembler has no provision for allowing the programmer to specify the short forms of the relative

jump and call instructions.

- 6. The symbol table that is printed at the end of the listing usually contains sections of garbled mess.
- 7. The printed symbol table lists all intermediate values of labels that are redefined many times with the DEFL statement.
- 8. Some errors in the source code cause the assembler to crash with no error message.
- 9. The linker crashes with a nonsense error message under some conditions that seem to have something to do with certain exact lengths of files.

Neil R. Koozer Kellogg Star Rt. Box 125 Oakland, OR 97462

In the September issue, we ran a comparative review of PC TEX and MicroTEX. At that time only PC TEX included INITEX, but we stated that MicroTEX was also scheduled to include INITEX in an August update. As of the time this issue went to press, MicroTEX still does not include full INITEX capabilities. We have been informed by Addison-Wesley that a new version of MicroTEX, including INI-TEX, will not be available until January 1, 1986. Check with Addison-Wesley before placing an order. Both Addison-Wesley and Personal TEX are already distributing the Textset laser printer and screen preview driver.-ed.

DDJ

```
/* element of a linked list of matching names */
struct file_list {
    struct file_list *next; /* ptr to next node */
    char name[1]; /* a file name */
};
struct file_list *
expand(wildname, add_path)
char *wildname; /* the name to match */
int add_path; /* non_zero if path of wildname should be
    prepended to matching names */
```

#### Table 3

```
char *
path_of(name)
char *name;
char *
root_of(name)
char *name;
/* the name from which to extract the path */
char *
ext_of(name)
char *name;
/* the name from which to extract the root */
ext_of(name)
char *name;
/* the name from which to extract the
extension */
```

#### Table 4

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#### VIEWPOINT

#### Inefficient C

This column is adapted from DTACK Grounded, The Journal of Simple 68000/32801 Systems, Issue 42.

I thas been apparent to me for two or three years now that complex programs (as opposed to the famous but simplistic "hello world" type) written in C consistently run a great deal slower than the same complex programs written in assembly. Examples of this rule can readily be found in the personal computer mass marketplace.

It is true that, in the last year or two, many intelligent and experienced programmers have asserted that C has little or no highlevel language (HLL) overhead.

There is an apparent conflict here.

Consider the complex problem of writing a BASIC interpreter. You can break this problem down into a large number of simple problems that can then be solved in assembly or, say, C. The speed with which each of the simple problems can be solved depends on how good a match can be made between the problem and the available control constructs, data types, and so on. In some cases the match with the constructs and data types available in

#### by Hal Hardenberg

C is as good as the best that can be devised using assembly, and you have zero HLL overhead. In fact, many simple problems can be solved as quickly (or nearly so) in C as in assembly.

So, the C supporters have no shortage of real-world examples of simple problems that can be solved as quickly (or nearly so) in C as in assembly, which leads them to claim that C has no high-level overhead.

Some problems, though, can be solved easily and quickly in assembly and are real bears in C, given the restricted range of operations, constructs, and data types available to C programmers. The large number of simple problems that comprise a BASIC interpreter will include some C-easy problems, some C-not-so-easy problems, and some C-bears. Therefore, any full-function BASIC interpreter written in C will always be a great deal slower than the same interpreter written in assembly.

You need not look for complex ways to analyze the HLL overhead of C (or any high-level language). The fact is, all high-level languages greatly restrict the range of operations, control constructs, and data types available, compared to assembly language. Thus, HLL programmers have a limited number of tools available for solving the large number of simple problems that comprise any solvable complex problem.

Obviously, the programmer with the most complete set of tools can always produce the fastest code. Real-world evidence such as Microsoft's MBASIC (written in assembly) vs. DRI's Personal BASIC (written in C) confirms this rule.

Why, then, do intelligent, experienced programmers claim that C has little or no HLL overhead when abundant real-world evidence contradicts such an assertion?

Well, suppose you have solved a complex problem using C. Then, keeping the 90-10 rule in mind, you go looking for ways to speed up the software. Do you rewrite the program from scratch in assembly, using the greater variety of available operations, data structures, and control constructs? No, you continue to use the data structures and algorithms that you developed in C. That is, you implicitly restrict yourself to the smaller toolkit that is already used by C! Not surprisingly, the resultant program is not much faster than the original. Hence you announce, "Hey, guys, I tried optimizing the problem using assembly and got little or no improvement over the original C. Obviously C has little or no high-level overhead.'

Obviously.

HLL inefficiency is acceptable in many environments. Where HLL inefficiency is not acceptable is in the personal computer mass marketplace wherever an efficient alternative is available, DRI's Personal BASIC is essentially a dead issue, being highly uncompetitive with Microsoft's assembly-based BASICs in terms of speed. In turn, the marketplace is widely avoiding Microsoft's Cwritten FORTRAN compiler.

The individuals who make up the mass marketplace and who vote with their wallets don't care how hard it was to produce a program or how long it took. They don't care about source documentation or maintainability because they never get to see the source documentation and somebody else maintains the code. They vote for those programs that work both swiftly and well (Lotus 1-2-3, original WordStar, MBASIC, or GW-BASIC) rather than those that run well but slowly (Context MBA, WordStar 2000, or DRI Personal BASIC).

The folks who assert that C has little or no HLL overhead are either talking about simple problems for which C is in fact well suited, or they are observing that, if algorithms written in C are duplicated in assembly, then the assembly versions have little advantage. To reason that C therefore has little HLL overhead is faulty logic.

Other things being equal, the more complete toolbox is always going to win provided the workman knows how to use those tools. That's why good, experienced assembly programmers cost more than HLL coders.

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#### C CHEST

#### A Unix-like Shell for MS DOS

ast month we presented a few support routines, part of a Unix-like shell for MS DOS. This month we're going to continue with the shell itself, describing how it works on a high level. Next month we'll look at it at a lower, functional level. The code itself, because it's so long, will be spread over the next three months.

The shell described here includes functions of the Unix C shell that I use most often, such as aliases, history, and command-line wildcard expansion (more on all this in a moment). It has batch-file capability and will permit nested batch-file execution (unlike MS DOS, which lets you chain from one batch file to another but won't let you return to the original batch file).

It supports a 2,048-byte command line with interactive editing. The long command line is passed to an executed program through an environment variable.

#### **Using the Shell**

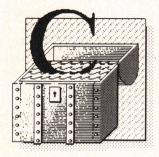
Commands are entered from the command line, just as in DOS. (Note that \ is a special character to the shell, so use slash (/) or \ \ to separate directory names.) DOS wildcard characters (\* and ?) are expanded before a command is executed. So if you say echo \*.c, the \*.c will be expanded to the names of all files having a .c extension before echo is invoked. Expanded names are sorted. Several semicolon-delimited commands may be executed from a single command line. For example:

cd foo; pwd; ls

#### by Allen Holub

changes the current directory to foo, prints the full path name, and then prints a list of the files in the current directory.

Command-line editing (as de-



scribed last month) is supported. To summarize:

Cursors—moves the cursor.

Home—gets to the beginning of the

End—gets to the end of the line.

Ctrl-right arrow and Ctrl-left arrow—get to the next and the previous word, respectively.

^H—is a destructive backspace.

Del—deletes the character under the cursor. (Typed characters will be inserted at the current cursor location, moving all characters to the left of the cursor over one notch)

^X—deletes the entire line.

Esc—does the same and aborts.

Return—causes the commands to be executed

There are several built-in commands:

alias—creates, modifies, or prints aliases (see below). There are two syntaxes:

> alias alias name <val>

The first prints all currently defined aliases, and the second creates an alias for *name* with the indicated value. *<Val>* may be anything on the command line, but you have to escape (precede with a \) any character that's special to the shell (or surround *<val>* with double quotes).

cd—changes a directory or disk: cd foo—changes to the subdirectory foo. cd..—changes to the parent directory.

cd a:—changes to the current directory on the a: drive.

cd a:/foo—changes to the /foo directory on the a: drive.

Cd must be used to change disks, although you can alias a: to cd a: if you like. The shell checks to see if a disk is in the indicated drive before the drive is logged on.

exit—terminates the shell. Either *exit* or *logout* must be used to leave the outermost shell. Subshells can be terminated with a \*C.

history—prints the history list (see below).

logout—like *exit*, terminates the shell; however, the file /logout-bat is executed before exiting.

pwd—prints out the current working directory (same as *chdir* with no arguments under DOS).

rem-does nothing. May take arguments (which will be ignored). This command is here only for DOS compatibility. The preferable method for commenting a batch file is to start comment lines with a # in the far left column. Note that rem is interpreted as if it were a command; that is, the line is put into the history list. Moreover, if a line starts with a rem but has a semicolon on it as well, text up to the semicolon will be ignored but the semicolon will be treated normally, and any text following the semicolon will be treated as a second command and executed.

set—creates, modifies, or prints a shell variable (see below). Its syntax is:

set name [=][value]

The first form prints all current shell variables, and the second creates or modifies an existing variable. Both the equals sign and the value fields are optional. If you omit the value, the alias will expand to a null string.

setenv—changes or creates an environment variable. Its syntax is:

setenv name [=][value]

Both the equals sign and the value field are optional. An example:

setenv PROMPT \$p>

sets the prompt to the current directory name followed by >. This command is very similar to the DOS set command. However, setenv with no arguments doesn't print the environment. Also, the equals sign is optional.

shift—shifts all the \$<num> arguments in a batch file left one notch. For example, if a batch file (foo.bat) consisting of:

echo \$0 \$1 \$2 shift echo \$0 \$1 \$2 shift echo \$0 \$1 \$2

is executed with the line:

foo first second

the following output will be sent to the screen:

echo first second first second second

(\$0 expands to the file name.) unalias—removes an alias. The syntax is unalias name.

unset—removes a shell variable. The syntax is *unset name*. The default variables (*verbose*, *echo*, *arg*) can't be removed.

#### History

The history functions let you examine, edit, and reexecute previous commands. The 32 most recently typed commands are remembered in a history list, and each command has a history number associated with it. Once the history list is full, the oldest command is discarded every time a new command is added to the list. The easiest way to see how this mechanism works is with an example. Assume that you've typed the

following five commands:

vi prog.c
cc -c prog.c > err
vi another.c
cc -c another.c > err
link another+prog,prog,prog,c.
lib+/lib/tools.lib

As you type the commands, they're entered into the history list, which can be examined by typing *history*. The following will be printed:

1: vi prog.c

2: cc -c prog.c >err

3: vi another.c

4: cc - c another.c > err

5: link another+prog,,,c.lib+/lib/ tools.lib

6: history

Note that *history* itself is also added to the history list. The numbers are the history numbers associated with each command. A command can be executed again by typing !#, where # is the history number. Typing !5 causes the link command to be executed again. !2 will reexecute the first



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#### C CHEST:

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cc command. You can also type !<pat>. In this case, the history list will be searched backward for a command starting with <pat>. So !! or !link would also redo the link. !cc would be the same as !4 (because the first matching command is used). !! will repeat the last command you typed. In the above example, !!, !6, and !h will all do the same thing. Commands are added to the history list every time they're executed, even if a command is an expansion of a history request.

Several additional non-Unix history commands are supported. !>file will write the current history list to the specified file. If no file is given, /histlist is used. The complementary command is !<file, which adds the commands in file to the current history list. The commands aren't executed, they're just added to the list. Again, if file isn't specified, /histlist is used. Neither !> nor !< will show up in the history list (they won't use up a history number,

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^ may replace !; that is, the commands ^^, ^#, and ^<pat> may be used in place of !!, !#, and !<pat>. The ^commands work just like the! commands except that the line is brought into an edit buffer that you can then manipulate in the normal way (with the cursors, etc.) before executing it. Unlike DOS, the command line is visible while you're editing it. Note that Esc will abort out of edit mode without executing the edited command line.

#### **Environments. Shell Variables.** and the Set Command

Shell variables are macros. They let you associate a body of text with a name, and when that name is used, the corresponding text is substituted. Shell variables are created with the set command and deleted with the unset command. They work something like Unix and DOS environments except that they can't be passed to a child process. Once a shell variable is created, it can be used anywhere in a command. For example, you can define a shell variable to represent a long directory spec with:

set HOME /usr/allen/src/shell

You can then use it on the command line.

cd \$HOME

Cd \$HOME will be expanded to cd /usr/allen/src/shell before the shell executes the line. Note that \$ must precede all uses of shell variable names but must not be in the definitions; % may be used instead of \$ if you prefer. There are several predefined shell variables (which can't be modified with the set command). These are:

- \$<num>—an argument to a batch file (\$0 is the file name, \$1 the first argument, etc.).
- \$\*-expands to all \$<num> variables concatenated together.
- \$p-expands to the current path name.
- \$!—expands to the current history number.
- \$s-expands to the current shell level.



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#### C CHEST:

(Continued from page 20)

The top-level shell is 0. If you create a shell within a shell, the second one will be at level 1. All batch files are executed in their own shells.

There are three other shell variables that can be modified with *set* but can't be expanded with *\$name*. These are *echo*, *verbose*, and *cmd*. If *echo* is set *(with either a set echo or a set echo = 1)*, then commands will be echoed to standard output just before they're executed. The default is *echo* 

off (unlike MS DOS), so you must set echo inside a batch file if you want to see what the batch file is doing. The echoed line will show all macro substitutions and all wildcard expansions; however, it will be truncated to the 127 characters permitted by DOS. You can use this feature to see if the command line has been truncated when you're invoking a program that doesn't know about the CMDLINE environment (see below).

Verbose shows the input as the shell receives it (before it's interpreted). If cmd is set (the default), then an

environment variable called CMDLINE will be created every time a file is executed. CMDLINE holds the entire 2,048-byte command line (which can't be passed via DOS). If cmd is cleared, then CMDLINE is still created but will have no contents. Echo, verbose, and cmd can be cleared with set < name > = 0. Note that setting echo or verbose has the same effect as specifying -x or -v on the command line used to invoke the shell.

Environment variables (or strings) are similar to shell variables except that they can be passed to a child process (a pointer to them is included in a child's PSP). Many compilers (at least the Aztec, Lattice, and Microsoft) have a getenv(name) function in their libraries that returns a pointer to the environment string corresponding to name. If you need to write your own getenv(), a good description of the PSP can be found in The Peter Norton Programmer's Guide to the Norton IBM PC (Bellevue, Wash:: Microsoft Press, 1985), pp. 260f.

Environment variables can be set from within the shell with the *setenv* command. Unlike DOS, they can be used on the command line just like a shell variable (precede the name with \$ or \$\%).

#### **Special Characters**

- \* and ?—have the same significance as in MS DOS. They are expanded by the shell to matching file names. Expanded names are sorted. For example, echo \*.c \*.obj would print a list of all the .c files in the current directory (sorted) followed by all the .obj files (also sorted).
- ;—is used to separate multiple commands on a single command line. Cd foo;pwd would change the current directory to foo and then execute the pwd command.
- \—is used to take away the significance of a special character. For example, \\* can be used to pass an asterisk to grep (the \* won't be expanded). \; can be used to define a multiple command alias (see below). \\ evaluates to a backslash. The \ will be stripped from the line before the line is passed to the child process.

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(") or single (') quotes won't be modified (wildcard characters aren't expanded, semicolons aren't interpreted as command delimiters, etc.). The quotes aren't removed unless the -q argument is given on the command line. Unlike Unix, there's no distinction made between single and double quotes.

#—when found in the far left column, signifies a comment. The remainder of the line is ignored, and the line isn't put into the history list.

#### Aliases

Aliases are another sort of shellmaintained macro. Unlike shell variables, a \$ is not needed to expand the name; rather the alias is expanded if its name is found as either the first word on a line or the first word following a semicolon on a multiplecommand line. Aliases have two uses: They can be used to change the name of a command and they can be used in place of batch files. Aliases are created with the alias command. Some examples: If you have a program called ls that prints the current directory but you also want to be able to type dir and get a directory listing, you can define an alias for dir as follows:

alias dir ls

Thereafter, when the shell finds dir as the first word on a line, it will substitute the string ls for the string dir, and the program ls will be executed. Only the first word of a command is modified so dir foo bar rat will be changed to ls foo bar rat. Aliases must be either the first word on a line or the first word following the semicolon when there are several commands on one line.

Aliases can also be used in place of batch files, provided that no arguments need to be expanded. Because aliases are memory resident, they will execute much faster than a batch file. Similarly, an alias doesn't execute under its own shell, as does a batch file, so much less core is needed to execute an alias than is needed to execute a batch file. A simple alias that is similar to a batch file is:

alias shell cd /src/util/shell

Now you can type the single word *shell*, which the shell will expand to *cd /src/util/shell* and then execute to move to the indicated directory; however, you can't expand arguments. The portion of the command line that follows the alias name will be concatenated to the end of the expanded alias. Let's look at an example: Polymake lets you specify a default rules file on the command line with a —B option, but you often have to make a target name as well. By defining the alias *m* with:

alias m make -B /lib/

you can then type *m foo.exe* and the shell will expand it to:

make -B /lib/builtins.mak

foo.exe

Because you can have multiple semicolon-delimited commands on one line, you can define an alias that will expand to several commands. For example:

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#### C CHEST:

(Continued from page 23)

alias m rm err \; make −B /lib/ builtins.mak

will expand to

rm err; make -B/lib/builtins.mak

thereby both removing the file err and executing make. Note that you must escape the semicolon (with a \) in the alias definition so that the shell won't intercept it. (You could also surround the definition with quotes.)

Aliases may use shell variables. Two such aliases are:

alias here set here = \$palias there cd \\$here

Here will set the shell variable here to the current directory. There will put you in the directory remembered with a previous here invocation. Note that \$ has to be escaped to prevent the shell from expanding it when the alias is defined.

A caveat about aliases: Aliases defined in terms of other aliases won't work. For example:

alias foo echo foo alias bar echo bar alias foobar foo;bar

won't work (foobar will expand to foo; echo bar). However, the command line foo;bar will work. Also note that commands are added to the history list before aliases are expanded.

#### Environments and Files

The default command line prompt is (\$s,\$!) (which prints the current shell level followed by the current history number). You can specify a different prompt with the PROMPT environment variable (use either the DOS set or the shell's setenv command). Any ASCII character may be used, and any of the \$ arguments will be expanded before the prompt is printed. For example:

setenv PROMPT \$p->

will change the prompt to the current directory name followed by and >. The prompt can be changed at any time (it doesn't have to be set when the shell boots). The SWITCHAR environment variable tells the shell what character signifies a commandline switch when it's the first character in a command-line argument (the default is -). Because environments are inherited by the child process, SWITCHAR is also available to a program if it chooses to use it.

The CMDLINE environment holds the complete, 2,048-byte command line (which can't be passed via DOS). It is changed every time the shell executes a command. Note that the command line (truncated to 127 characters) is also passed to a child process in the normal way (via the DOS command line buffer at offset 0x80 from the child's initial code segment). When the shell spawns a subshell to execute a batch file, it uses CMDLINE.

The SHLEV environment is set to the current shell level. The outermost shell is at level 0. All shells created from within another shell (including those used to execute batch files) will have higher numbers, de-

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pending on the level of nesting.

Several files are used by the shell. If it exists, the file /shrc.bat is executed (in a manner analogous to autoexec.bat) every time a shell is created. Because batch files are executed in their own shells, /shrc.bat will be executed every time a batch file is executed. A second file /login.bat is executed only once, when the lowest-level (level 0) interactive shell is created. Shrc.bat is executed before login.bat, and both files are executed before the environment is examined. My own login file is shown in Table 1, at right. My logout file is simply !> and lets me leave the shell without losing the current history list. Note that !< in login.bat lets me reenter the shell without losing the list. You could also use !< (without the !>) to read in a list of commonly used commands when the shell boots.

#### **Shell Invocation Syntax**

There are several ways to get into the shell from the command line. The easiest way is to type *sh* (with no arguments), putting you into *interactive* mode. You can't get out of the lowest-level interactive shell with a °C. Use either *exit* or *logout*.

- sh —c string—invokes the shell in nonresident mode. It will execute the command contained in <string> as if it had been entered from the command line and then terminate.
- sh filename args...—executes a batch file. \$0, if found inside the batch file, will be expanded to the *file name*. The arguments can be fetched with \$1, \$2, etc. % can be used instead of \$ if you like.

Four other command-line switches are available:

- -i—puts the shell into interactive mode even if arguments are listed on the command line. Normally, if command-line arguments are present and -c isn't specified, the shell will try to execute a batch file. sh i  $arg \dots arg$  will create an interactive shell, putting the arguments into \$1, \$2, etc. \$0 will hold the string -i.
- q—causes quotes to be stripped from commands before they're passed to a child process. The

```
setenv PROMPT=[\$s:\$!]
alias a alias
a h history
a book
                 cd /text/book
a ddi
                 cd /text/ddj
                 set here = \$p
a here
                 cd /src/class
a sclass
                 cd \$here
a there
                 cd /lib/tmac
a tmac
                 cd /src/tools
a tools
a type
                 cat
                 "rm err; make -B builtins.mak"
a m
```

Table 1: A login.bat file

quotes will still protect wildcard characters, etc., from expansion.

- -v—verbose mode, commands are echoed to stderr as they're read by the shell. This is the same as a set verbose command.
- -x—commands are echoed to stderr just before they're executed. All \$ arguments and wildcard characters will have been expanded at this point. This is the same as a set echo command.

#### Redirection

The shell itself doesn't support redirection; however, because command.com is still resident, redirection is available if you need it. There are two ways to use command.com for redirection. A nonresident shell can be invoked from DOS with a line such as:

sh -c grep pattern \*.c > foo

Grep will be executed, and the shell will expand the \*.c to the names of all files in the current directory having a .c extension before grep is invoked. Grep's output will be redirected to foo in the normal way. Because this command is executed from MS DOS and not from the shell, it won't be added to the history list.

The second method also lets you enter redirected commands into the history list. From inside the shell, type:

command /c grep pattern \*.c > foo

The /c argument to command.com works like the -c argument to sh, so it will execute the following string as if it had been typed. The \*.c will again be expanded by the shell before command.com is executed, but the > will be interpreted by com-

mand.com (which will put the output into *foo*). Be careful here of command-line truncation. Because command.com doesn't know about the CMDLINE environment, it has no way to get to the extended command line, so it will work on only the first 127 characters.

#### **Support Routines**

In order to minimize the size of the shell, I've tried not to build in commands that aren't essential. I've found the following programs to be useful.

cat.c—prints (concatenates) files to stdout.

cp.c—copies a file to another file or disk. Copies a group of files to another directory or disk.

echo.c—echoes command line to stdout.

grep.c—searches for pattern in file. ls.c—lists a directory.

mkdir.c-creates a directory.

mv.c—renames a file or moves a group of files to another directory.

printenv.c—prints the current environment.

rm.c—removes a file or group of files.

rmdir.c—removes a directory.

#### Availability

This column is part of a four-part series describing the entire shell. A reprint of all four parts along with a disk containing the listings is available for \$29.95 from *Dr. Dobb's Journal*, 2464 Embarcadero Way, Palo Alto, CA 94303. Please direct inquiries to the The Shell. Prepayment is required.

#### (Listing begins on page 84) Reader Ballot

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# PL/68K Becomes 68000 Assembly Language

by Edward K. Ream

ne day not long ago, I became embroiled in an old debate with another programmer named Charlie....

"The programming team I manage is about to start a big project," I said, "and I must decide which language to use."

"Really? Which languages are you considering?"

"C and 68000 assembly language. The product will have strong competition, and great performance is crucial, so it's reasonable to consider assembly language. On the other hand, C is so much easier to use."

"Why don't you program in C and recode in assembly language as needed?" Charlie asked.

"Of course I've considered that. It might work as far as execution speed is concerned, although I'm not sure. C doesn't let you allocate registers globally, and that's a big handicap. Speed is not the only problem, though. The code must be compact, but our C compiler produces code that is 50 percent larger than assembly language. No, there's no doubt about it—eventually the program will have to be written in assembly."

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I asked myself, suppose the program produced by the C compiler and the program produced by the assembler were semantically equivalent? Suddenly PL/68K became not just another assembly language but a new way of using C.

"Do the initial prototyping in C. That's the right way," Charlie persisted. "When the program is finished, recoding in assembly language will be much easier."

"Hmmm. I'm not convinced. Recoding is going to be expensive; we'll end up debugging the whole program twice. There

might even be pressure from higher management not to recode and come out with an inferior product."

Charlie just snorted and walked away, muttering something about assembly language being a throwback to the Dark Ages.

#### Writing in Both C and Assembly

Fortunately, my friend John overheard this conversation. John and I have worked together for 15 years, and we enjoy discussing problems that come up on the job. John laughed, "Charlie is more interested in being right about C than in solving your problem."

"You sound more sympathetic."

"Well, your choice is crucial. Which language you use determines, to a large extent, how your project will turn out."

"Yes. What bothers me most is that I've got to choose now, but I won't know until the project is almost over whether the choice was correct."

"I think I know a way around this dilemma—it's a language I invented called PL/68K."

"John, my only options are C and 68000 assembly language."

"Don't be fooled by the name. PL/68K isn't really an independent language but a way of using C to do assembly-language programming."

"John, you are not making sense!"

"Let me explain. You can think of PL/68K as being either C or assembly language—either/or. But in fact, you can run a program written in PL/68K through both the PL/68K assembler and any standard C compiler. PL/68K is both C and assembly language at the same time."

"Wait a minute. You are going much too fast," I said. "First of all, you can't possibly compile an assembly-language program with a C compiler! Assembly language doesn't look anything like C—the C compiler will spit out a thousand error messages!"

"PL/68K doesn't resemble 'traditional' assembly language. Forget what assembly language usually looks like and ask yourself, 'What are the characteristics of assembly language?' "

"Go on," I replied. "You tell me."

"First, assembly language allows full access to all machine resources—all registers, all locations in memory (including the run-time stack), all I/O ports, all privilege modes, and all machine instructions. Second, there is a one-for-one correspondence between the source code you write and the object code produced by the assembler. You always know what code a particular assembly-language construct generates; assemblers neither rearrange code nor 'optimize' code away nor add anything extraneous. Assemblers are very literal-minded. Thus, assembly language ensures zero time and space overhead."

"You're saying that assembly language gives you complete control over the machine, without a compiler getting in the way."

"Exactly. Now, suppose we say assembly language is any language that (1) allows complete access to all machine resources, (2) provides a clear correspondence between source code and object code, and (3) imposes zero time or space overhead."

#### Semantic Identity

"Hmmm," I mused. "This definition doesn't say what assembly language looks like. It could even look like C. But I still don't understand. If you run a PL/68K program through an assembler, you will get one program. If you run the same source through a C compiler, you will get a second program. The two programs are not going to do the same things—similar things, maybe, but not the same things. The fact that the source code is the same doesn't matter. To put it another way, given a result desired from a specific PL/68K program, we would still have to choose between assembling the program with the PL/68K assembler or compiling it with a C compiler."

"You've stated the problem very well," John said, "but I have discovered that it's possible to design PL/68K so that the program produced by the PL/68K assembler will work in the same way as the program produced by the C compiler."

"That sounds impossible!"

"I don't think so. Let's turn the problem around. Suppose we design PL/68K according to what might be called 'the principle of semantic equivalence.' This principle states that a program, when assembled by the PL/68K assembler, must work in the same way as when it is compiled by a standard C compiler. Now let's ask, 'What needs to be eliminated from PL/68K to guarantee semantic equivalence?' '' (See Figure 1, page 40.)

"Tell me," I said, "how much of C is left after the principle of semantic equivalence takes its toll?"

"Surprisingly, almost all of it. The preprocessor is identical to the C preprocessor. All declarations and structure statements are present. Functions do not return values but otherwise are unchanged, as are Boolean and relational operators and expressions. The biggest restriction is that arithmetic operators and expressions must be severely curtailed in order to make PL/68K expressions mean the same thing as C expressions."

"You keep talking about PL/68K being assembly language," I said. "How is it possible to produce code the quality of assembly language from a language that is a subset of C?"

"I haven't shown you the whole language yet. Two other rules guide the design of PL/68K. These rules, together with the principle of semantic equivalence, determine the form and content of PL/68K. The two rules are 'the code selection rule'—the assembler for PL/68K does no code selection, and all arithmetic operations in PL/68K correspond to unique 68000 machine instructions; and 'the register allocation rule'—the assembler for PL/68K does no register allocation, and all operations in an assignment statement are performed in the location specified by the left side of that assignment statement.

"In short, these rules say that an assembler for PL/68K never has to make any significant decisions. Because the PL/68K assembler knows how to select code and allocate registers, it will never need any of the fancy techniques used by optimizing compilers, but it will be able to produce code that is just as good as assembly language.

"I like to think of the assembler for PL/68K as a simple compiler, consisting of a parser and straightforward code generator and possibly a peephole optimizer. The whole job should take about a year to complete rather than the 10 to 20 programmer years for a typical optimizing compiler. Using compiler technology to write an assembler was the initial idea that started me thinking about PL/68K."

Now that I've presented the general ideas behind PL/68K, I'll drop this dialogue format and these fictional characters and look at the details of the language.

#### **Specifying Registers**

Because PL/68K is both assembly language and C, some way must be found to deal with assembly-language constructs such as registers, address modes, and individual machine instructions, while at the same time remaining compatible with C. These assembly-language constructs are represented by reserved words, shown in Table 1, page 29.

As for the registers of the 68000, d0 through d7 stand for the data registers, a0 through a7 for the address registers, pc for the program counter, ssr for the status regis-

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#### PL/68K

(Continued from page 27)

ter, and ccr for the condition code register, which is the lower byte of the ssr.

Standard aliases are also defined. Register a7 can also be called sp, ssp, or usp to denote the stack pointer (or system stack pointer or user stack pointer). The reserved words r0 through r7 are synonyms for d0 through d7, and the names r8 through r15 are synonyms for registers a0 through a7.

All registers on the 68000 are 32 bits long (except the status registers), but not every instruction uses all 32 bits of a register. Besides long (32-bit) operations, byte-length (8-bit) and word-length (16-bit) operations are permitted on data registers, and word-length operations are permitted on address registers. To represent the length of an operation, the name of any data register can be followed by a b to denote byte length or a w to denote word length. Thus,  $d\theta$  stands for the long register d0,  $d\theta w$  stands for the word-length register d0, and  $d\theta b$  stands for the byte-length register d0. Address registers are treated in a similar manner, except that byte-length operations are not permitted.

#### Address Modes

The 68000 has 12 different address modes, or means of accessing operands. (See Table 2, page 29.) The address modes are represented in PL/68K by five of C's operators, namely &, \*, ++, --, and ->.

Let's look, for example, at the Address Register Indirect with Postincrement address mode. (It's a lot easier to use than to say.) This mode uses the contents of an address register as the address of an operand. After the operation is performed, the address register is incremented by 1, 2, or 4, depending on the size of the operation. In traditional assembly language, that mode applied to address register a0 would be written as (a0)+. In PL/68K, that address mode is represented by \*a0++. For example, you would write d0b=\*a0++; in PL/68K instead of move.b (a0)+,d0b. Constructions such as \*++a0 are not allowed because of the code selection rule. The 68000 has no addressing mode of the form +(a0), so \*++a0 is not part of PL/68K.

The word *primitive* denotes what is called an effective address in machine-language terms. A primitive describes an operand, which may be in a register, on the run-time stack, or in static memory. In PL/68K, the valid forms of primitives are determined by the address modes I've just discussed.

#### **Declarations**

While I am talking about operands, I'll say a few words about how those operands are declared. Declarations in PL/68K are just the same as in C, except that functions do not have types. If you think about it for a moment, this means that PL/68K declarations have no parentheses. Declarations can never become unreadable as they can in full C.

In effect, declarations produce DC (define constant) and DS (define storage) pseudo-operations. (See Table 3, page 29.) Although declarations produce no executable code,

they determine what code gets produced by arithmetic operators. For instance, if a is an integer, the assignment statement a \*= b, which multiplies a by b, generates a MULS (signed multiply) machine instruction, but if a is an unsigned integer or pointer, the assignment statement generates a MULU (unsigned multiply) instruction.

As another example, the assignment  $a \neq b$ , which adds b to a, generates an ADD.B (byte length add) instruction if a is a char, but it generates an ADD.W (word length add) instruction if a is an int and generates an ADD.L (word length add) instruction if a is a long word or a pointer.

#### Assembly-Language Instructions and Pseudo-operations

Reserved identifiers also stand for 68000 machine-language instructions and pseudo-operations. A library of pseudofunctions must be linked with a PL/68K program when it is translated with a C compiler. This library, called the ops library, contains declarations and functions that allow C programs to simulate the effect of 68000 machine instructions and pseudo-operations. (See Table 4, page 32.)

The pseudofunction btst(), for example, simulates the BTST (bit test) machine instruction. In PL/68K, you would write btst(1,d0b); in those places where you would write btst.b #1,d0 in traditional 68000 assembly language.

Other pseudofunctions allow PL/68K programs to refer to assembly-language pseudo-operations. The PL/68K assembler translates the  $org(\ ), even(\ ), bss(\ ), text(\ ),$  and  $data(\ )$  pseudofunctions to the ORG, EVEN, BSS, TEXT, and DATA pseudo-operations. Similarly, the PL/68K assembler translates the  $dcb(\ ), dcw(\ ), dcl(\ ), dsb(\ ), dsw(\ ),$  and  $dsl(\ )$  pseudofunctions to the DC.B, DC.W, DC.L, DS.B, DS.W, and DS.L pseudo-operations. None of these pseudofunctions has any effect when a C compiler translates a PL/68K program. In other words, the corresponding pseudofunctions in the ops library do nothing.

You may be wondering why I keep calling these routines pseudofunctions. After all, they are perfectly good functions when compiling a PL/68K program with a C compiler. When you turn the program through the PL/68K assembler, though, it translates pseudofunctions directly into 68000 machine instructions or pseudo-operations.

#### **Expressions and Assignment Statements**

I've covered the components of low-level assembly language—registers, address modes and effective addresses, machine instructions, and pseudo-ops. Now let's see how you put these components together to make expressions and assignment statements.

Expressions are quite restricted; they are just primitives or parenthesized constant expressions. Assignments are restricted to the forms *primitive aop expression* or *primitive aop* ( *assignment* ), where *aop* stands for one of the assignment operators of the C language, namely, =, +=, -=, \*=, /=, &=, !=,  $^-=$ , >>=, and <<=. The regular arithmetic operators in C, namely, +, -, \*, /, &, |,  $^+$ , >>, and << are allowed only in constant expressions, which must be parenthesized.

You are probably wondering why all these restrictions exist. There's a short answer and a long answer. The

```
Data registers
d0, d0b, d0w, ..., d7, d7b, d7w
r0, r0b, r0w, ..., r7, r7b, r7w

Address registers
a0, a0w, ..., a7, a7w
r8, r8w, ..., r15, r15w
sp, spw, usp, uspw, ssp, sspw

Status registers
ssr, ccr

Program counter
pc
```

**Table 1:** Reserved words corresponding to 68000 registers

PL/68K	Traditional assembly language	
123	#123	
abc	abc	
&abc	#abc	
*(&abc+1)	abc+#1	
abc.25	abc+#25	
*(0x80)	\$80	
a0	a0.l	
a0w	a0.w	
d0	d0.l	
d0w	d0.w	
d0b	d0.b	
*a0	(a0)	
*a0++	(a0)+	
*a0	-(a0)	
a0 → 5	#5(a0)	
$a0 \rightarrow d0$	#0(a0, d0.l)	
$a0 \rightarrow (d0w)$	#0(a0, d0.w)	
$a0 \rightarrow (d0+5)$	#5(a0, d0.l)	
$a0 \rightarrow (d0w + 5)$	#5(a0, d0.w)	
pc → 5	#5(pc)	
$pc \rightarrow d0$	#0(pc, d0.l)	
$pc \rightarrow (d0w)$	#0(pc, d0.w)	
$pc \rightarrow (d0+5)$	#5(pc, d0.l)	
$pc \rightarrow (d0w + 5)$	#0(pc, d0.w)	

Table 2: Representing the address modes of the 68000

```
Code Generated
Declaration
                       abc: ds.b
                                    1:
char abc:
char c1 = c;
                       c1: dc.b
                                    'c';
                       cp: ds.l
char * cp;
                                    1:
                        xyz: ds.l
                                     1;
long xyz;
char a[] = "abc";
                        a:
                            dc.b
                                    'abc',0;
                        a2: ds.w
                                    25:
int
      a2 [25];
struct s1 {
          long sl;
          char *s2;
          int s3;
                        no code generated.
}:
                        s2: ds.b
struct s1 s2 [25];
union u1 {
           int ui:
           char uc;
           long ul;
                        u1: ds.b
                                    4;
};
```

Table 3: Code generated for declarations

(Continued from page 29)

short answer is that it's the only way to reconcile all of the design rules behind PL/68K with all of C's rules concerning operator precedence and arithmetic conversions. Take my word for it: You cannot have full C expressions in PL/68K and still retain the principle of semantic equivalence. For the long answer to this question, see the Appendix section.

There are additional restrictions concerning assignments to pointers. Suppose, for example, that a0 has been declared to be a pointer to some object X whose size is larger than 1. In this case, the only assignments allowed to a0 are assignments of the form

a0 + = constant:

(or constant expression) or

a0 -= constant;

To conform to C's scaling rules, these constructions are equivalent to

ADDA CONSTANT \* SIZEOF(X), A0 SUBA CONSTANT \* SIZEOF(X), A0

If you do not want this scaling to take place, you can use the <code>adda()</code> or <code>suba()</code> pseudofunctions. In general, no scaling, optimization, or other kind of tampering is ever done to the arguments of pseudofunctions. This rule has no exceptions.

Two special forms of assignment statements are permitted: location + +; and location - -;, where location stands for an address register, data register, or the name of a memory variable.

A few words about why PL/68K does not have function calls and array references—let's consider function calls first. Functions don't return values in a designated location because no single spot is right in all cases. Returning the results of a function in a register would conflict with the register allocation rule. Thus, the programmer must specify how functions will return values. Note, however, that you can pass arguments to a function in the normal way, as you shall see shortly.

The array operator [] conflicts with both the register allocation and the code selection rules. There is no suitable spot in which to evaluate subscript expressions, and picking a spot at random would violate the register allocation rule. (Actually, no single location would suffice because general C expressions can require arbitrarily many temporary locations to evaluate.) In addition, the "natural way" to access arrays on the 68000 is with pointers, which is how an assembly-language programmer would probably do it. Thus, retaining the array operator would violate the spirit behind the code selection rule.

Assignment statements are simple, but that shouldn't hurt much—assignment statements usually have only one or two operators anyway.

As the code selection rule requires, PL/68K defines what code arithmetic expressions will generate. (See Ta-

ble 5, page 32.) The += operator generates the ADD instruction, the -= operator generates the SUB instruction, and so on. Many machine instructions on the 68000 have several variants—for example, ADDA adds address registers, ADDI adds literal data, and plain ADD adds to data registers and memory locations. The assembler has to do some code selection, but the choice is easy; even traditional assemblers do that much.

#### **Boolean Expressions**

Boolean expressions in PL/68K are similar to Boolean expressions in C. All the Boolean operators !, !, and && and all the relational operators ==, !=, <, <=, >, and >= are allowed. Of course, arithmetic expressions are restricted in Boolean expressions just as they are in assignment statements.

Boolean expressions in PL/68K are not general expressions, as they are in C. Boolean expressions can only appear in the appropriate part of *if*, *do*, *while*, and *for* statements. The code fragment a = b; is not valid in PL/68K outside a structure statement. This restriction eliminates a whole class of hard-to-find errors (the programmer almost certainly meant to say  $a = b_i$ ).

Because Boolean expressions appear in limited contexts, much less work is required to evaluate them. (See Tables 6 and 7, page 32.) Boolean expressions of the form primitive or primitive relop 0 or 0 relop primitive (where relop denotes one of the relational operators) generate the TST (test operand) instruction followed by some form of the  $B_{XX}$  (conditional branch) instruction. The  $B_{XX}$  instruction chosen depends on the relop and the declared type of the operand being tested.

Expressions of the form *primitive relop nonzero constant* or *nonzero constant relop primitive* generate the CMP (compare) instruction followed by a *Bxx* instruction.

The *NOT* operator generates no code at all but instead simply reverses the "polarity" of one or more *Bxx* instructions. For example, the statement

if (a = 0)

generates

TST A BNE

while the statement

if (!(a = = 0))

generates

TST A BEQ

Similarly, the H and && operators generate no extra code. Incidentally, you can use parentheses in Boolean expressions to affect the order of binding of Boolean operators. For instance,

if (a && !(b < 5!|b > 20))

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#### Pseudofunctions corresponding to 68000 machine instructions abcd, add, adda, addi, addg, addx, and, andi, asl, asr bcc, bcs, bchg, beq, bge, bgt, bhi, ble, bls, blt, bmi, bne, bpl, bvc, bvs bchg, bclr, bra, bset, bsr, btst chk, clr, cmp, cmpa, cmpi, cmpm, dbcc, dbcs, dbeq, dbf, dbge, dbgt, dbhi, dble, dbls, dblt, dbmi, dbne, dbpl, dbt, dbvc, dbvs, divs, divu eor, eori, exg, ext, jmp, jsr, lea, link, lsl, lsr move, movea, movem, movep, moveq, muls, mulu nbcd, neg, negx, nop, not, or, ori, pea reset, rol, ror, roxl, roxr, rte, rtr, rts scc, scs, seq, sf, sge, sgt, shi, sle, sls, slt, smi, sne, spi, st, svc, sbcd, stop, sub, suba, subi, subq, subx, swap tas, trap, trapv, tst, unlk Pseudofunctions corresponding to assembly-language pseudo operations dcb, dcw, dcl, dsb, dsw, dsl org, data, text, bss, even

Table 4: Pseudofunctions

Operator	Gene	rated c	ode
a += b	add	b, a	(or adda or addi or addq)
a-=b	sub	b, a	(or suba or subi or subq)
a *= b	muls	b, a	(or mulu)
a /= b	divs	b, a	(or divu)
a %= b	divs	b, a	(or divu)
	swap	а	
a>>=b	asr	b, a	(or ror)
a <<= b	asl	b, a	(or rol)
a &= b	and	b, a	(or andi)
al= b	or	b, a	(or ori)
a ^= b	eor	b, a	(or eori)
a + +	addq	#1, a	(or addi)
a	subq	#1, a	(or subi)

Table 5: Code generated by arithmetic operators

Boolean	Gener	ated code	е
if (Z)	bnz	false	
if (a)	tst bz	a false	(or cmpa
if (a < b)	cmp bge	b,a false	
if (!a)	tst bnz	a false	
if (a && b)	tst beq tst beq	a false b false	
if (!(a && b))	tst beq tst bne true:	a true b false	

Table 6: Code generated by Boolean expressions

Sign	ed Con	parisons	<b>Unsigned Comparisons</b>		
		if $(c1 = = c2)$			
cmp	c2,c1		cmp c	2,c1	
bne	false		bne fa	alse	
		if (c1!= c2)			
cmp	c2,c1		cmp c	2,c1	
beq	false		beq fa	alse	
		if $(c1 < c2)$			
cmp	c2,c1		cmp c	2,c1	
bge	false		bcc fa	alse	
		if $(c1 < = c2)$			
cmp	c2,c1		cmp c	2,c1	
bgt	false		bhi fá	alse	
		if $(c1 > c2)$			
cmp	c2,c1		cmp c	2,c1	
ble	false		bls fa	alse	
		if $(c1 > = c2)$			
cmp	c2,c1		cmp c	2,c1	
blt	false		blo fa	alse	

Table 7: Code generated by Boolean comparisons

Масі	o N	lame	Pseudo- function	Meaning
Z	or	EQ	CC_Z	zero
	or	NE	cc_nz	not zero
С	or	CS	cc_c	carry
NC	or	CC	cc_nc	no carry
V	or	VS	CC_V	overflow
NV	or	VC	cc_nv	no overflow
GT			cc_gt	greater than
GE			cc_ge	greater or equal
LS			cc_ls	less than
LE			cc_le	less than or equal
HI			cc_hi	high .
LO			cc_lo	low
MI			cc_mi	minus
PL			cc_pl	plus

Table 8: Condition code constants

S	yntax
	(1) if ( boolean ) { statement list } (2) if ( boolean ) { statement list 1 } else { statement list 2 }
С	ode generated for (1)
	\$ Evaluate boolean. If false, jump to label 1 \$ statement list \$ label 1:
С	ode generated for (2)
	\$ Evaluate boolean. If false, jump to label 1 \$ \$ statement list 1 \$ bra label2; label1: \$ statement list 2 \$
	label2:

Table 9: The if statement

### PL/68K (Continued from page 30)

is valid and generates

TST A
BEQ
CMPI 5, B
BLT
CMPI 20, B
BGT

You can specify condition code values directly. (See Table 8, page 32.) For instance, the statement if(Z) tests the current value of the zero bit in the condition code register and generates the BNZ (branch not zero) instruction. Z is a macro in C, defined in the ops library, which expands to a call to the pseudofunction  $cc\_z($ ).

#### Structure Statements

I said earlier that PL/68K has all C's structure statements—if, do, while, for, and switch. They look exactly like C code, but curly braces are required surrounding statement lists in structure statements. In other words, structure statements have the form

$$\label{eq:continuity} \begin{split} & if(.\ .\ .) \left\{ statement\ list \right\} \\ & if(.\ .\ .) \left\{ statement\ list \right\} \ else \left\{ statement\ list \right\} \\ & while(.\ .\ .) \left\{ statement\ list \right\} \ while (.\ .\ .); \\ & for (.\ .\ .) \left\{ statement\ list \right\} \\ & switch(.\ .\ .) \left\{ statement\ list \right\} \end{split}$$

In my opinion, allowing curly braces to be optional is a big flaw in C. In this example:

 $if (abc < 5) \\ xyz = 5; \\ if (abc < 6) \\ xyz = 6;$ 

the indentation is misleading and will probably cause a bug. This kind of error can be extremely difficult to find.

Let's see what code PL/68K's structure statements produce. In the accompanying tables, the dollar sign denotes code that corresponds to some language construct. For example, \$ statement list \$ stands for whatever code is generated for the statement list. The statement list could be arbitrarily complicated—for instance, it could contain nested structure statements. The notation

\$ Evaluate boolean. If false, jump to label1 \$

indicates that code is generated for the Boolean expression such that a jump to *label1* is taken if the Boolean expression is false. Otherwise, control falls through to the following code. Labels are indicated in the usual way, by identifiers followed by colons. All generated labels are, of course, unique, even though they may have identical names in the tables.

Table 9, page 32, shows the *if* statement. When an *if* statement contains no *else* clause, the Boolean expression

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```
Syntax
```

(1) while ( boolean ) { statement list } (2) while (1) { statement list }

#### Code generated for (1)

bra continue\_label; label1: \$ statement list \$ continue label: \$ Evaluate boolean. If true, jump to label 1\$ break\_label:

#### Code generated for (2)

continue\_label: \$ statement list \$ bra continue\_label; break\_label:

#### Table 10: The while statement

#### Syntax

(1) do { statement list } while ( boolean ); (2) do { statement list } while(1);

#### Code generated for (1)

label1: \$ statement list \$ continue\_label:

\$ Evaluate boolean. If true, jump to label 1 \$ break\_label:

#### Code generated for (2)

continue\_label: \$ statement list \$ bra continue\_label; break\_label:

#### Table 11: The do statement

#### Syntax

(1) for (assignment list 1; boolean; assignment list 2) statement list }

(2) for ( assignment list 1;; assignment list 2) { statement list }

(3) for ( assignment list 1; 1; assignment list 2) { statement list }

#### Code generated for (1)

\$ assignment list 1\$ bra label0; label1: \$ statement list \$

continue\_label:

\$ assignment list 2\$

\$ Evaluate boolean. If true, jump to label 1 \$ break\_label:

#### Code generated for (2) or (3)

\$ assignment list 1\$ label1: \$ statement list \$ continue\_label: \$ assignment list 2\$

bra label 1; break\_label:

Table 12: The for statement

#### PL/68K (Continued from page 33)

is evaluated, and control either falls through to the then clause or a jump is made to the end of the statement.

Similar code is generated when the if statement contains an else clause. The Boolean expression is evaluated, and control either falls through to the then clause or a jump is made to the else clause. A BRA (branch always) instruction following the then clause skips around the else clause.

The code generated for the while statement, shown in Table 10, page 34, might be a little controversial. The first instruction is a branch to the end of the loop so that the loop test occurs at the bottom. This produces the fastest code unless the while loop is executed less than once on average. In the rare cases in which this jump is unwanted, the programmer must simulate the loop in some way.

Notice the labels called *continue\_label* and *break\_label*. These are used as target labels for the break and continue statements. In other words, within a while, do, or for statement, the effect of a continue instruction is to generate a branch to the continue\_label defined for that statement. Similarly, a break statement generates a jump to the appropriate break\_label. As in C, you can also use the break statement inside a switch statement.

The while statement generates different code if the Boolean expression is a nonzero constant. This is a common idiom in C, and the definition of PL/68K ensures that there is no time penalty for using it.

The code for the do statement, shown in Table 11, page 34, is similar to the code produced by the while statement. The code for the for statement (see Table 12, page 34) is more interesting. If the loop test in a for statement is nontrivial, the code for it appears at the bottom of the loop. Note also that the syntax of the for statement is more restricted than in C.

The switch statement, shown in Table 13, page 37, generates a jump table—i.e., a table of addresses. Code is generated that jumps through that table to the proper case statement, based on the contents of a register. Note that the switch statement destroys this register.

It is sometimes better to generate a sequence of tests rather than a table jump, but the case statement always generates a table jump. Remember, each language construct in PL/68K stands for a particular sequence of code—if you want a sequence of tests, use a sequence of if statements; if you want a table jump, use a switch statement.

Many compilers generate jumps to jumps when they translate structure statements, but the definition of PL/ 68K requires that all jumps to jumps (and jumps to return statements) be eliminated. The assembler can do this in several ways. For instance, if the assembler creates a parse tree for an entire function before any code is generated, it's easy for the code generator to look at the target of any jump to see if it is another jump or a return instruction. Alternatively, the assembler can use a standard peephole optimizer.

#### **Function Calls**

Functions in PL/68K can have formal parameters and lc

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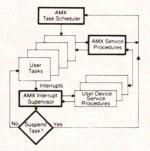
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PL/68K

(Continued from page 34)

cal arguments, just as in C. The code shown in Table 14, page 37, is generated by function calls. Code is generated to push all arguments on the stack, a JSR (jump to subroutine) instruction is generated, and, if necessary, an ADD instruction is generated to pop arguments off the stack. One long word is always reserved on the stack for the first argument, which shortens the calling sequence when there are less than two arguments.

The ADD instruction can be eliminated by having the called program, instead of the calling program, pop the arguments off the stack, but the sequence shown in Table 14 is the fastest. If you eliminated the ADD instruction and pushed the arguments in the same way (that is, above the return address), the called program would need to do much more work to pop off its arguments. You could also push the actual arguments below the return address, but that way actually increases the length of the calling sequence. In order to push arguments below the return address, you would have to use an instruction such as move arg,n(sp), which is 2 bytes longer than move  $arg_{\bullet} - (sp)$ .

Unlike standard C, PL/68K does specify the order in which arguments are pushed, namely in reverse order. Thus, a function that takes a variable number of arguments—printf() for instance—will find its first argument on the top of the stack.

Of course, it is often best to pass arguments in registers, but PL/68K doesn't need a separate mechanism to do this. Suppose you have a function called g() whose two arguments are passed on the stack. To change g() so that it will take its arguments in registers, you just define the following macro

#define g(a,b)d0=a; d1=b; g1()

and change g's name to g1. Notice that a statement such as

if (...) g(x,y);

cannot cause problems in PL/68K because you must write

if  $(...) \{g(x,y);\}$ 

instead. (If braces were omitted, after macro expansion, the code would be

if (. . .) d0 = a; d1 = b; g1();

and only the assignment d0=a would be part of the if statement.)

This macro might generate redundant code. Suppose it were called with d0 as the first argument, for instance. The macro would expand to d0=d0 and generate the instruction move d0,d0. To handle that problem, the PL/ 68K assembler eliminates redundant moves. If you must generate such a redundant move for some reason, use a pseudofunction. Pseudofunctions are never secondguessed by the assembler.

```
Syntax
   switch (req) {
   case constant 1: statement list 1:
   case constant 2: statement list 2:
   case constant n: statement list n;
                       default statement list;
Code generated
   if (reg < min || reg > max) {
          goto default_label;
   else (
      $ goto the routine whose address is at table [reg] $
   table:
   dc.I label imin;
   dc.l label imin+1;
   dc.I label imax-min+1;
   label1
   $ statement list 1$
   label 2:
   $ statement list 2$
   label n:
   $ statement list n $
   default_label:
   $ default statement list $
   break label:
```

### Table 13: The switch statement

```
No arguments
           function
   jsr
One argument
   move arg, (sp)
   isr
           function
Two or more arguments
   move argn, (sp)
   move arg<sub>n-1</sub>, -(sp)
   move arg2, -(sp)
           arg<sub>1</sub>, -(sp)
   move
            function
   isr
            # size of all arguments except arg,, sp
   add
```

Table 14: Code generated for function calls

```
Entry

movem.l all_local_registers, -(sp)
link a<sub>n</sub>, # - size of local auto variables - 4

Exit

unlk a<sub>n</sub>

movem.l (sp)+, all_local_registers
rts
```

**Table 15:** Code generated for entry/exit of  $A_n$ -based functions

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### Accessing Variables Within Functions

Now that I've covered the passing of arguments to a function, I'll explain how the function gets hold of those arguments. This is a complicated subject, so before getting involved in some messy details, let's handle the easy cases.

First, PL/68K keeps its hands off all registers declared outside any function. These global registers can be accessed from within functions, but PL/68K never generates code to save or restore them. Consequently, you can prevent PL/68K from interfering with any register simply by declaring that register outside a function. By the way, the register keyword is not valid outside functions, but that does not prevent you from declaring register variables anywhere you wish, either in C or in PL/68K. For instance, you can declare a0 to be global simply by saying char \* a0; outside any function.

Second, except for these global registers, all registers used in a function are saved on entry and restored on exit from the function. The assembler uses the MOVEM.L (move multiple register, long) instruction for this purpose. Accessing register variables declared in a function is, of course, easy.

Third, if a local variable is declared to be static, memory is allocated to that variable in static memory, not on the stack. No extra code is needed to access that variable, either on function entry or exit.

Static internal variables are not very useful; because the values of static internals are retained between invocations of a function, static internals are destroyed by recursive function calls. Also, on many machines (including the 68000) accessing a variable in static memory is more expensive than accessing a "local auto" variable—that is, a local stack variable. At any rate, generating code for static internal variables is straightforward and is not affected by the following complications.

Functions need to access two kinds of nonregister variables: formal parameters and local auto variables. Both types are allocated on the stack. You have three choices

```
If one or more functions called within this function

Entry

movem.l all_local_registers, -(sp)
suba # size of local auto variables + 4, sp

Exit
adda # size of local auto variables + 4, sp
movem.l (sp)+,all_local_registers
rts

If no function called within this function

Entry
movem.l all_local_registers, -(sp)

Exit
movem.l (sp)+, all_local_registers
rts
```

**Table 16:** Code generated for entry/exit of sp-based functions

for how PL/68K generates code for these stack variables. You make your choice using one of three pseudofunctions:  $base(A_n)$ , base(sp), or nobase(). If none of these appears in a function, the default is base(sp).

First, you can have PL/68K access stack variables via an address register that is different from the stack pointer, as shown in Table 15, page 37. If the  $base(A_R)$  pseudofunction appears anywhere in the function (where  $A_R$  is any address register except the stack pointer, a7), PL/68K generates a LINK (link and allocate) instruction on function entry and an UNLK instruction on function exit. All stack variables are accessed via the address register named in the  $base(A_R)$  pseudofunction.

Second, you can have PL/68K access stack variables via the stack pointer, as shown in Table 16, page 38. The assembler generates this kind of code if the base(a7) or base(sp) pseudofunction appears anywhere in the function. The code generated for sp-based functions is more complicated, but more efficient, than that for  $A_R$ -based functions. If an sp-based function contains no function calls, the stack pointer is not incremented on function entry and local auto variables are accessed using positive offsets from the stack pointer. If the function does contain other function calls, however, space is reserved for local auto variables by incrementing the stack pointer on function entry, and local auto variables are accessed using negative offsets from the stack pointer.

Third, you can access stack variables by hand. If the <code>no\_base()</code> pseudofunction occurs anywhere in the function, no code is generated on function entry or exit, declarations of stack variables are ignored, and no explicit references to stack variables are permitted. The <code>no\_base()</code> pseudofunction is useful when you must play some kind of game with the stack.

*Sp*-based functions are somewhat dangerous. If the stack pointer is changed using a pseudofunction, the offsets used to access stack variables are going to get out of sync. To handle this problem, several pseudofunctions, shown in Table 17, page 38, allow you to indicate that the offsets should be changed.

The push(arg) and pop(arg) pseudofunctions are equivalent to the move(arg, \*--sp) and move(\*sp++,arg) pseudofunctions, but in addition, they tell the PL/68K assembler to adjust the offsets used to access stack variables in sp-based functions. The addsp(n) and subsp(n) pseudofunctions are equivalent to the addi(n,sp) and subi(n,sp) pseudofunctions, but again they tell the assembler to adjust offsets. Finally, the adjsp(n) pseudofunction generates no code but tells the assembler to adjust the offsets.

Pseudofunction	Meaning	
base(a <sub>n</sub> ) base(sp)	use a <sub>n</sub> basing for suse sp basing for s	
nobase()		ariables "by hand"
push(arg) pop(arg)	move arg, $-(sp)$ move $(sp)+$ , arg	and adjust stack offsets and adjust stack offsets
addsp(n)	adda #n, sp	and adjust stack offsets
subsp(n)	suba #n, sp	and adjust stack offsets
adjsp(n)	adjust stack offset	ts

Table 17: Pseudofunctions for stack operations

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### PL/68K

(Continued from page 38)

These pseudofunctions have no effect on  $A_n$ -based functions.

### Summary

To summarize the strengths and weaknesses of PL/68K, the syntactical restrictions that make PL/68K a strict subset of C are listed as follows:

- All operand/operator combinations must correspond to an instruction in the 68000 instruction set.
- Functions neither have types nor return values.
- The array operator ( ) is eliminated, but arrays may be declared.
- · Assignment statements are severely restricted.
- Statement lists must be surrounded by curly brackets.
- There are no floating points, or double constants, or operators.

There are no additions or changes to C syntax. PL/68K and assembly language are semantically identical; the advantages of PL/68K over assembly language are all syntactical:

- · C syntax makes programs easier to read.
- C syntax eliminates many common coding errors.
- Far fewer visible labels are required.

PL/68K is not the successor to C nor is it superior to C for most programming projects. The advantages of PL/68K over C apply only in limited circumstances, but when performance is paramount, PL/68K stands out:

- All machine resources and instructions are available.
- Global allocation of registers is possible.
- Total control over generated code is possible.
- · Generated code is smaller and faster.

Figure 2, page 42, shows the relationship between PL/68K and its two parent languages. PL/68K is only a subset

of C; not all of C is included. On the other hand, C must be augmented by the ops library in order to simulate all the machine resources of assembly language.

Listing One, page 101, shows an example of a PL/68K function. This function finds a name in a symbol table and returns information about that name in registers. Listing Two, page 101, shows the code generated by the PL/68K assembler for that function.

### Conclusion

PL/68K started life as a C-like assembly language for the 68000 chip. High-level assemblers are not new—a paper by Niklaus Wirth (*Journal ACM* 15, January 1, 1968) described a high-level assembly language for IBM 360 machines.

My early thinking about PL/68K was influenced by Wirth's paper and by the register allocation and code selection principles. At that stage, I was interested in using compiler technology to create better assemblers. I saw no particular reason to make PL/68K a subset of C.

I felt dissatisfied with the initial version of PL/68K; it was doomed to be "just another computer language." Not wanting to add another minor dialect to the Babel of computer languages, I decided to make PL/68K's syntax identical to C's. This was a lucky decision.

From the first, PL/68K had to be able to name all machine resources, including machine instructions and assembly-language pseudo-ops. Early versions of the language allowed "raw" lines of assembly code to be interspersed with C-like lines of code. Although this approach had some merit, my decision to make PL/68K's syntax compatible with C convinced me to use functional notation to represent assembly-language features. This change was also fortunate.

I began to speculate about what would happen if a PL/68K program were compiled with a C compiler. I asked myself, suppose the program produced by the C compiler and the program produced by the PL/68K assembler were semantically equivalent? Suddenly, PL/68K became not just another assembly language but a new way of using C.

The principle of semantic equivalence guided a rede-

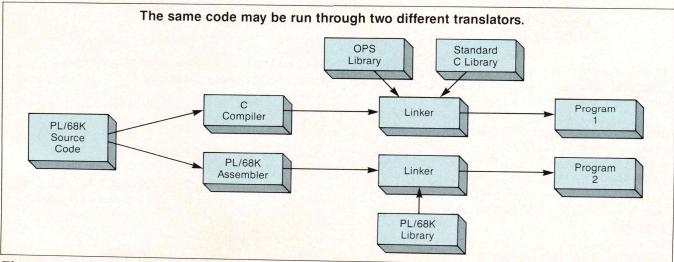
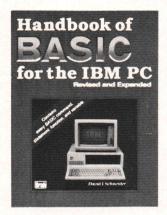


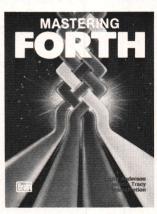
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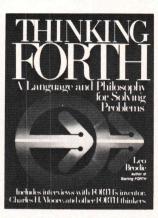
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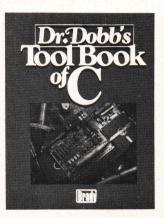
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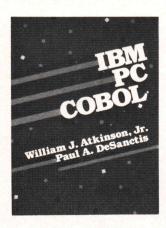
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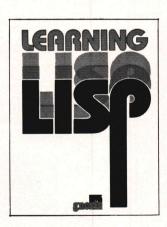
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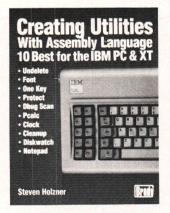
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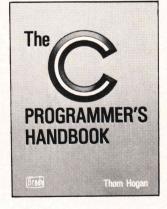
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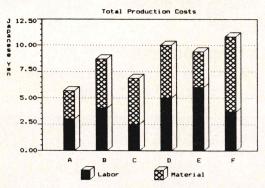
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PL/68K (Continued from page 40)

sign of the language; any construct that violated this principle had to go. In addition, I invented the ops library so that C programs could simulate 68000 machine instructions. Along with the ops library, the concept of pseudofunctions was born, and the language was complete.

PL/68K allows me to sidestep a question that has been haunting me ever since I started programming—whether to program in assembly language and accept the resulting inconveniences or to program in a high-level language and accept a final product that is larger and slower than it could be. PL/68K solves that dilemma.

You can design and write a program in C, keeping in mind the possibility that you will convert it to PL/68K eventually. You can write difficult parts of the program, or parts of the program that will not appear in the final version, using all of C's features. After you have debugged the C program, you can produce a PL/68K version of the program, if desired, by rewriting or excluding those parts of the program that use full C. You then test the PL/68K version and improve its performance as much as you require.

PL/68K hits precisely the right level of abstraction for systems programming. All the features of C allow you to design and code programs easily, but when you need to do low-level work, nothing stands in your way.

A final thought—you could transfer most of the syntax and all the design rules of PL/68K to a similar language for other machines. In this limited sense, PL/68K is a machine-independent assembly language—PL/68K programs are much more portable than programs written in traditional assembly language.

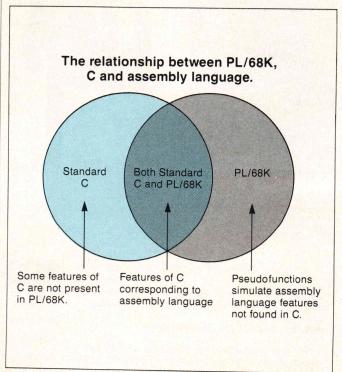


Figure 2

### Note

A compiler for the PL/68K language is available from the author.

# Appendix: Why Expressions Are Restricted in PL/68K

The outline of the argument is as follows. The register allocation rule conflicts with C's precedence rules (which govern the order in which operators are applied) and C's rules for converting operands from one type to another (known as the "usual arithmetic conversions" in the *C Reference Manual*). Thus, by the principle of semantic equivalence, all constructions in PL/68K that involve C's precedence and type conversion rules must be eliminated.

Consider a typical C expression, such as a=(b \* c) + (d \* e);. The register allocation rule says that this expression must be evaluated in location a, but that is not possible. At least two separate locations are required—one to hold the subexpression (b \* c) and another to hold (d \* e). In general, any expression involving parentheses may require more than one location to evaluate.

The only way to make a single location suffice for the evaluation of all expressions is to require left-to-right (or right-to-left) evaluation of operators. Alas, this changes the precedence of operators, except in cases such as a = ld \* c + c;, which is evaluated left to right in C anyway.

Even such limited expressions cannot be allowed, though. C's rules for type conversions state that conversions are made during the evaluation of the right-hand side of the expression. Only after the right side is completely evaluated does the assignment take place. PL/68K has only one place to make those conversions, however—the left-hand side of the expression. The operands themselves cannot be converted because that would change their value. Thus, only one operand can be converted and only one operand can be allowed on the right side of an assignment statement.

DDJ

(Listings begin on page 101)

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The instruction set

for the 68000 family

is nearly orthogonal.

This is important for

a system on which a

lot of assembly-lan-

guage development

work is to be done.

or the past year, we at Terra Nova Communications have been involved in a development project that requires a simple, fast, clean 32-bit microprocessor operating system. After a great deal of research, we were unable to find a commercial system that met our stringent requirements of extremely fast response time (even under a load of 20 users), low price (less than \$10,000 for both system and software), compact code size (we wanted a system kernel, including all the utility routines discussed in this article, that required less than 20K of object code), and simple programming of applications. After some brainstorming, we created a 68000 multitasking kernel that met and even exceeded our expectations of speed and compactness. Released from hardware requirements by our decision to write the kernel ourselves, we decided to use the VMEbus hardware configuration because of its standardization, complete hardware specification, relatively low price, ease of expansion, and the availability of lots of highspeed hardware devices. We were also impressed by the reliability and ease of use of the Eurocard connectors used with the VMEbus.

Why Not Use an Existing OS? Our requirements for speed and

compactness stemmed primarily from the need to handle a large number of I/O tasks over serial lines with-

out incurring large overheads for interrupt handling, disk access, and context switching-that is, we needed to be able to do significant amounts of I/O without slowing the system. Our kernel should eventually be able to handle 20 real-time users over serial lines at 1,200 bps, including full-speed block transfers, with no perceptible response-time delay and only minor slowing of the byte-transfer rates. We needed a system that would degrade gracefully; rather than pausing in midstream as one task takes over the system for an appreciable fraction of a second or as tasks are paged in and out, it should slow down gradually as the load increases, always providing steady and uniform output, even if it's at a reduced byte transfer rate. None of the commercial systems we examined

without significant degradation. Examination of the code used in several of the commercial kernels we sampled showed some interest-

had this property, nor were they

able to handle the load required

ing reasons for this: Most kernels contained code designed to handle all sorts of unlikely circumstances that might arise in an environment where you don't know what sorts of programs might be running. Not only did this code add significantly to the size of the kernel but it also slowed down the process of context switching between tasks in many cases. We needed the fastest possible context switch in order to guarantee that minimum system time was spent on this. Fortunately, we knew exactly which applications would be running on our system, and we were able to design a complete application/system interface to make application coding easy. Because we knew that all application code running under our kernel would be "polite" (would follow all the rules of the interaction between application and kernel) and that all source code would be available for debugging, we were able to dispense with a lot of the error-handling code usually present in commercial kernels.

We discovered that another contributing factor to the time required for a context switch was the magnitude of the context that was switched. In most of the systems we examined, all the machine registers were saved and restored, and full status information was saved, both of which took up a significant amount of processing time. As we'll explain later, we fixed this in a rather unorthodox way.

Further, many commercial kernels required the use of a memory manager chip and spent significant amounts of time paging users in and

Nicholas Turner, 10 McGranahan Court, Boulder Creek, CA 95006, (408)

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out to compensate for a small system memory. We opted against memory management, mainly because it solved no problems for us. We didn't need any sort of memory protection; in fact, one of the most important criteria was that all tasks must be able to quickly read and write data belonging to any other task or to the system itself. Also, our memory requirements were not large (minimum 512K, expandable to several megabytes) and because of the amount of code sharing between tasks, the structure of our data heap, and the heap's interaction with the disk system, there was no need for hardware paging of memory.

# Why Program the Kernel in Assembly?

It was clear from the start that, in order to get the kind of performance we wanted from the system, the inner kernel had to be written in native code. Even compiled C or Forth would have had to be manually "tuned" in assembly source code form. Further, we could see several difficulties with compiled language-we needed to do so many 'tricky'' things to extract the last few cycles from the system kernel that writing it in compiled code was out of the question. By putting the entire kernel in machine code, we simplified the effort required to make major changes (it's all in the same language) and made possible a far more complete and integrated tuning.

Eventually we expect to put up at least a C compiler and a Forth interpreter for faster development, but for the moment all development is in assembly for speed and compactness. Because assembly was the language of choice, selection of the target processor was the next important issue.

### Why Use the 68000?

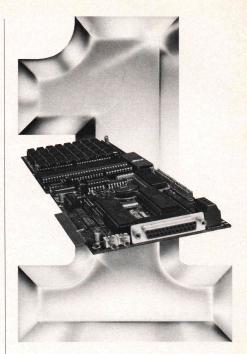
The eventual goal of our project is to provide a responsive telephone dialup system that is able to support up to 30 or 40 simultaneous calls without significant performance degradation. Such a task requires a truly powerful processor, even if the whole system is written entirely in native code. For several reasons, we have chosen the 68000 family of processors for our base hardware. The most important reason is that the instruction set is extremely versatile and powerful. It is in many ways a true 32-bit instruction set, although, unless you are using a 68020, you must put up with slightly slower memory access for 32-bit reads and writes because of the 16-bit data path.

The instruction set for the 68000 family is nearly orthogonal—that is, almost every instruction can be used with any of the 12 addressing modes. This is important for a system on which a lot of assembly-language development work is to be done because the programs become much easier to generate, read, and debug. Unfortunately, even the 68000 is not perfect. Several times we've encountered annoying restrictions—for example, we've often cursed our inability to do a PC-relative store.

The 68000 also has another advantage for assembly-language programmers: The memory architecture in its native mode, without the extras added by a memory management chip, is perfectly flat. That is, the address space is completely continuous from \$00 0000 all the way to \$FF FFFF-or \$FFFF FFFF if you have a 68020. For an assembly hacker, this is far more desirable than the segmented architecture required with the Z8000 or 80286 or with any 8- or 16bit processor. For people writing in a high-level language, this is not an issue because they never deal directly with the memory at all. But for us, it's nice to be able to chop up that big address space in any way we like. As I'll explain later in this article, we have chosen to make it into an enormous heap and virtual disk area, thus making the fullest possible use of each and every byte.

Finally, after it became clear that the VME hardware bus was the best bet for our needs, the 68000 processor family was the logical choice because of the large number of 68000-oriented products available for the VME architecture.

In this article I'll describe briefly how our operating system fits together, and then I'll get to the fun stuff: the tricks and shortcuts we used to get such incredible performance out of the 68010 in our system. If you are already familiar with how a multitasking operating system fits together, you might want to skim down to the tricks and shortcuts.



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PROGRAMMERS' PFANTASIESTM

MULTITASKING OS (Continued from page 45)

### **Memory Structure**

RAM memory is divided into two areas: the system zone and the heap, as shown in Figure 1 (page 47). The system zone begins at memory address 0 with the 68000 vector list and continues up from there. It's quite small and contains only a few data structures. Figure 2 (page 47) diagrams the structure of the system zone.

The 68000 vector list contains all the hardware vectors required for processing of interrupts and exceptions. It requires \$300 bytes. The data in the vector list is first set up by the initialization section and modified thereafter by the I/O manager as device interrupts are added or deleted.

Above the vector list is a small zone containing miscellaneous system variables and pointers. The time of day and date are kept here, along with pointers for the various linked lists maintained by the kernel, plus some other information that needs to be quickly accessible via absolute short addressing (the fastest way to get at a memory location from the 68000). Several I/O devices also have data here, where it can be accessed by interrupt routines without the overhead of following pointers through memory.

Following the system variable zone, we have a rather strange beast in today's world: an old-fashioned jump table! This table contains more than 100 absolute-long address mode JMP instructions at the momenthundreds more are planned.

The next structure in the system zone is the task control buffer (TCB) table, which is an array of data structures that are linked via pointers into several doubly linked lists. Each task is associated with one TCB in the TCB array. When a task releases control to the next task, the context switcher reads the pointer from the outgoing task's TCB to get the address of the next task's TCB. This prevents unnecessary overhead while reading TCBs belonging to inactive tasks in the array because they are not part of the active task linked list. It also makes possible an extensible TCB array: If the primary array is full when another task is about to be spawned, the task manager can allocate a nonrelo-

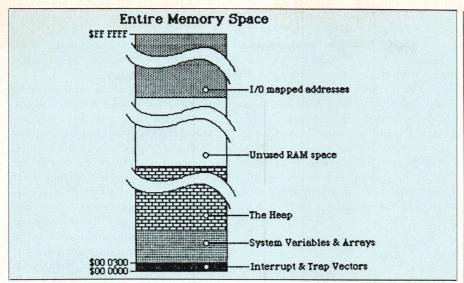


Figure 1: Basic organization of the 68000 memory space

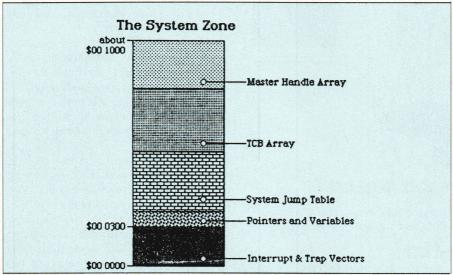


Figure 2: The system variables and arrays

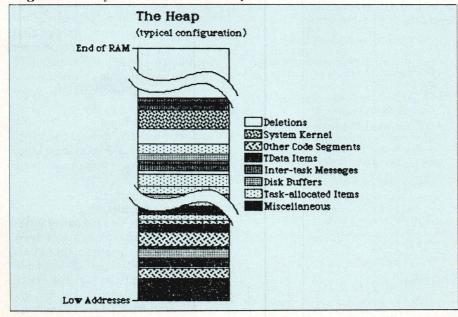


Figure 3: Typical structure of the heap



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catable item in the heap and continue the TCB array there.

Another extensible array follows the TCB table: The master handle array contains handles to relocatable heap items. (A handle is the address of a pointer.) All accesses to relocatable heap items must be dereferenced (followed through the handle to the pointer to the actual heap item) before a task can use the item. That way, if an item is relocated by the heap munger (see description of the heap munger, later) during its background heap-optimization, any tasks that own the relocated item will still be able to find it because the heap munger always fixes the master handle so it is correct after moving a heap item. This master handle is often located in the master handle array, although it doesn't need to be. Wherever it is, though, it must be in a nonrelocatable place so that the heap munger can follow a pointer back

from the heap item to its master handle.

After the last master pointer is an end mark. At the next 32-byte boundary, the heap begins.

The heap (Figure 3, page 47) takes up all the available RAM beyond the system zone. It is a single data structure composed of chunks called items. Every single byte of the heap belongs to an item of one sort or another. An item can be a deletion, or it can contain actual information. Items that contain information can be allocated (owned by one or more tasks) or unallocated. Unallocated items can be purged by the heap munger (see later) if it needs to make more room for a memory request from a task. For speed and ease of programming, every single heap item begins and ends on a 32-byte boundary.

Virtually all system information that doesn't have to be addressed absolutely is stored in the heap, in various heap items owned by the system kernel. In addition, the heap contains all executable code, including the system kernel itself, in chunks called code items.

Each heap item contains a 32-byte heap header record, followed by zero or more 32-byte data blocks. The header record contains the information necessary for the heap manager and the heap munger to identify, move, and validate each item. The heap is organized, like much of the rest of the system, as a doubly linked list. As the heap munger scans through it, it follows the pointers forward or backward to verify that all is in order. If it finds anything that is not completely kosher, it immediately stops the system, takes over the system console, and enters the debugger with a descriptive error message. When a bug occurs, it is often the heap munger that detects the problem (in the form of a messed-up heap header) before anything happens.

Division of Labor

The kernel is divided into several distinct blocks of code. They are of three types: one-time routines (initialization), system calls (routines available from every task), and discrete tasks (self-contained programs that run under the context switcher, just as do application tasks). Some of the system



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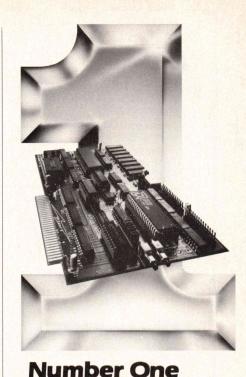
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Call name	Description
	(Task Manager Calls)
Spawn	Create a new task
Kill	Destroy a task (by task number)
Suicide	Kill the calling task
	(Heap Manager Calls)
HeapGimme	Allocate an item in the heap
HeapDel	Release (delete) a heap item
FillZero	Re-initialize a heap item
GetMaster	Assign a master handle for an item
	(Message Manager Calls)
SendMsg	Send a copy of a block of memory to another task
Del1Msg	Delete message from top of incoming queue
DelMsgs	Delete entire incoming message queue
TxtMsg	Send a message of type "TEXT"
GetMsg	Fetch next message in queue
HandleMsg	Analyze incoming message and handle if standard type
	(Character I/O Manager Calls)
DevReq	Request a character I/O channel
DevDemand	Demand a character I/O channel (usually impolite)
DevRel	Release a character I/O channel
PrToStd	Select this task's standard character I/O device
PrToMem	Select the MemPrt device (see text)
	(Text Manager Calls)
GetCommand	Input a command line and parse it, passing control to the appropriate routine based on the command.
AddCmdTab	Add a set of commands to the existing command set
DoCommand	Parse and execute a command already stored in memory
GetPSW	Input and encrypt a password (to be compared with an encrypted password from the user file)
MoveString	Move an ASCII string
GetLine	Input a line of ASCII text from the character I/O device
PrLine	Print an ASCII string on the character I/O device
Print	Print a line of text. The line is expected to immediately follow the JSR Print instruction.
CompString	Compare two ASCII strings
	(Miscellaneous System Calls)
Random	What system would be complete without random numbers?
Sqrt	Square root of 32-bit integer

Table 1: Some of Terra Nova's system calls

TRAP-oriented calls		
Instruction	Cycles used	Description
TRAP #n	38	call the system routine
MOVE.L 2(SP),A0	16	point to word argument
MOVE.W (A0)+,D0	8	fetch the argument
MOVE.L A0,2(SP)	16	update return address
(variable)	(variable)	decode argument word
		useful code
RTE	24	return to caller
	104 (+ decode)	total cycles for overhead
JSR-oriented calls		
Instruction	Cycles used	Description
JSR Label.W	18	call low memory entry poin
JMP Label.L	12	call actual routine
		useful code
RTS	16	return to caller

Table 2: Comparison of TRAP- and JSR-oriented system calls



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### MULTITASKING OS (Continued from page 48)

calls we've developed are listed in Table 1 (page 49). This is not intended to be a complete list, only to give some of the system's flavor.

The initialization section is the block of code that gains control before anything else happens. It starts with a brute-force approach: It grabs control from whatever operating system invokes it, and then it sets up the bare-bones data structures for the system. I'll go into greater detail about the initialization section later, when I talk about tricks and shortcuts.

The most important low-level code segment is the context switcher, which is the routine that receives control from one task and passes it on to the next. It is extremely small and extremely fast, and it makes a lot of assumptions about the tasks as it does its job. This is by design: By making assumptions and forcing the tasks to adhere to them, a lot of overhead is eliminated. Again, I'll go into detail about the context switcher in the section on tricks and shortcuts.

The task manager is composed of a group of system routines available to every task. They allow a task to create, destroy, and manipulate other tasks or itself.

The heap manager is a collection of routines that allow any task to request, release, lock, enlarge, or otherwise manipulate heap items.

The message manager is a collection of system routines that make possible a clean and well-defined message-passing protocol between tasks. Simply by pointing at a block of memory and calling a system routine, any task can send a copy of any piece of data to any other task. Reading queued messages from other tasks is similarly easy.

The character I/O manager is a set of system calls and interrupt service routines. Together with the text manager, it makes possible a simple I/O structure, in which each task can select any physical or logical device for I/O by passing a device number. Serial input is interrupt-driven, and serial output is polled. Device drivers can be added or removed from the I/ O manager with another set of sys-

The text manager is a collection of

system routines that ease the processes involved in talking to humans. It includes powerful routines to get and parse command strings as well as text-manipulation routines such as case conversion, context-sensitive string comparisons, and so forth. The powerful parsing calls make it easy to create a tiny machine-language task that includes a complete command interpreter and syntax error handler. This is very important if a significant amount of development is to be done in assembly language.

The trap manager handles all system traps except I/O interrupts, which go directly to the I/O manager. Error traps always cause the system to come to a complete standstill. This is important to us because of the close interaction between the tasks, which must always be in intimate communication to fulfill the purpose of the system. When an error trap occurs, all task switching and interrupt processing stop and the task in which the error occurred takes over the system console and enters the debugger. The human in charge can then take corrective action and restart the system with minimal damage. Obviously this approach would be completely unacceptable in a commercial operating system, but for us, it is ideal because we have all the source code for the entire system and can often correct bugs as soon as they occur.

No assembly language development system is complete without a debugger, of course, especially a multitasking one. Our debugger is a command parser that any task can invoke, either automatically (in response to an error trap) or directly. It is capable of running alongside other active tasks, even multiple copies of itself, and it allows the user full manipulation and examination of memory.

The heap munger is a distinct task, always present, always active, whose original job was to survey the contents of the heap continuously and maintain it as an efficient data structure (by using a background task for this, we avoided many complexities). The heap munger has turned into quite a bit more than just a trash compactor, however. Its responsibilities are many and varied, from checking the TCB array for in-

tegrity and waking up sleeping tasks when their ships come in to responding to messages from other tasks that want to know what the system loading is so that they can adjust their own CPU usage to increase the overall performance of the system. In fact, the heap munger is also capable of responding to text messages sent to it by a human who is operating a user task, in which case it responds by sending a plain English message back to the source task, which then displays it for the human to read. As a general rule, the heap munger performs any systemwide activity that must be performed at frequent, regular intervals. With all these responsibilities, it is the largest single code

segment in our kernel, weighing in at about \$A00 bytes.

The disk munger, like the heap munger, is a distinct task that is always running except when it's waiting for an I/O completion. Because its structure and function are application specific, I won't go into it in great detail here. I would like to point out, however, that by allocating a single task to handle each disk device, a large number of problems related to data contention between tasks can be avoided. The disk munger is entirely message-driven: As each task requires a disk access, it sends a message to the disk munger for the device it wants to access. Each I/O request gets added to a queue of such

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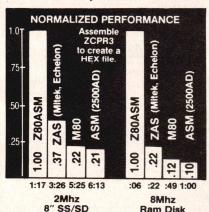
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MULTITASKING OS (Continued from page 51)

messages. When the disk I/O is completed, the disk munger sets the "wake-up" flag for the requesting task, and the heap munger wakes it up on its next pass. The requesting task then looks in its TCB for the completion code from the disk munger. Because the I/O requests are highlevel calls, usually implicitly including the open, read/write, and close of the file, there are few file contention problems.

On start-up, a discrete task called the system console manager initially gains control of the system console and puts up a system monitor prompt. Until another user task is spawned, it is the only task that has access to a character I/O device (and thus, to a human). The system console manager can then give way to any other user program.

### Tasks

Everything that gets done on the system, with the single exception of the context switch between tasks, is done by a task. Three tasks are always present in the system: the system console task, which initially runs the system console manager; the heap munger; and at least one disk munger. The disk munger is actually optional, although it's hard for me to imagine doing any useful work without using a disk.

Each task, whether it's active or not, possesses exactly one task data item (TData item) in the heap. The TData item is crucially important to the proper functioning of the system because it's where each task stores its most important local information. Every task's TCB entry contains the master pointer to its TData area in the heap. During normal execution of a task, the 68000 register A5 always points to the base of the TData item. Before the context switcher passes control to a task, it always sets up register A5 for the incoming task.

The TData item is also the location of the task's local data stack: The stack goes from the top of the TData item down, and the task's local data goes up from the bottom. Note that it is the responsibility of each task to ensure that its stack and TData areas do not collide.

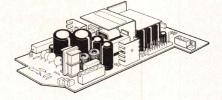
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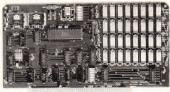
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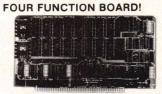
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### MULTITASKING OS

(Continued from page 52)

Each task's TData item contains spaces for various pointers and vectors associated with its current character I/O device, if any. The vectors include various standardized routines such as InCheck, which checks to see if a character is available; InWait, which waits for a character and returns with it; OutCheck; OutWait; plus several others. The pointers include the addresses of the destination for the next input byte (if any), the source of the next output byte, and so forth.

Character input is generally interrupt-driven. Because a task that is awaiting input is often "asleep" (not in the active TCB list), it is necessary for the interrupt routine to set a flag that tells the system to wake up the owner of the device. Then, after the interrupt has been handled, the heap munger, which is always awake and active, catches the set flag the next time it gets control and performs the actual manipulations to return the task's TCB to the active list. In order to

provide an input time-out, the heap munger also awakens each task once every ten seconds so that the task can test for the time-out

Each task may request to be assigned to a character I/O device. If a device that is requested is currently assigned to another task, the requestor will usually have to wait until the device is available. For special cases such as error traps, however, there is a system call that allows the task to demand to be connected to a device. When a demanded device is released, it reverts to the task (if any) to which it was attached originally. When a requested device is released, it always becomes available (unattached) again. Some devices can be attached to multiple tasks at the same time-for example, there is a device called MemPrt that reads and writes to memory as if it were a character stream from/to a serial device. This "virtual" device can be simultaneously active for different tasks, each with its own set of pointers in the TData item for reading and writing.

I won't go into great detail about our I/O drivers or the low-level structure of the I/O routines. This sort of information is readily available in several forms in computer bookstores.

## The Good Stuff: Tricks and Shortcuts

The first thing our operating system does when it gains control of the processor is to deviously remove the existing operating system. This it does by using a trick to get into supervisor mode (see the listing, page 102). First, it shoves a new address into the privilege exception vector, which is in the hardware vector table in low memory. This address happens to be that of the second instruction following the one that does the store to the vector. The next instruction is a privileged but otherwise harmless one. Thus, when it tries to execute it, it traps to the privilege exception vector and thence to the next instruction in the program. If through some quirk we happen to be in privileged mode already, the processor harmlessly executes the privileged instruction and falls through. Now we are in privileged mode, and we can quickly grab the rest of the system.

Next, we turn off all the interrupts in the system as quickly as possible. It is important to do this before clearing memory because a stray interrupt might happen before we can turn it off and it must still vector properly. After all interrupts are out of commission, we copy our own set of vectors into the interrupt table.

The jump table (discussed later) is next moved into place in low memory. It is read from a section of object code within another assembled module. It's nothing more than a sequence of more than a hundred absolute-long JMP instructions.

Next, we clear the rest of memory—except the kernel, of course—to zeroes. This is both a general safety measure and an aid in debugging: If a chunk of memory is nonzero, we can be certain that something we did caused it to be that way. Also, it's nice to be able to assume that unused memory is always zeroed out; it makes for much faster initializations later on, once the system is running.

The system zone requires a certain amount of initialization. The task control blocks must be set up and the end marks for the TCB and master



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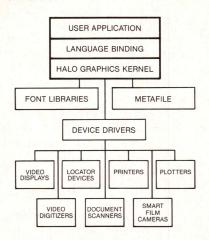
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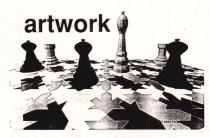
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This picture was created with "Artwork", a HALO-based application written by West End Film of Washington, DC.

### **De-Facto Standard**

First released in 1982, HALO has an established base of over 40,000 end users and over 100 corporate clients. Numerous HALO-based CAD, solids modeling, presentation graphics, art packages, mapping and other vertical market software applications are commercially available.

As early as June 1984, PC World featured an article entitled, "HALO A new software library leads the way toward graphics standardization and portability." The July 9, 1985 issue of PC Week" featured users stories from both Rockwell International Corp., and Lawrence Livermore Labs on how they use HALO to save development time and money. Most recently, Mini-Micro Systems August 1985 issue stated, "Widely used, HALO has attained de facto-standard status. It's certainly the most widely used library." HALO has achieved this status because it provides a complete device independent graphics environment for software developers. Since the HALO interface rarely changes, compatibility with a new device is achieved simply by adding a new device driver.

By incorporating all the HALO device drivers into their applications, licensed commercial software developers solve distribution and compatibility problems.



This picture was created with "Design Board 3-D", a HALO-based application written by MEGA-CADD of Seattle, Washington.

### **Powerful Graphics Toolbox**

HALO has over 170 functions. Standard functions include: point, line, arc, circle, ellipse, polygon, etc. Coloring control functions are available for color selection, palette management, dithering, and textures.

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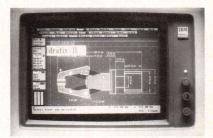
There are currently over twenty bit mapped and stroke text fonts available for HALO. The object oriented type faces provide user definition of line width, size, proportional width, angle, and interior fill style.

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In benchmark tests by independent companies evaluating HALO and GKS, HALO was determined to be three to ten times faster depending on the function tested. These benchmark results made HALO a requirement for many government agencies and large corporations. This gives HALO compatible hardware and software significant inroads into these agencies and corporations.



This picture was created with "drafix II", a HALO-based application written by FORESIGHT Resources Corp. of Overland Park, Kansas.

### Easy to Learn

HALO comes with a computer aided tutorial. This tutorial, entitled LEARNHALO, teaches users to become proficient graphics developers in less than a day by providing stepby-step instructions through graded examples.

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Graphics hardware manufactures receive distinct marketing advantages by having their products supported by HALO. HALO is the gateway to an extensive body of commercially available software.

### **Pricing Information**

A single copy of HALO for one language costs \$250.00. A second language binding ordered at the same time is an additional \$150.00. Author, Corporate and Site licensing and OEM agreements are negotiated on a case by case basis

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335-D Washington St., Norwell, Mass. 02061 617-659-1571 MULTITASKING OS

(Continued from page 54)

handle arrays must be set in place. The values in the miscellaneous system data area must also be initialized.

After the system zone is in place, the heap is defined from the next 32-byte boundary to the end of memory. Three heap items are initially set aside: a deletion from the beginning of the heap to the start of the kernel's code, a fixed (immovable) code item for the kernel, and another deletion from the end of the kernel to the end of memory. Soon, the other tasks will be carving up the big deletions for their own use.

Each time a task is spawned, the task manager creates a new TCB and a new TData item. The first task to be spawned is the system console manager. The system task manager allocates an item of the proper size from the heap, creates a TCB for the new task, and adds the TCB to the current linked list of active tasks' TCBs.

After the system console task is spawned, the heap munger and disk munger are also spawned. Note that none of them begins to execute until the initialization code jumps into the middle of the context switcher.

We discovered that one of the biggest sources of "extra" system overhead in commercial operating systems is the need to manage tasks that raight possibly get out of hand and take over the system. In most cases, such "runaway" tasks are avoided by using a hardware timer to assure that any given task will only be able to run for a preset time. If a task spends too much time without releasing to the system, the timer interrupt occurs and vectors the CPU through to the supervisor program. This program then saves the existing task's registers and status and restores those of the next task in line, which then is off and running with a newly reset timer. All of this requires a lot of tricky and complex code to ensure that no task can "run away" and lock up the system. Also, the constant saving and restoring of registers and status, which is necessary because the context switch is interrupt-driven, adds measurably to the overhead required for a context switch.

We debated for some time about the best way to reduce this overhead.

Finally, we settled on the answer: We developed our kernel as a nonpreemptive task controller. This means that there is no hardware timer to interrupt each task after some preset interval, and no routines are required to service such an interrupt. Instead, we use a simple context switcher, one that doesn't even bother to save registers or the previous task's status bits. It is the task's responsibility to call the context switcher often enough to ensure smooth system operation and to make sure that it saves any registers that it needs (other than A5 and the stack pointer). For our application, this is perfect because most of our tasks spend most of their time waiting for input or output in the inactive task list, during which time the other tasks can run unhindered.

When a task has finished with the CPU and is ready to let the next task in the active list run for a while, it simply calls the context switcher as a subroutine using a JSR instruction to a low-memory JMP instruction whose address is fixed regardless of the location of the context switcher. Even with the added overhead of the extra JMP instruction, this method of calling is considerably faster than a TRAP instruction—the usual method of calling a context switcher.

Once the switcher has control, it checks to see if the system tasking is stopped (see the listing). If it is, then the calling task immediately gets back control through a simple RTS. Otherwise, it gets ready to call the next task.

The address of a task's TCB is in its TData area. This address is loaded into A0. Then the current TData base address (in register A5) is subtracted from the stack pointer, yielding a relative displacement, which is stored in the TCB. Now we're ready to move on to the next task.

The address of the next task in the circular linked list of active tasks is fetched from the old task's TCB into A0. Register A5 is set to point to the new task's TData area by moving its address from the TCB, and the stack pointer is restored from the relative displacement by adding A5 to it. Then a simple RTS returns control to the task.

This otherwise trivial scheme has one slight complication: Frequently, a task will release control with the intention of going to sleep for a while. This happens, for example, when the RAM buffer has no input characters for the character device attached to a task that is waiting for input. When the input finally happens, the interrupt routine sets a flag that causes the heap munger to wake up the task when it next checks for such a situation. The problem is that the context switcher I have just described has no provision for putting the old task to sleep: It assumes that both tasks want to stay awake.

So, we created an alternate context switcher that removes the outgoing

task from the active list before calling the new one (see the listing). When a task wishes to go to sleep, it simply calls the alternate context switcher. The primary difference is that before it moves on to the next task, the alternate switcher does some standard and fast list manipulation.

Most commercial operating systems use one or more of the TRAP instructions to perform system calls, usually going on the premise that they are there for that purpose and that the TRAP instructions allow programs to be more general and more easily relocated. Unfortunately, TRAP instructions on the 68000 cause a ma-

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57

### MULTITASKING OS

(Continued from page 57)

jor problem: They take a long time to execute, as do the instructions to decode their arguments. The RTE instruction, which returns from a TRAP, shares the same problem.

We could understand the importance of relocatable programs, certainly, but we felt the long-lasting TRAP instruction was too much to ask of an ultra-fast, real-time system. We therefore designed a different, faster way to call system routines: the JSR

instruction and an old-fashioned jump table in low memory (see Table 2, page 49, for a timing comparison). Instead of using a TRAP instruction with an argument following in the next word, we simply call one of many entry points that are at absolute locations in low memory. Each entry point consists of a single JMP instruction to the actual system call entry point. Not only do we save the extra time required for the TRAP and RTE instructions but we also avoid having to extract the argument word from the bytes following the TRAP in-

struction and having to add two to the return address to jump around the argument because the argument is implicit in our choice of which routine to call. (With TRAP-based system calls, an argument is required to specify which system call to use because there are only 16 traps. With JSR calls, there can be hundreds of separate entries, so no argument is required to specify which call is to be used.) By using subroutines instead of traps, we shaved more than 100 machine cycles from every single system call, which makes a measurable difference in a machine that uses lots of system calls.

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### Conclusion

With the two context switchers just described, a small set of carefully designed system routines, a somewhat unusual system calling procedure, and a certain amount of cooperation from the application programs, we have vastly increased the throughput of our system. Our approach is obviously not well suited to most projects as it requires a considerable amount of skill and cooperation on the part of the programmers. Furthermore, because of its nonstandard nature, it is poorly suited to any applications that are written for commercial systems—at least until we get a C compiler running! If you need an extremely fast multitasking system for a specialized real-time application and are strapped for funds, however, this approach can turn a relatively inexpensive microcomputer into an amazingly powerful system. To date we have done just that for four different hardware configurations.

I'd be happy to discuss the philosophy and tricks in greater detail with anyone who is interested. For more information on our operating system, write to Nick Turner at the following address:

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# Bringing Up the 68000— A First Step

Adapted by permission of the publisher, Prentice-Hall Inc., from the forthcoming book Designing and Troubleshooting a 68000 Microcomputer System by Alan D. Wilcox, available 1986. No part of this adaptation may be reproduced, in any form or by any means, without written permission from the publisher.

here are two ways to bring up a new 68000 microcomputer system: the hard way and the easy way. The hard way is to use the traditional approach of designing the hardware and then using a development system along with test software and some in-circuit emulation. Given enough hours of testing and correcting problems, the 68000 system has a good chance of running successfully. In contrast, the easy way is to design, build, and test the hardware module by module using the 68000 as a free-running processor.

The impact of the free-running technique on hardware development is quite startling. The 68000 kernel shown in Figure 1 (page 61) can be made to run so easily that a logic probe can test it. You don't need to use sophisticated digital development tools such as a logic analyzer or a development system with an in-circuit emulator. If troubleshooting is necessary, you need only a common dual-trace oscilloscope.

Free running the 68000 means that the processor is allowed to execute a do-nothing instruction continually. This is accomplished by breaking the by Alan D. Wilcox

The impact of the free-running technique on hardware development is startling. The 68000 kernel can be made to run so easily that a logic probe can test it.

normally closed loop between the 68000 and its memory, as shown in Figure 2 (page 61). Instead of carrying program instructions from memory, a one-word instruction (call it a NIL instruction) can be jammed onto the data bus. (The mnemonic NIL is not part of the 68000 instruction set per se; I coined it as a simple expression of the instruction used for free running.) Thus, when the 68000 reads the data bus for an instruction, it fetches the NIL word, executes it, increments the address, and reads the next NIL. This cycling repeats over the entire 16-megabyte address range; when the processor reaches the end of the 16 megabytes, it simply starts over again.

The strategy for bringing up the 68000 is to design, build, and test the 68000 kernel shown in Figure 2. Next, design and build one additional module, connect it to the kernel, and test it while the processor is free running. Add yet another module and test it while free running. You can free run the 68000 all the way through the construction of a complete CPU board. In fact, if a processor

board fails, you can usually free run it to help speed troubleshooting. The only part of the system that you cannot test easily while it is free running is the data bus itself, because the NIL instruction is forced on the bus.

### The Steps

Several steps are involved in bringing up the 68000 using the free-running technique. The intent of this article is to describe the necessary first step: how to get the kernel running. Once you have the kernel in operation, the rest is fairly straightforward. Here is a brief overview of the entire scenario to complete a working 68000 CPU board:

- 1. Bring up the kernel. Design, build, and test the power system, the 68000 clock and drivers, the reset and halt module, and the 68000 module.
- 2. Add a wait state and data transfer acknowledge (DTACK\*) generator module.
- 3. Add RAM and EPROM decoding circuits, connect address and control bus circuits.
- 4. Write a simple looping program for a pair of EPROMs. Remove the NIL instruction and close the broken loop between the EPROMs and the 68000 data bus. The processor should now be able to read its stack and program counter vectors from the EPROMs and execute the loop program.
- 5. Add the RAM connections to the data bus. If the 68000 is still running successfully with the simple loop program, add more code to include reading and writing RAM. If the code is a very tight loop, an oscilloscope will synchronize easily and you can use it to check the timing of the various control lines to all the memory in the system so far.

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- 6. Modify the reset and EPROM-control circuits so that the EPROMs do not have to be decoded at address 0 except during reset. Normally, the low memory addresses should be RAM so that exception vectors can be altered dynamically; EPROMs should be elsewhere.
- 7. If at least 4K of RAM has been decoded starting at address 8, and the EPROMS are decoded for 0 to 7 and \$8000 to \$BFFF, then the system can use the Motorola TUTOR EPROM set. When you restart the system, assuming that it does not halt, you can use an oscilloscope to see the activity on the various processor lines while the monitor waits for a console keystroke.
- 8. Add a 6850 ACIA decoded at \$010040 to serve as a console port and test it with the TUTOR EPROM set. Run various memory-testing commands and scope loops to check operation of the new system.

### The First Step

As stated earlier, the first step is to bring up the kernel in the free-running mode. It seems a bit overwhelming when you first try doing it, but it really is quite simple. Unless you have made a wiring error, the 68000 is virtually guaranteed to come alive and begin executing the NIL instruction. To bring up the kernel, you need the power system, the clock and drivers, the reset and halt module, and the 68000 module.

The size of the power supply depends on the nature of the system and what loading it will have in the final configuration. In my own case, I intended to use the 68000 processor board in an IEEE-696 (S-100) system, so I needed on-card regulation from an 8-volt system supply. A common 7805 circuit was adequate for the processor and its RAM and EPROMs; I used a second 7805 circuit for the rest of the LS-TTL logic on the CPU board.

Watch the Motorola data manual for footnotes. In the case of the 68000, although the data indicates a power requirement of 300 mA or so, that is not the whole story. The fine print at the bottom of one page casually mentions that the 68000 might require a peak current of some 1.5 A. Make sure the power supply can handle the peak current without falling out of regulation. Likewise, power and

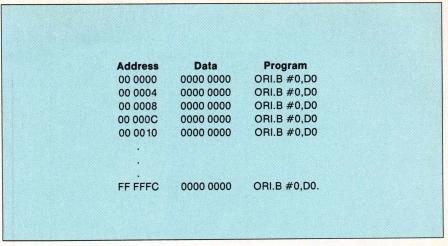


Table 1

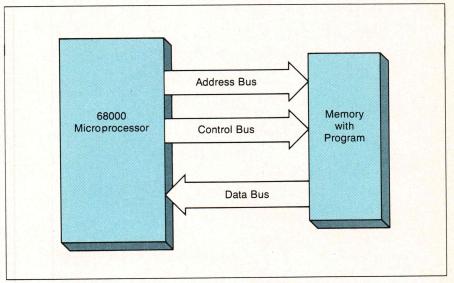
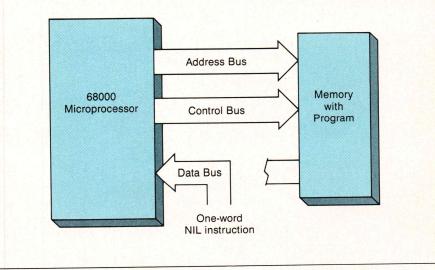


Figure 1: The 68000 kernel is the essential hardware for program execution.



**Figure 2:** To free run the 68000, the normal feedback path from memory is disconnected, and a NIL (do-nothing) instruction is substituted.

ground leads to the 68000 need to be heavy, say #24 wire rather than #30 wire-wrap wires. Locate bypass capacitors close to each of the power and ground connections.

You can design and build the clock

circuit using a crystal, some resistors and capacitors, and a 7404 or similar; doing this is hardly worth the effort, though. For prototype work, being able to change the clock frequency easily without redesigning the circuit is a distinct advantage, so using a DIP oscillator is appropriate. I used a 6-MHz oscillator in my S-100 proto-

type to keep within the bus specification; only after I had finished the system did I run it up to 10 MHz and later to 12 MHz.

Also, you will use both the clock signal and its complement in the final circuit design. The complement clock could be derived from a 74LS04, but that would introduce a skew be-

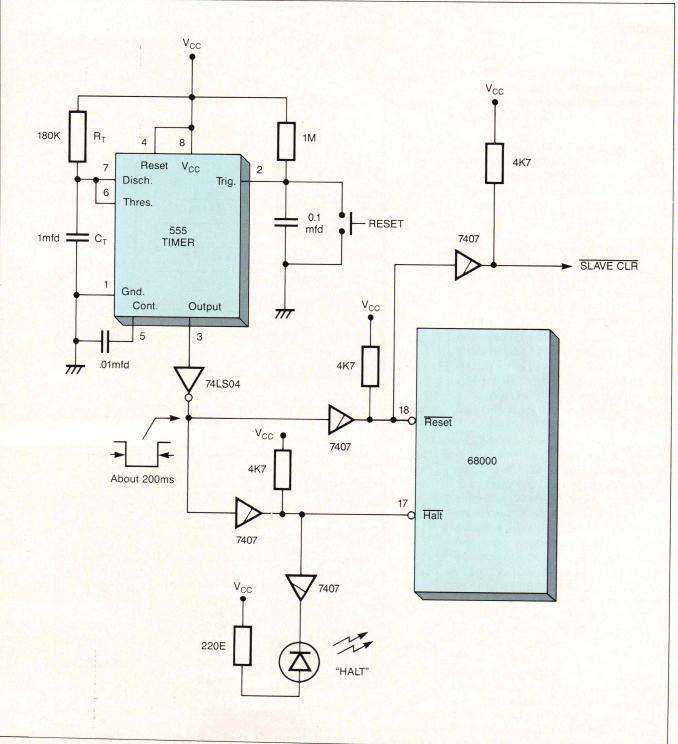


Figure 3: Circuit diagram of a simple power-up and reset timer circuit for a 68000 processor. Note the use of open-collector devices on the bidirectional HALT\* and RESET\* controls of the 68000.

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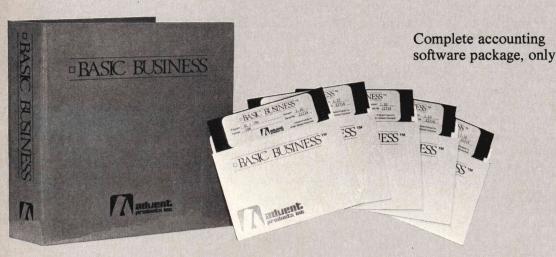
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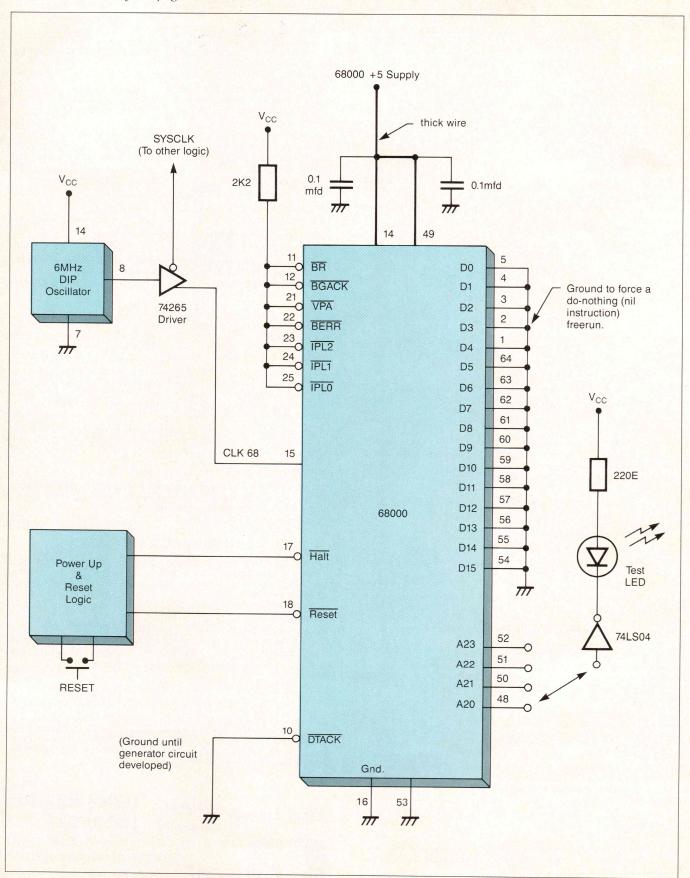


Figure 4: Circuit diagram of the minimum 68000 system for free-run test

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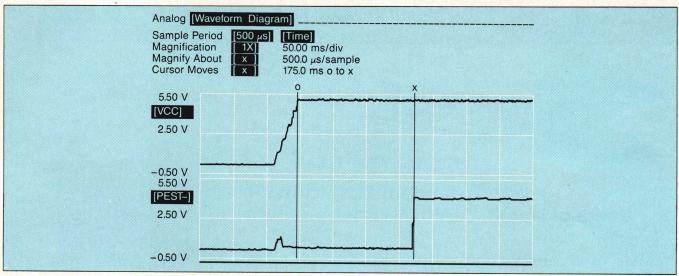
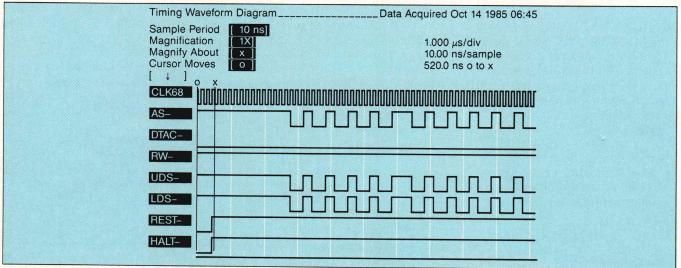


Figure 5: Power-up performance of the 555 timer circuit. On power-up, the 555 timer with the parts given in the schematic provides about 175 ms RESET\* to the 68000.



**Figure 6:** Typical free run starting from a complete RESET\* and HALT\* asserted low. The clock is running at 6 MHz. DTACK\* is grounded in this example.

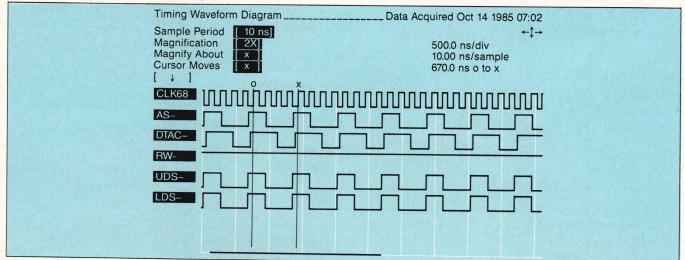
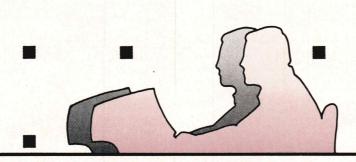


Figure 7: Typical free run with the DTACK\* circuit enabled. The clock is 6 MHz, and there are no wait states inserted.



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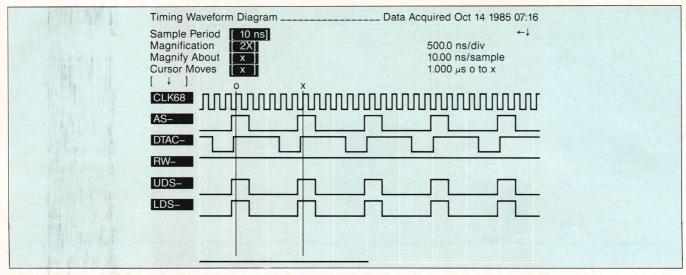
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**Figure 8:** Typical free run with DTACK\* enabled. This timing shows DTACK\* delayed enough to cause two waits in each bus cycle.

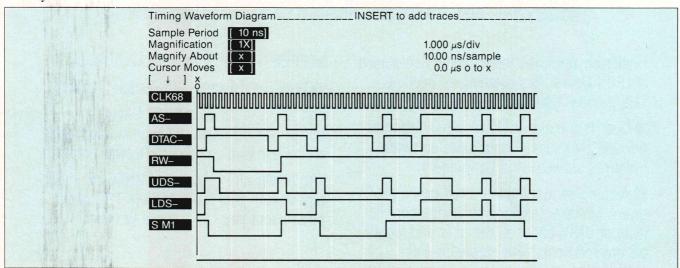


Figure 9: A view of the bus activity when the TUTOR EPROM set runs at 6 MHz. The CPU board was set for eight waits on I/O and three waits otherwise.

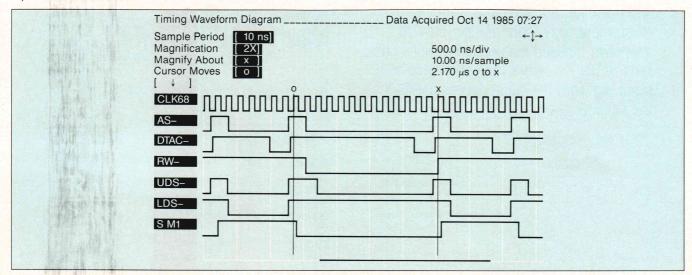


Figure 10: A closer look at the bus controls when TUTOR executes a write bus cycle

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68000 (Continued from page 62)

tween the two clocks of some 10 to 15 ns, depending on loading. Although this amount of skew seems slight, it can cause severe timing difficulties when the clock speed gets above 10 MHz. The 74265 quad complementary-output element with a worst-case skew of 3 ns is a good selection; in my 12-MHz prototype, this selection has worked out well.

The reset and halt module has two basic functions. One task is to hold the 68000 HALT\* and RESET\* lines low for at least 100 ms on power-up. Its other function is to pull the same two lines low for at least ten clock cycles for a reset button press at any time after the power has been on.

The circuit in Figure 3 (page 62) provides a simple and reliable reset function for the 68000. It provides a reset pulse either on power-up or whenever you press the reset button. Open-collector devices are required because the 68000 HALT\* and RESET\* controls are both bidirectional. For example, the 68000 can itself drive the RESET\* line to reset any peripherals if the software reset instruction is issued. Also, the 68000 can force the HALT\* line low if the system cannot continue processing. A single "halt" LED connected as shown is valuable in helping bring up the processor for the first time.

The last module in the minimum system is the 68000 processor shown in Figure 4 (page 66). By now, you should have checked the power, clock, and reset modules for proper operation and connected them ready for the 68000. If the processor is wired as shown, it should begin free running immediately. On power-up the HALT light should flash briefly, and then the TEST light will begin flashing on and off.

Earlier I referred to my so-called NIL instruction. As you see in the circuit, the data bus is completely grounded so the NIL has an opcode of 0000. In the context of its use in free running, it acts like a no-operation or NOP. The 68000 does have a NOP opcode (\$4E71), but this NOP will not work as a free-running instruction.

A critical constraint on the opcode precludes using the NOP instruction in free running: Whatever is wired to the data bus for the 68000 to read upon reset must be even. The reset sequence is this: The 68000 will do four 16-bit reads to get the initial SSP and PC vectors; then it will fetch its first opcode at the address in the PC. If the PC is not aligned on an even address, the 68000 detects an address error and immediately begins illegaladdress exception processing. It tries to push its status on the stack at the beginning of the exception, but the stack is also an illegal address (the same noneven number as in the PC). The result is the fatal double bus fault that stops all processing and asserts the HALT\* output.

The opcode 0000 does in fact correspond to a real instruction in the 68000 set. It is the mnemonic ORI.B #0,D0, and it was selected for free running for two reasons: first, because it was even; and second, because connecting all grounds to the data bus was simpler than making sure one or two data lines had a logic 1 on them. When the instruction is considered in its free-running environment, the appearance of its memory is as shown in Table 1 (page 61).

You can calculate the execution time of this "program" easily. Each instruction takes eight clock cycles (two read bus cycles), so for a 6-MHz clock, the execution time is 8×167 ns or approximately 1.33 microseconds. A complete sweep through the entire 16 megabytes of the 68000 address range takes 1.33×4 megabytes or about 5.59 seconds. If you connect the TEST light to the top address bit, A23, it will be on for 2.8 seconds and then off for 2.8 seconds. I connected the TEST light to A20 permanently. It stays on for 0.35 seconds and off for 0.35 seconds—a reassuring flash rate during development work and not nearly as unsettling as a constant red HALT light.

#### Results

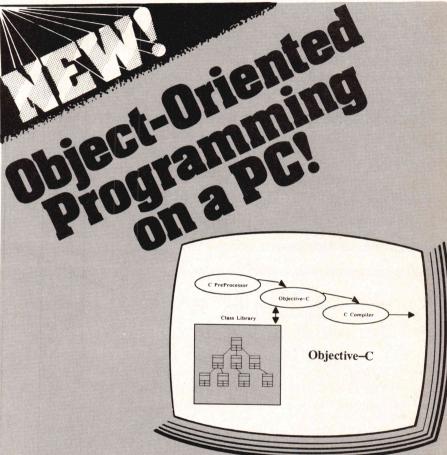
Figure 5 (page 68) shows the performance of the power-up timer circuit. The top plot is the main system power as it comes on and eventually regulates at 5 volts on the CPU board. About 175 ms after the supply voltage is valid, the RESET\* and HALT\* lines go high to successfully complete a full 68000 reset.

The effect of this reset operation is shown in Figure 6 (page 68). The last

two lines on the timing diagram are the RESET\* and HALT\* controls for the 68000. After its internal start-up time, the 68000 asserts its address strobe (AS\*) and its data strobe lines (UDS\* and LDS\*) in the first read bus cycle. After four read bus cycles, the processor PC begins execution at address 0, as discussed above.

DTACK\* is the asynchronous bus

control line that normally comes back from memory or peripherals to tell the 68000 to complete the current bus cycle. During the initial free run of the processor, there is nothing connected that will acknowledge a data transfer, so the control is grounded. The timing diagram shows this line at a logic low. The timing diagram also shows the read/



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(Continued from page 73)

write control, R/W\*, as constantly high because the 68000 only does a read bus cycle when it is free running.

Figure 7 (page 68) shows the freerunning processor with a DTACK\* generator in operation. Notice the o and x markers bracketing a single bus cycle. The normal read bus cycle has a total of four clock cycles. If DTACK\* is delayed for two cycles, as shown in Figure 8 (page 70), then the bus cycle is lengthened and two "waits" are inserted into each bus cycle. When you interface memory or peripherals to the 68000, you can design each external module to hold back DTACK\* until its unique timing requirements are met.

As an example, the timing diagram in Figure 9 (page 70) shows the system while *not* free running: It is executing the monitor program (TUTOR) and waiting for a keystroke. The system was set to provide eight waits for I/O read operations, nine for writes, and three waits otherwise. Figure 10 (page 70) shows a close-up look at the bus cycles. The lower timing line, marked sM1, is the S-100 bus status indicating an opcode fetch.

Summary

Bringing up the 68000 using the freerunning technique is very different from the more traditional approaches to getting a processor running. You can see, though, just by the brief description of this first step in bringing up the 68000 kernel, that you do not need sophisticated equipment to get started.

There is more to be said about all the steps beyond this first free-running processor; no doubt many questions remain unanswered. From my experience, though, the understanding you can get from doing a free-running 68000 is very valuable, and it can help you go on to design and build a complete system successfully.

Availability

The TUTOR firmware is available directly from the author.

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### COM: An 8080 Simulator for the MC68000

fter I bought a 68000-based S-100 system, I found that I needed to run several CP/M 2.2 (alias CP/M-80) software packages, none of which were available in equivalent forms for my machine. I was faced with two options-either buying another processor board so that I could run these programs or writing an 8080 simulator. The software approach seemed infinitely preferable. This also seemed like an opportune time to find out just how fast my 68000 really was. So I set out to write an 8080 simulator (which I named COM because it interprets .com files).

In principle, writing a simulator is simple. You just set up a set of fake registers and start picking up opcodes and interpreting them. Unfortunately, it's the little details that get you. This simulator took about twice as long to write as I expected. As it is, it isn't perfect. In fact, it would slow down considerably if it were. However, most programs aren't bothered by the imperfections, and the speed difference would be significant because the simulation is already on the slow side of usable. I wrote COM to run as fast as possible.

The simulation speed is approximately that of a 1.4-MHz Z80 processor based on a sample assembly with MAC. (My 68000 system is an 8-MHz CompuPro/Morrow single-user hybrid running CP/M-68K 1.2, with 16-bit no-wait-state memory.) Simulation time can vary widely, as some 8080 instructions aren't easily simu-

Jim Cathey, ISC Systems Corp., TAF-C8, Spokane, WA 99220.

#### by Jim Cathey

Simulation time can vary widely, as some 8080 instructions aren't easily simulated with the 68000.

lated with the 68000. MAC was chosen as a typical example program. LU, the public-domain library utility, is one of the worst performers. It spends vast amounts of time calculating CRCs. This instruction sequence isn't very efficient on the 8080 (using a C arithmetic library), and the simulation magnifies any inefficiencies. I didn't time LU for comparison, but I thought that the simulation had crashed because nothing seemed to be happening for long periods. Another poor performer is WordStar. I tried simulating it because it is so popular and because it gives the simulator a thorough workout. The simulation works, and it does so at an acceptable performance level most of the time.

#### About the Program

The program source is broken into four files. Listing One (page 104) is the first file, which contains the start-up, command-line, CP/M-2.2 simulation, and trace routines. Listing Two (page 112) contains opcode simulation subroutines and flag tables. The rest of the listings will be continued in the March issue.

The program starts out by prompt-

ing for the trace end points if that code has been included. (There are several conditional assembly trace features in COM.) It then builds the 8080 environment in a 64K buffer (biggest TPA yet!) and initializes the 68000 simulation registers, with the 8080's PC set to 100H into the buffer. It then calls a subroutine that loads the specified .com program into the buffer and transfers the 68000's second FCB to the 8080's first FCB and passes the remains of the command tail to the 8080's DMA buffer. If all goes well, the program then enters the main loop at label MLOOP. Here it fetches the next 8080 opcode and uses it to index into a table of 256 68000 subroutines (one per opcodebig, but fast) and jumps to the selected subroutine. Each opcode subroutine then picks up what parameters it needs and plays with the fake registers appropriately. Each subroutine then jumps back to MLOOP, which repeats the process. This continues until a service request is picked up or until an illegal instruction is found.

The service request used by the simulation is the HLT opcode (\$76). HLT is followed by a 1-byte parameter telling the 68000 which action to take. All BIOS/BDOS functions are implemented as service requests. After the service is performed, execution continues at MLOOP (or at the byte following the parameter—it depends on your point of view). Refer to the 8080 Environment section for more details on how service requests are used by the 8080's BDOS/BIOS.

Register dumps are caused by illegal opcodes or are done during a trace. They are easily interpreted (registers S0-S3 are the top four en-

tries on the stack) as the current instruction is also disassembled. Illegal opcodes terminate the simulation after the register dump.

Flag simulation is done with two tables. Because any 8080 logical operation that sets the parity flag also clears the carry bit, these flag results are based solely on the value in the accumulator. The simulator uses a 256byte flag lookup table for these operations. Similarly, anything that conditionally sets carry (an arithmetic operation) doesn't need to set the parity flag if the code is intended to run on a Z80. (This describes all CP/M-2.2 software I was interested in.) Another 16-byte table can therefore be used for arithmetic flag results. The 4-bit flag field of the 68000's status register is used as the index for this second table. The 68000 does have an overflow flag, so this is substituted for the parity bit of the 8080 (exactly as in the Z80). This causes one problem that is discussed in the Known Faults section. Treatment of the half carry bit is also discussed later because it doesn't fit into either of the tables.

The CP/M-80 environment simulation is greatly simplified (to my great disappointment) by the strong resemblance of CP/M-68K to CP/M-80. Most of the calls are directly translatable. There are a few exceptions, though, and they require the bulk of the code.

- 1. Any call referencing an FCB requires that the byte order of the Random Record field be switched, if the call uses that field.
- 2. CP/M-68K can't open a file in any but the base extent. You have to change such requests to an open in the base extent and then do a Random Read to the point you wanted.
- 3. Direct console I/O (BDOS #6) under CP/M-80 returns a null flag if no character is available. CP/M-68K waits for a character. A status check is performed first, and if a character is ready, it is fetched and returned to the simulation. Otherwise just a null status is returned.
- 4. Any call referencing an address (DMA or otherwise) needs to have that address translated to point into the 8080's code buffer. (This is a problem inherent to the simulation because the 8080 buffer cannot be placed in system memory at address 0.)

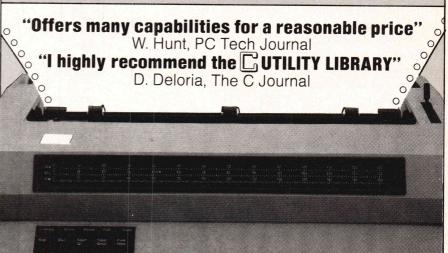
Some instruction simulations don't

do what you think they might. Specifically, the EI and DI instructions do not translate to the equivalent 68000 sequences. The object of using DI/EI pairs in an 8080 program is to prevent time-critical code from being interrupted or to prevent resource contention between interrupt routines and background processes that share resources. Because there are no 8080 "interrupts" possible under COM, there is nothing to block. (We won't even talk about simulating time-critical code!)

A few Z80 simulation routines are used in COM because of the overflow/ parity flag problem discussed below. These are just an extension of the table approach used for the 8080-a large jump table and a bunch of subroutines. Extension to a full Z80 simulation is straightforward but would require a lot of code if the present jump table technique is kept.

Another approach to simulating instructions involves dividing the opcodes into classes, e.g., all MOVs handled by one subroutine that figures





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out what to move where by examining the opcode in further detail. Though this is much smaller codewise than having the 63 similar subroutines that I used, it also suffers a speed penalty that I just couldn't tolerate. I didn't buy a 68000 just to slow it down!

#### **Known Faults**

There are two problems with the simulation of the 8080's flags. The first is that the flags are more like those of the Z80 than the 8080 in that the parity flag reflects overflow status after arithmetic operations rather than parity. This fools some dynamically selected run-time packages such as the one used for BDS-C. There is minimal Z80 support in COM to handle the few extra instructions that BDS-C wants to use (LDIR and LDDR) so that programs compiled by it will run. (CPIR was also required by another program for the same reason.) This Z80 support could be extended as much as needed-up to and including full Z80 simulation.

The other flag problem is with the half carry bit. Simulating it would take a lot of overhead because there is no similar flag in the 68000. Therefore, assuming the only need for the flag is for the DAA instruction, this instruction is treated specially. Instructions that set the half carry bit meaningfully (ADDs, ADCs, and INR As) have additional code to store away the two operands and CY (if used) in special locations. The DAA simulation then recreates the HCY bit out of these stored values. A problem can arise when the flags are pushed and then pulled before the DAA is executedan incorrect HCY is created if another addition-type operation occurs while the flags are supposedly saved. In practice, I only ran into this problem when using the 8080 DDT to trace through a DAA instruction. Be forewarned. If required, the simulation could be extended to eliminate this problem by proper simulation of the HCY bit, but this would slow the arithmetic routines down quite a bit.

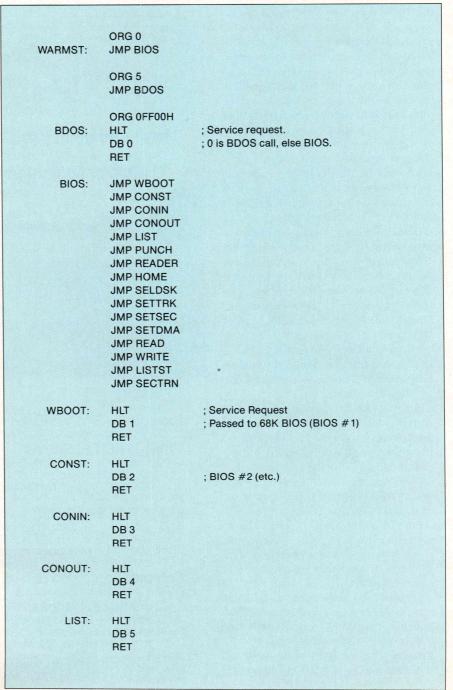
BIOS calls to the disk drivers aren't allowed. I did this for safety reasons more than anything else-I didn't want possibly buggy simulations Table 1

playing with my hard disk without the protection of my BDOS. This limitation could easily be removed.

BDOS call #31 (Get DPH address) isn't supported. I didn't need this for any program I wanted to use so it was left out. Including it would involve getting the table from the 68000, copying it to somewhere where the "8080" could find it, and then returning a pointer to this copy of the table. Similarly function #27 (Get ALLOC vector) isn't supported either. Using either of these calls will cause an abort and an appropriate error message.

The IOBYTE and LOGIN vectors at 3 and 4 aren't supported.

Only one parsed FCB is supplied in the 8080's base page. The normally present second name at \$6C isn't parsed. (Note that CP/M-68K parses two full FCBs when COM is invoked. The first is the name of the .com program to run, and the second is used as this program's first FCB. CP/M-80's second "FCB" isn't one—it is only an-



other name field in the first FCB. I didn't find anything that needed it, so I didn't go to the trouble of picking another name out of the 68K's command tail.)

An additional complication inherent to the simulation arises when you try to use programs running under COM to drive DMA devices. The hard disk controller in my system is a Morrow HD/DMA. In order for the 8080 simulation to be able to drive this controller, all addresses passed to the

board's DMA circuitry must undergo translation so that they point to the buffers in the "8080" program space, not the 68000's! Handling this tends to require a lot of code specific to each device supported, but it can be done. (COM was originally written for two reasons: to develop firmware for the 8085 used in my system's keyboard and to run the Morrow FORMATMW [hard disk formatter] program. Tracing this program pointed out the error in the HD/DMA documentation

that kept my own formatting program from working.) There is a conditional assembly flag in COM to include the support for the Morrow controller. You probably will never want this option, but I left the code in to serve as an example of extending COM to support a DMA device.

#### 8080 Environment

The 8080 Environment is a 64K buffer, of which all but 512 bytes are available as TPA. This is probably the only real advantage of the simulation over real execution. The fake BIOS/BDOS (FDOS) starts at 8080 address \$FF00 and is in the form of a jump table followed by a service request table. The warmstart and BDOS jumps in the low page of the buffer point to these tables. There is no CCP because COM takes its place. The 8080 form of the FDOS is shown in Table 1 (page 78).

The service request handler performs a BIOS call to the 68000 if the parameter following HLT is not zero or a BDOS call if the parm is zero. In either case the appropriate parameters from the fake 8080 registers are translated (if required) and passed to the 68000 FDOS. The return values are then translated (if required) and stuffed into the 8080's pseudoregisters, and the simulation is continued.

#### **Using COM**

CP/M-80 programs are run by inserting the word COM in the normal command line. Examples of use are:

A>COM WS TEST.ASM
A>COM MAC TEST
A>COM LOAD TEST
A>COM DDT TEST.COM
A>COM LU -O JUNK -A TEST.COM
A>COM LDIR JUNK
A>COM MBASIC FFT.BAS

COM may be assembled with several optional trace facilities. Normally I create a separate version called COMT because the presence of the trace code slows down the simulation. A trace is specified by giving the full normal command line and then answering the prompts for the start and stop trace addresses. COMT will check each address before it simulates the opcode at that address for a match with either of the two limits and turn on or off the register dump appropriately. An example is shown

#### (Continued from page 78) PUNCH: HLT **DB 6** RET READER: HLT DB 7 RET HOME: HLT ; Normally blocked by **DB8** : the simulation. RET SELDSK: HIT ; Ditto, etc. DR 9 RET SETTRK: HIT **DB 10** RET SETSEC: HLT **DB 11** RET SETDMA: HLT **DB 12** RET READ: HLT **DB 13** RET WRITE: HLT **DB 14** RET LISTST: HLT ; This one is allowed. **DB 15** RET SECTRN: HLT **DB 16** RET END

Table 1

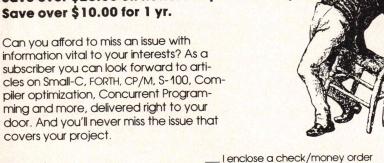
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#### 8080 SIMULATOR (Continued from page 79)

in Table 2 (page 81). The tracing addresses in this figure will trace every BDOS call made by the program.

The code for address prompting is rather stupid-it doesn't use linebuffered I/O. Because I almost never use tracing now that COM works, I didn't bother to fix this up. (The code was pulled out of my ROM monitor, where I didn't necessarily have any RAM for a buffer, which explains its strange structure.)

Other tracing code may be included in COMT. There is support for dumping (to the printer) FCB calls to the BDOS and for including a register dump at this time. All of these trace options were useful in debugging the simulation. You probably won't need them, but they illustrate one advantage of the simulation over the real thing-you can monitor events at any level of detail you want, provided you don't need real-time execution.

#### Teaching the Assembler Tricks (With a Hammer)

This section describes some of the tricks I used to make the ALCYON assembler (AS68) that is distributed with the CP/M-68K package do what I wanted when writing COM. Also described are some other tricks that I have used successfully in other assembly programs. (AS68 is also the assembler distributed with the Atari 520ST developer's package.)

At the beginning of Listing One are the register definitions used by COM. The form of these definitions was picked out of the AS68INIT file that you use the first time you run AS68. These definitions allow you to refer to the 68000 registers with more meaningful names than just "D3" and so on. The only real disadvantage of using names is that it is easy to forget which registers are in use and accidentally use one as a temporary that should have been saved first. Proper documentation helps in this. You also cannot use these new names in a "reg" directive.

One problem with these names (and all other "equ" defined symbols) is that you can't define them as global and use them in another file. The declarations must be repeated in each source file.

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- B. 500 to 9,999
  - C. 100 to 499
  - D. 10 to 99 E. less than 10
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- A. more than once per day.
- B. once per day.
- C. once per week
- D. less than once per week.
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253	254	255	256	257	258	259	260	26
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- B. design software.
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- D. don't design software or write code.

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Because of the conditional assemblies in COM, I needed to place some labels on lines by themselves, particularly MLOOP, the main looping point of the opcode simulation. Though the assembler allows this (provided you put a colon after the label), I was unable to make the label global using the ".globl" directive. However, if you do something like this:

.globl mloop mloop:

~ ~ mloop: ; Why? I don't know!

you can get the declaration to work. This trick was found by examining the code produced by the C compiler. I believe that if the label is in uppercase you don't need this extra statement. Note that

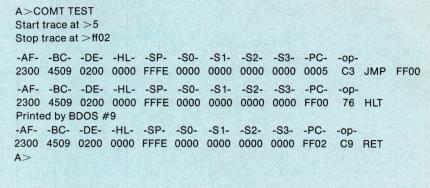
.globl mloop mloop: equ\*

won't work because of the problem with "equ" described earlier.

The "offset" directive is extremely useful for generating data storage areas to be used with indexing. It would have been useful in COM, but it doesn't work. "Equ" must be used, and the programmer must count bytes in the storage areas to make sure nothing overlaps. Grrrrr. "Offset" wasn't accepted by the CP/M-68K Release 1.1 assembler. The Release 1.2 assembler works just fine-but neither linker will accept symbols defined in offset sections. Another "feature" to fix. Doesn't DRI test anything?

Another trick that may be useful in 68000 assembler programming is using the ".opd" directive to generate your own opcodes. These are particularly useful as stand-ins for less meaningful ".dc.w constant" sequences when generating data tables, although you may define instructions also. COM uses this trick to define the 68000's BDOS and BIOS trap instructions. Look at the file AS68INIT for more examples of its use. The only restriction for use is that the new opcode must follow the addressing rules of one of the existing 68000's instructions. Oh, for a true macro assembler....

The C preprocessor, CP68, may be used as a poor man's macro assembler. If the assembly file (let's say TEST.MAC) looks something like this:



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#### 8080 SIMULATOR (Continued from page 81)

#include "MACRO.H" label testmac(1,2,3)

and the file MACRO.H looks like this:

- \* File produced by using CP68 as a macro
- \* assembly preprocessor on a .MAC file.

#define testmac(x,y,z) dc.w x \ .dc.w y \ .dc.w z

and if the files are processed like this:

A>CP68 TEST.MAC TEST.S ; CP/M 1.1

or

A>CP68 -P TEST.MAC TEST.S; CP/M 1.2

then the output file (TEST.S) will look like this:

- \* File produced by using CP68 as a macro
- \* assembly preprocessor on a .MAC file.

.dc.w 2 .dc.w 3 end

The <cr> lines are additional blank lines, one for each of the lines of a #define in the .H file.

#### Notes

The principal reference for the 8080 model was the MCS-80/85 Family User's Manual by Intel Corporation.

This program is released to the public domain with the stipulation that it be used for noncommercial purposes and that appropriate credit be given in any upgrade.

DDJ

(Listings begin on page 104)

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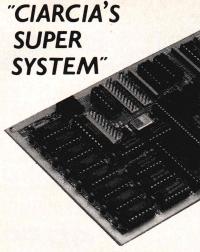
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#### C CHEST LISTING

(Text begins on page 18)

```
#include <stdio.h>
#include <dir.h>
#include <process.h>
#include <errno.h>
#include <fcntl.h>
        #include <signal.h>
#include <dos.h>
                 SH.C: A shell for MSDOS.
                                 Copyright (C) 1985 Allen I. Holub. All rights reserved.
  11
  14
                           This file contains an MSDOS shell that operates in conjunction with command.com. It has several built-in commands (see below). It recognizes several semi-colon seperated commands on a line & does wild card expansion. It can take commands interactively or from the command line (but command.com will intercept the re-direction if you have any on the command line). It supports simple batch files and will expand several $ arguments.
  16
17
18
  19
  20
21
22
 23 24 25
                            It does not support redirection yet.
 26
27
                                             sh [-cvx] <args>
 28
                  Invocation:
  29
                                                           shell entered in interactive mode. If -i is given any command line arguments that follow will be assigned to \$0 \$1 etc. \$0 will point at the leftmost argument following the -i.
  31
  32
33
  34
35
                                                           commands are read from string. If several strings are present they are concatanated together before execution. The 'c' may be upper or lower case.
  36
                  sh -c <string>
  39
                                                           . commands are taken from <file>. Args are expanded to correspond to $0 $1 etc inside the file. For DOS compatability %0 %1 etc are also recognized. $0 will be <file> itself.
  40
41
42
                   sh <file> args..
  43 44 45
                                                           Strip quotes from quoted argument strings. Usually they're left in so that a spawned process can assemble its argv correctly. Print input lines to the shell as they are read (same as "set verbose=1"). Print lines as they are executed. (same as saying "set echo-1" in the shrc.bat file.
                   sh -q
  48
                   sh -v
  50
                   sh -x
  53
               Enviornment variables:
CMDLINE (set)
                                                                          Holds the complete, 2048 byte, command line
that can't be passed via MSDOS.
Defines the prompt string. Any character
or $<arg> may be used. Default prompt is
  56
57
58
59
                            PROMPT
                                                                           [$s:$p].
Current shell level (0 is outermost). Remember that batch files are executed in their own
  60
61
62
                            SHLEV
                                                                           shell.
                                                                           Use first character in line to designate command line switches.
  63
64
65
66
67
68
69
70
71
72
73
74
75
77
77
78
80
81
                            SWITCHAR (used)
                                                                   Executed every time a shell is created.
Executed when a level 0 shell is created.
Executed when "logout" command is executed
                             /shrc.bat
                             /login.bat
/logout.bat
                Built in commands:
                                               [name [wordlist]] 
<directory name>
                            alias
                             cd
                            exit
history
                            logout
                             rem
                             setenv
set
                                               <name> [=] <value>
[{cmd|echo|verbose|$arg} [= val]]
                             shift
                             unalias <name>
  82
83
84
85
86
87
88
89
                                                 <name>
                             unset
                                             !<num> !<pat> ^<num> ^<pat>
                             !> [name]
                             !< [name]
            * Pre-defined shell variables (may use either % or $) :
  90
91
92
93
94
95
96
97
                                               one argument in batch file. $0 is file name.
All the $<num>s concatantated with spaces between them.
Full path name of current directory.
Current History number.
Nesting level of current shell. 0 is the outermost
                             $p
$!
  98
99
                 Special characters:
                                               Used to delimit two commands on one line.
101
                                                Name containing these is expanded to matching directory
103
                            Special characters aren't recognized in quoted strings or when preceeded by a backslash.
All lines with # in the left-most column are ignored.
```

```
107
108
         */
                                                                                             ); /* in stdio library
                                                         (char*, int
(int, int, int
110 extern int
                                        access
       extern int
                                                         (int,
(char*
                                        bdos
                                        chdir
                                       errno;
*fopen
fseek
113 extern int
114 extern FILE
115 extern int
                                                         (char*, char*
(FILE*, long, int
116 extern long
117 extern char
                                        ftell
                                                          (FILE*
                                        *getenv
*getcwd
                                                         (char*, int
118 extern char
119 extern char
120 extern int
                                        *malloc
                                                          (unsigned
                                       putenv
*signal
                                                          (char*
121 extern int
                                                         (int,int(*)()) )()
122 extern int
123 extern char
                                       strlen (char*
*strpbrk (char*, char*
124
                                                                                                        source is in:
/src/tools/cpy.c
/src/tools/dir.c
                                                         (char*, char* ); /*
(DIRECTORY * ); /*
(char*, DIRECTORY*); /*
(char*, int, FILE*); /*
(int ); /*
126 extern char
                                       *cpy
del_dir
dir
*efgets
127 extern void
128 extern void
                                                                                                        /src/tools/dir.c
/src/tools/efgets.c
       extern char
129 extern char *eigets (char*, int, Filb*); /* /src/tools/eigets.c 130 extern DIRECTORY *mk dir (int ); /* /src/tools/dir.c 131 extern char *next (char**, int, int); /* /src/tools/next.c 132 extern char *skipto (int, char*, int); /* /src/tools/skipto.c 133 extern char *strsave (char* ); /* /src/tools/strsave.c 134 extern int unargv (int,char**,char*,int,int); /* /src/tools/unargv.c
                                                                                                        /src/tools/strsave.c */
136 extern void
137 extern int
138 extern void
                                     unsetvar( char* );
setvar( char*, char* );
printalias();
                                                                                                  /* ./var.c:
139 extern void
140 extern int
                                     printvars();
getvar(char**, char**, int*);
141
142 extern void
143 extern int
144 extern void
                                     print hist(FILE*);
                                                                                                  /* ./hist.c
                                    get_hnum();
history( char*, int);
147
148 #ifdef DEBUG
                      static int Lev = -1; define TRACE(p) printf("%*s{ entering %s\n" ,++Lev * 4, "", p) define END_TRACE(p) printf("%*s} exiting %s\n" ,Lev-- * 4, "", p) define DIAG(f,a) printf(f,a)
150 #
151 #
152 #
153 #else
                      define TRACE(p)
define END TRACE(p)
define DIAG(f,a)
156
157 #endif
160 #ifdef STR CMDS
161 # define PSTR(subr,str) printf("%s <%s>\n", subr, str);
162 #else
163 # define PSTR(subr,str)
167
                                            "1.0" /* Version number
/* Maximum line number permitted by DOS
171 #define MAXLINE (2048+1) /* Largest input command line in bytes +1
172 #define MAXDIR 128 /* Largest number of objects on cmd line
173 #define CNTL Z ('Z'-'')
174 #define COMMENT '#' /* Deliniates Comment
168 #define VER
175 #define ISQUOTE(c)
177 #define ISWHITE(c)
178 #define SKIPWHITE(p)
179 #define ISVAR(c)
                                                     ((c)-'"' | (c)--'\'')
((c)--' ' | (c)--'\t')
while(ISWHITE(*p)){p++;}
((c)--'$' | (c)--'$')
                                                     /* Possible modes in which shell can operate
                                                     /* Get input from a file
/* Get input interactively from stdin
/* Get input from the command line
182 #define FILEMODE 0
183 #define INTERACTIVE 1
184 #define COMMAND
186 #define PMODE() ( Mode—COMMAND ? "COMMAND" : \
(Mode—FILEMODE ? "FILE" : "INTERACTIVE") )
189 /*
190
                    Token definitions for built-in commands.
191
192
193 typedef enum
195
                       ALIAS.
196
197
                       CD,
CMD,
198
                       EXIT.
                       HISTORY,
LOGOUT,
199
200
201
                       PWD,
202
                       REM,
                       SET,
                       SETENV.
204
205
                       SHIFT,
UNALIAS,
206
207
                       UNSET
208 } TOKEN;
210 /*--
211
                  Global variables:
213
214 static char** Numv ; /* Vector array for expanding $<num> vars
215 static int Numc = 0; /* count of valid entries in the above
216 static char Ibuf[MAXLINE]; /* Input buffer
```

(Continued on next page)



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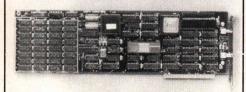
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#### C CHEST LISTING

```
(Listing Continued, text begins on page 18)
            218 static int
                                                              = 1; /* Generate CMDLINE env with spawned proc
='-'; /* Designates command line switches
                                                             - 1 ; /* Generate CMDLINE env with spawned proc
-'-'; /* Designates command line switches
- 0; /* Print input lines as they're read -v
- 0; /* Print commands as they're executed -x
- 0; /* Strip " or ' from quoted args -q
- 1; /* Nesting level of the current shell
- 0; /* Full path name of file specified in
/* Filemode input.
            219 static int
                                         Switchar
                                         Verbose
            220 static int
                                         Echo
            222 static int Noquotes
223 static int Shlev
224 static char *Filename
                    * Set up input mode for file mode processing. Note that Last_posn and * Ateof are used by file_input().
            228
            230
            231
            232 char
233 static int
                                                *file input();
                                               Mode
*(*Ifunct)()
Last_posn
Ateof
                                                                             - FILEMODE
                                                                             - file_input;
- OL;
- 0;
            234 static char
235 static long
            236 static int
            237
            238 reset fileinput()
            239 1
            240
241
                                 Last_posn = 0L
Ateof = 0
Mode = FILEMODE
            242
           243 244 }
                                                   - file_input
                                  Ifunct
            246 /*-
247 *
                                 Making isdigit into a subroutine makes processing marginally easier in exp_vars (below).
            248
249
            250
            251 digit(c)
                                 return( '0'<= c && c <= '9' );
            253
           254 };
            256 /
                                                                                                         Gets input from keyboard
Gets input from cmd line
                    * Input functions:
                                                              interactive_input()
                                                              command input()
            258
                                                              file input ()
                                                                                                         Gets input from a file.
                                 Only one of the three will be used depending on the way that the shell was invoked. All three return 0 on end of input, a pointer to the beginning of the current input line on success. The input line will have been loaded into the global array Ibuf, which is assumed to be dimensioned to MAXLINE characters.
            261
            264
           265 *
266 */
267 /*--
                   char
                                 *interactive_input()
            270
                                 register char *rval - NULL;
                                 TRACE ("interactive input");
            273
                                 if ( efgets (Ibuf, MAXLINE, stdin) )
rval = Ibuf;
            276
                                 END TRACE("interactive_input");
            279
            280
281 }
                                  return rval;
            282
            284
                                  *command_input()
                                 static int have been_called = 0;
register char *rval = NULL;
            287
            288
            290
                                 TRACE ( "command input" );
                                  if (!have been called )
            293
                                                unargv( Numc, Numv, Ibuf, MAXLINE, ' ');
rval = Ibuf;
            294
295
            296
                                                have_been_called++;
            298
            299
                                  END_TRACE( "command_input" );
                                 return rval;
            301 }
            304
            305
                  char
                                  *file_input()
                                                Get input for file mode. This kludge is required because a child process will close any open files when it exits. Consequently we open the file, get a line, remember the position within the file, and then close the file on each call. On the next call we'll return to the position we remembered in the previous call.
            307
            311
            312
            314
                                  register char *rval register FILE *fp;
            315
                                                                                - NULL :
                                  TRACE( "file_input" );
            318
                                  if (!Ateof)
             321
                                               if( !(fp = fopen(Filename, "r")) )
    fprintf(stderr, "Sh: Can't open batch file <%s>\n",
            322
                                                                                                                        Filename);
```

```
else if( !fseek(fp, Last_posn, 0) )
                                              Get a line from the buffer. Note that if the file doesn't have it's last line terminated with a carriage return, efgets will return true even though we're at end of file thus the call to feof.
 327
 330
 332
 333
                                          *Ibuf = 0;
                                         rval = efgets(Ibuf, MAXLINE, fp) ? Ibuf : NULL ;
if( !(Ateof = feof(fp)) )
Last posn = ftell( fp ) - 1;
fclose( fp );
 335
 338
339
340
341
342
343
344 }
                    END_TRACE( "file_input" );
return( rval );
345
346 /*
347
                    has_wild( bp )
       register char
                                  *bp;
350 {
351
352
                                  Return true if the string has a * or ? in it
                    353
354
355
356
357
358
359 }
                    return 0:
360
361 /*
362
                    strip( src ) *src;
363 int
364 reg
365 {
      register char
365
367
368
369
370
                                 Take care of special characters in a string (* and ?).
                                 Copy src onto itself, stripping out back-slashes. If a * or ? which isn't preceded by a backslash is found return 1, else return 0. If the first character in the string is a quote then * and ? aren't special.
371
373
374
375
                                             *dest = src;
special = 0;
376
                    TRACE ("strip");
378
                    while ( *src )
380
381
                                  if( *src == '\\' )
                                               /* Copy the char. following the \ into
 * dest and then, if we aren't at end of
 * string, advance src to point past the
 * escaped character
383
386
387
388
389
                                               if( *src )
                                                            ++src:
392
                                  else if ( ISQUOTE (*src) )
394
                                               /* Copy a quoted string verbatum, removing
    the quotes if Noquotes (-q) was given on
    the command line.
395
397
398
400
                                               if ( Noquotes )
401
403
                                               while ( *src && !ISQUOTE(*src) )
404
                                                            if( src[0] == '\\' && src[1] )
    *dest++ = *src++;
406
408
                                                             *dest++ = *src++ ;
409
411
                                               if( *src ) /* Then *src is a quote */
                                                             if ( Noquotes )
414
415
416
                                                                           *dest++ = *src++ :
417
418
                                  else
420
                                               /* Just do the copy. Set special to true
* if we copy a special character.
421
423
424
425
426
                                               if( *src == '*' || *src == '?' )
special = 1;
427
428
                                               *dest++ = *src++;
430
431
                    *dest = '\0';
433
                    END TRACE ("strip");
```

(Continued on next page)

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#### C CHEST LISTING

#### (Listing Continued, text begins on page 18)

```
return ( special );
436
437 }
439 /*-
440
                         nextarg( lp )
442 char
                       **1p;
442 Cl
443 {
444
445
446
447
                                       Get the next, space delimited, argument from the string pointed to by *lp. Return a pointer to the argument and update *lp to point past it. Leading white space is
                                       skipped.
448
                       register char *start, *line = *lp;
450
451
                       TRACE("nextarg");
SKIPWHITE( line );
 453
454
                       if( !*line )
    start = (char *)0;
 455
 456
457
                       else
 460
                               line = skipto( ISQUOTE(*start) ? *start : ' ', line , '\\');
461
                               if ( ISQUOTE (*line) )
 463
                               if( *line )
    *line++ = '\0';
466
467
468
469
470
                               *lp = line;
471
472
                       END TRACE ("nextarg");
474
                       return start:
475 }
477 /*-
478
479 int
                        exp_dir( buf, maxcount )
*buf;
 480
        char
481
482
                                      Remake buf, expanding any wild card characters into their proper names. That is, if buf containg the string: "foo * bar" and the current directory contains the files A, B and C then on exit, buf will point at the string "foo a b c bar". Expanded entries are sorted. Return 0 if a wild card that couldn't be expanded was found, else return 1.
483
484
485
486
487
 489
490
491
                       register DIRECTORY *dp = 0;
register char *arg, *sbuf, *p;
int i, rval = 1;
 492
493
494
495
                       TRACE("exp_dir");
 496
497
                       498
 499
500
                        dp->files = 1;
dp->dirs = 1;
dp->path = 1;
dp->sort = 1;
                                                                       /* Get all files
/* and all directories
/* and prepend the path name if given
/* the list should be sorted
 501
 502
503
 504
                                        Get the arguemts from the input src, one at a time, putting them into the dirv array of a DIRECTORY structure (one argument per dirv entry).
 505
 507
 508
                                       If the argument has special characters (* or ? not preceded by a \setminus or enclosed by quotes) then expand to a directory using dir(), otherwise just put the argument into the dirv array directly.
 510
 511
                                        Buf will grow larger if anything is expanded but it won't get larger than maxcount.
 514
515
 516
517
518
                         for ( sbuf - buf; (arg - nextarg(&buf)) && dp->maxdirs > 0; )
 519
 520
                                        if ( strip(arg) )
                                                        i = dp->maxdirs;
  522
  523
                                                        dir( arg, dp);
  525
                                                        if ( dp->maxdirs -- i )
  526
527
                                                                /* dir() didn't do anything
  530
                                                                fprintf(stderr, "Sh: Can't expand <%s>\n", arg);
  531
                                                                goto abort;
  534
  535
536
                                                        /* Add a command to dirv. First malloc space
* for it. Then copy it the arg into the
* malloc()ed space & put it into dirv.
* We use malloc in order to make it easier
* to free the space used by the DIRECTORY
  537
  538
```

```
* structure.
 542
543
544
545
546
547
                                             if(!(p = malloc(strlen(arg)+1)))
                                                         fprintf(stderr, "Sh: out of memory !!\n");
 548
549
550
                                             strcpy( p, arg );
                                             *( dp->lastdir )++ = p ;
--( dp->maxdirs );
++( dp->nfiles );
 553
 554
555
556
 558
 559
                    i = unargv(dp->nfiles+dp->ndirs, (char **)dp->dirv,sbuf,maxcount,' ');
                    561
 562
 564 abort: if ( dp )
 565
                                del_dir( dp );
                   END_TRACE("exp_dir");
return( rval );
 567
 568
569 }
 570
 573 char
                    *search( fname, ext )
*fname, *ext;
       char
 575 1
576
577
578
                    /* Search for fname.ext in the current PATH. Return a pointer * to the full path name if you find it (0 if you don't).
 579
                                            pathname[80], pbuf[129] ;
*p;
*paths ;
                    static char
register char
 581
 582
583
584
                                Assemble the pathname by concatanating fname and ext
 585
586
                   TRACE ("search"):
 587
                   if( strpbrk(fname, ".") )
     ext = "";
                                                                     /* If file name already has an */
/* extension don't add another. */
 590
 591
592
 593
                   sprintf( pathname, "%0.32s.%0.3s", fname, ext);
594
595
596
                   if (access (pathname, 04) < 0 )
597
598
                                            The file doesn't exist in the current directory.
                                           If fname contains the characters \ or / or if
the PATH enviornment isn't set, return a NULL,
else search for it along the path. strpbrk() is
a microsoft and Lattice library function. It'll
return true if fname contains a / or a \.
599
600
602
603
605
                               if( strpbrk(fname,"\\") || !(p = getenv("PATH")) )
     *pathname = '\0';
607
                               else
608
609
                                           strncpy( (paths = pbuf), p, 129);
610
                                           while ( p = next( &paths, ';', -1 ))
612
613
                                                  sprintf(pathname, "%0.50s\\%0.20s.%0.3s",
                                                                                p, fname, ext);
615
616
617
                                                  if( access( pathname, 04 ) >= 0 )
                                                        break;
618
                                                  else
619
620
                                                        *pathname = '\0';
621
                               }
622
623
624
625
626
                   END_TRACE("search");
                  return ( *pathname ? pathname : NULL ):
627 }
628
629 /*-
630
631 int
                   execute ( buf )
632 char
                   *buf:
632 CI
633 {
634
635
636
637
                               Execute a command. Return the programs exit status or -1 if the program didn't execute. This routine assumes that buf is at least DOSMAXLINE characters long.
638
                  register int
register char
639
640
                                           rval = 0;
*name;
641
642
                  static char
                                          envstr[MAXLINE+8] = "CMDLINE=";
643
644
645
                  TRACE("execute");
PSTR( "execute(1) buf=", buf );
                              Create the CMDLINE environment variable if Cmd it true (it can be set false with a "set cmd=0." If cmd is false then generate a "CMDLINE=" with no argument. This is necessary because MSDOS doesn't support a unset command.
649
                  if ( Cmd )
```

(Continued on next page)

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#### C CHEST LISTING

```
(Listing Continued, text begins on page 18)
                                                      strncpy( &envstr[8], buf, MAXLINE );
             654
                                                      envstr[8] - '\0';
             656
                                      setenv( envstr . 0 ):
             657
                                                     Truncate the command line itself (not the enviornment variable, at 127 characters and then echo it to the screen if Echo is set. Then extract the program name portion of the string (with the next() call).
             659
             660
             662
             663
                                     buf[DOSMAXLINE-3] - '\0';
             665
              666
             667
668
                                                      puts ( buf );
             669
                                      name = next(&buf, ' ', -1 );
              671
             672
673
                                                      Now try to spawn a new process. Suspend the shell until
              674
                                                      task returns.
              676
                                      if( (rval = spawnlp(P_WAIT, name, name, buf, NULL)) < 0 )
              677
                                                                      If we can't find a .com or .exe file then see if we can find a batch file with ole right name. Otherwise print an error message Note that we have to modify the CMDLINE enviornment string so that the leftmost argument (will be argv[0] in the child process) is the string "sh"
              679
              681
              682
              684
              685
              686
                                                       if( errno == ENOENT )
              688
                                                                       sprintf(envstr, "CMDLINE=sh %s %s", name, buf);
rval = spawnlp(P_WAIT, "sh", "sh", name, buf, NULL);
              690
              691
              693
              694
                                                                       printf("sh: Can't execute %s %s", name, buf);
perror(" ");
              696
              697
              699
              700
701
                                      END TRACE ("execute");
                                      return rval;
              702 }
              703 /
              705
              706 int
707 cha
                                       exp_vars( dest, src, maxcount, mode )
*dest, *src;
                      char
              708 {
                                                      Copy src into dest, expanding shell variables as appropriate. Shell variables all have a $ or a % as their first character. The global pointers Numv and Numc keep track of arguments for $<num> variables.
               710
              711
712
713
714
                                                       If mode — 1
If mode — 2
If mode — 3
                                                                                      aliases are expanded sargs are expanded both are expanded.
               715
              716
717
718
                                                      return the size of the expanded string. A mode 1 or 3 call will return with the high bit of src
               719
               721
               722
                                                                       num;
                                                                        *p;
*start_dest = dest;
               724
                                       register char
               725
726
                                       TRACE("exp_vars");
DIAG("exp_vars, input = <\s>\n", src );
               727
               728
729
               730
                                                      Expand aliases. First, remember the original start of the target array. Then, set the high bit of *src so that getvar will look for an alias. Then actually expand it, Then clear the high bit of the first character of dest (which will be set if no alias was found). There is a second-order recursion here in that getvar() makes a mode 2 exp_vars call.
               733
               736
737
738
               739
                                       if ( mode & 1 )
               741
742
743
                                                       *src |= 0x80;
getvar( &src, &dest, &maxcount );
DIAG("exp_vars, expanded aliases <\s>\n", start_dest);
                747
                                                       Now, expand shell variables. If we see and escaped character copy both the \ and the character to dest. else if the character isn't a shell variable (doesn't have a leading $ of $) just copy it. Else, try to expand the shell variable.
                750
                751
                754
                                         while ( *src && maxcount > 0 )
                                                        if( *src == '\\' && src[1] )
                                                                        /* If the character following the \ is a
```

```
* $ of % then strip the backslash and copy
* the $ or % to the dest array. Otherwise
* copy both the \ and the character that
* follows.
 760
 761
762
  763
  765
                                              src++:
                                                                        /* Skip past the \
  768
                                              if( !ISVAR(*src) )
 769
770
                                                           if( --maxcount <= 0 )
  771
                                                           break;
*dest++ = '\\';
  773
                                              if( --maxcount <= 0 )
 776
                                              break;
*dest++ = *src++;
 777 778 779
                                 else if( !ISVAR( *src ) || !(mode & 2) )
 780
781
                                              /* Either character is just a normal character
* (not a shell variable) or we're told to not
* expand shell variables.
 782
 783
784
 785
786
787
                                              *dest++ = *src++;
if( --maxcount <= 0 )
 788
789
790
                                 else if( digit( *++src ) )
 791
792
793
                                                          Expand a $<num> arg. first extract the <num> from source, then copy the correct vector out of the Numv array. If the Numv entry is NULL, don't put anything in the dest array.
 794
795
 796
 798
                                              for ( num = 0; digit(*src); )
    num = (num * 10) + (*src++ - '0');
 799
 801
 802
                                              if ( num < Numc )
                                                    for(p = Numv[num]; *p; *dest++ = *p++)
if(--maxcount <= 0)
 804
                                                                 break;
 806
 807
                                 else
 808
 809
                                              /* We've found a $ not followed by a number */
                                              switch( *src++ )
 812
                                             case '*':
 814
                                                           num = unargv(Numc, Numv, dest, maxcount, ' ');
                                                                             num;
                                                           dest += num;
maxcount -= num;
 817
                                             case 'p':
                                                          ;;
getcwd( dest, maxcount );
for(;*dest; ++dest, --maxcount )
    if( *dest =- '\\')
        *dest = '/';
 820
821
823
824
825
                                                          break;
                                             case '!': num = get hnum(); goto skip;
case 's': num = Shlev;
skip: if( maxcount > 5 )
826
827
828
829
830
831
                                                                       sprintf( dest, "%d", num );
for(;*dest; ++dest, --maxcount)
832
833
834
                                                             break:
835
                                             default:
837
                                                           --src
838
839
                                                          getvar( &src, &dest, &maxcount );
                                                          break:
840
843
                                *dest = '\0';
844
845
846
847
                   *start dest &= 0x7f ;
                   DIAG("exp_vars: on return <%s>,", start_dest );
DIAG(" returning %d\n", dest = start_dest );
848
849
850
851
                   END_TRACE("exp_vars");
                   return ( dest - start_dest );
854 }
856 /
857
      prompt()
859
                   /* Print a prompt using the PROMPT enviornment variable
 * Return true.
860
861
862
863
                                            buf[50];
*p;
865
                   register char
866
                   if ( Mode -- INTERACTIVE || Echo )
867
                               if( !(p = getenv("PROMPT")) )
  printf( "[%d:%d] ", Shlev, get_hnum() );
870
```

(Continued on next page)

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#### C CHEST LISTING

(Listing Continued, text begins on page 18) else 873 874 exp\_vars( buf, p, 50, 2 );
printf( buf ); 875 878 return 1: 880 883 int docmd ( cmd ) 885 { Do one command from line. 886 Return an exit status (or -1 on error). 888 889 890 register DIRECTORY \*dp; register int rval = -1; 891 TRACE("docmd");
PSTR( "docmd(1) cmd=", cmd ); 893 /\* If ther're no wild cards in cmd then just strip backslashes
\* and quotes. Else, expand the wild cards ( exp\_dir will strip
\* backslashes etc.) Then execute the command. 896 897 898 899 if(!has\_wild(cmd))
 strip(cmd); 901 /\* Just strip backslashes \*/ 902 else if( !exp\_dir(cmd, MAXLINE) )
 goto abort; 904 905 906 rval = execute(cmd) ; 907 909 abort: END\_TRACE("docmd"); 910 return ( rval ); 912 913 /\*----\*/ 915 rcopy ( dest, src, maxcount ) 916 register char \*dest; 917 char \*src; 917 char 918 { register char \*tbuf; \*tail; \*sd = dest; 921 922 DIAG("rcopy (top of proc): src = <%s>\n", src ); if ( !\*src || maxcount <= 0 ) 925 return: 927 if( !(tbuf = malloc( maxcount )) )
 fprintf(stderr, "Sh: out of memory\n"); 930 else 931 /\* 1) expand the src buffer into tbuf
\* 2) advance tail to point past the next; and replace
\* the; with a null. 932 934 935 \* 3) copy everything up to the tail to dest and add a \* ; to dest if the tail is non-null. \* 4) repeat this process using the tail as source. 936 936 937 938 939 940 941 exp\_vars( tbuf, src, maxcount , 3 ); /\*1\*/ 941 942 943 944 945 946 947 tail = tbuf;
next( &tail, ';' , '\\');
SKIPWHITE( tail ); /\*2\*/ for(src = tbuf; \*src && src-tbuf < maxcount; \*dest++ = \*src++)</pre> 948 if( (maxcount -= (src-tbuf)) && \*tail )
 \*dest++ = ';'; 949 951 952 \*dest - '\0': 953 DIAG("rcopy (before recursive call): dest = <%s>\n", sd ); rcopy(dest, tail, maxcount); /\*4\*/ 956 958 DIAG("rcopy (after recursive call): dest = <%s>\n", sd ); 959 962 963 char \*next\_cmd() 966 { Get a line from input, split off one command, and then return a pointer to it (or NULL if at end of input. History processing is done here too. Semicolons inside quoted strings or preceded by a '\'do not seperate commands. Comments and blank lines are absorbed (ie. next cmd won't return until it gets a real input line). 969 973 static char Cmdbuf[MAXLINE]; /\* Should initialize to zeros \*/

```
*src = Cmdbuf ;
*p;
*tbuf;
977
978
                      static char
register char
979
                      TRACE ("next cmd");
981
982
                      DIAG("next_cmd: src is <%s>\n", src );
                      while( !*src )
                                    do
                                                   prompt();
989
990
                                                   if( !(*Ifunct)() )
992
                                                                 END_TRACE("next_cmd" );
993
                                                                  return NULL;
995
                                                   }
996
997
                                                   DIAG("next_cmd: got <%s> from input\n", Ibuf);
998
                                                   src = Ibuf;
SKIPWHITE( src );
1000
1001
 1002
                                       while( !*src || *src -- COMMENT );
1003
                                      history( src, MAXLINE - (src-Ibuf) );
if( *src )
1004
 1005
 1006
1007
                                                     rcopy( Ibuf, src, MAXLINE );
src = Ibuf;
1008
1010
1011
                        p = next( &src, ';', '\\');
1013
                       DIAG("next_cmd: buffer <\s>\n", p );
DIAG("next_cmd: on next call buffer will be <\s>\n", src );
 1014
 1015
1016
1017
                            rcpy( Cmdbuf, p );
                        if( Verbose )
    printf("[sh %d input] <%s>\n", Shlev, Cmdbuf );
1019
 1020
                       DIAG("next cmd: returning <%s>\n", Cmdbuf);
END_TRACE("next_cmd");
1021
 1022
1023
1024
1025
                        return ( Cmdbuf );
1026
1027 /*-----
1029 char
                       *errmsgs[] =
1030
1031
                                       shell entered in interactive mode\n",
enter interactive mode\n",
commands are read from string. If several strings\n",
are present they are concatanated together before\n",
execution. The 'c' may be upper or lower case.\n",
commands are taken from <file>. Args are expanded\n",
to correspond with $0 $1 etc inside the file.\n",
for DOS compatability $0 $1 etc are also recognized\n",
Strip quotes from quoted argument strings. Usually\n",
they are left in so that a spawned process can\n",
exsemble its argv correctly.\n",
Print input lines to the shell as they are read.\n",
Print lines as they are executed\n",
         "sh
         "sh %ci
1032
1033
         "sh %cc <string>
1035
1036 "sh file args...
 1038
1039 "sh %cq
1040 "
1041 "
1041 "sh %cv
1043 "sh %cx
1044 ""
1045 };
1047 usage( c )
                                      Print the usage error message.
1049
                        /*
*/
1050
1052
                                                   **pp;
                       register char
1053
                        fprintf(stderr, "Sh: illegal argument <$c>: Usage is \n\n", c); for( pp = errmsqs; **pp ; fprintf(stderr, *pp++, Switchar) )
 1054
1055
1056
                        exit(1);
1058 }
1059
1060
1061 /
         int
1063
                                       setargs(p)
        register char
1064
1065
                                      Set various global flags base on command line arguments. This routine will also be used by "set" which can't recognize the -c switch. So, ignore -c if c enabled is false. p should point at the - on entry. Processing will terminate when a space or tab or end of string is found.
1066
1067
1069
1070
1071
1073
1074
                       TRACE ("setargs");
1075
1076
1077
1078
                       if( *++p )
                               for(; *p ; p++)
1079
1080
                                       switch(*p)
1082
                                       case 'c':
                                      case 'C':
case 'q':
case 'v':
case 'x':
 1083
                                                          Mode - COMMAND:
                                                                                                               break:
                                                          Mode = Cornel.
Noquotes = 1;
Verbose = 1;
Fcho = 1;
1085
1086
                                                                                                               break;
```

(Continued on next page)

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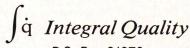
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#### C CHEST LISTING

```
(Listing Continued, text begins on page 18)
                                     case 'i': Mode
                                                                - INTERACTIVE;
        1088
                                     case
                                     case '\t'.
        1089
                                                                                           goto abort;
break;
        1091
        1092
                         }
        1094 abort:
        1095
                          END_TRACE("setargs");
        1096
        1097
        1098
        1100 doargs(argc, argv)
1101 char **argv;
        1102
                                                                    /* Check for comand line args & *,
        1103
                          if( -- argc <= 0 )
                                                                    /* use interactive mode if none */
/* are found. */
                                    Mode - INTERACTIVE;
        1105
        1106
                                     Ifunct = interactive input;
        1108
        1109
                                                                    /* There's at least one arg,
/* skip to second arg and call
/* set args if the first char
/* is a - Then skip past the
/* second arg too.
                          if( ** (++argv) - Switchar )
        1111
                                     setargs ( *argv++ );
                                      -argc;
        1114
        1115
1116
                         Numv = argv ;
Numc = argc ;
        1117
        1118
                          if( argc <= 0 || !**argv )
        1120
                                     fprintf(stderr,"Sh: missing file name, ");
fprintf(stderr,"use -i for interactive input\n");
        1121
        1123
                                    exit(1):
                          else if ( Mode -- COMMAND )
        1125
        1126
                                    Ifunct = command input;
        1128
                          else if( Mode == FILEMODE && !(Filename = search(*argv, "bat")))
        1129
        1130
                                     fprintf(stderr, "Sh: can't find <%s>\n", *argv );
        1131
        1132
                                     exit(1):
        1134 }
        1135
        1136 /
        1137
        1138 setenv ( env, allocate )
                          *env;
               char
        1140 {
                                  Set an enviornment. Env may be "name-contents" or the - may
        1141
                                  be a space.

If allocate is true allocate space, otherwise acutally use env (which must be static) as the string.
        1142
        1143
        1144
1145
        1146
                          register char
        1147
1148
                                    Look for either a space or a = in the string. If you find a space, replace it with an =.
        1149
        1150
1151
        1152
                          for ( p = env ; *p && *p != ' ' && *p != '=' ; p++ )
        1153
1154
                         if( *p = ' ' ' )
 *p = '=' ;
        1155
         1156
        1157
                                     Now set the enviornment to the indicated value
         1159
                          if( p = allocate ? strsave( env ) : env )
    if( putenv(p) != -1 )
         1161
         1162
         1164
                          fprintf(stderr,"Sh: setenv failed, out of memory\n"); if( allocate && p )
         1165
         1167
                                    free(p);
         1168 }
         1170 /
         1171
               register char *str;
         1174
                           1176
                            */
                           register int
                                                i;
*name;
         1180
                          char
         1181
         1182
1183
1184
                                     Get the name, replacing the trailing blank or - with a null. Print variables if there is no name.
                            */
         1185
                           if( !*str )
         1187
1188
                                     printf("%-8s: %s\n", "echo", Echo ? "ON" : "OFF");
printf("%-8s: %s\n", "cmd", Cmd ? "ON" : "OFF");
printf("%-8s: %s\n", "verbose", Verbose ? "ON" : "OFF");
         1189
         1191
                                     printvars();
```

```
1194
1195
                                for(name = str; *str; ++str)
if(*str == '=' || *str == ' ')
1196
1197
                                                       *str++ = 0;
1199
                                                       break:
1200
                                           Get the command tail, skipping past any =
1202
1203
                                           or blanks to get there.
1205
1206
                               while( *str == '=' || *str == ' ')
                                            ++str ;
1207
1209
                                           Process the command:
1211
                               i = *str ? atoi(str) : 1 :
1213
                                if(   !strcmp("echo", name))
else if( !strcmp("verbose", name))
else if( !strcmp("cmd", name))
else   setvar(name, str);
1214
1215
                                                                                       Echo
                                                                                       Verbose = i;
Cmd = i;
1216
1217
1218 1219 }
1220
                    Alias support uses the same tables as shell varialbes. However, the top bit of the first character of the alias name is set to indicate its function.
1222
1223
1225
1226
       unalias ( str )
1228 char
                    *str:
                    if( *str )
1230
1231
                                *str |= 0x80 ;
unsetvar( str );
1233
1234
1235 }
1236
1237 alias( str )
1238 char *str;
 1238 char
1239 {
1240
1241
1242
                    if( !*str
                                printalias();
                    else
1242
1243
1244
1245
                                *str |= 0x80 ;
set( str );
1246 1247 }
1248
1249 /
1251 void
                    pwd()
                                Print working directory
1253
1254
                                nbuf[80];
                    char
1256
1257
                    if( !getcwd( nbuf, 80 ) )
    fprintf(stderr, "sh: path too long to print.\n");
 1258
 1259
1260
                    printf("%s\n", nbuf );
 1262 }
1263
 1265
 1266 disk_present(id)
1268 {
                                Return true if there's a disk plugged into the indicated drive, else print an error message and return 1. This routine assumes that drive C has a disk in it so it will return true if id — 'c' or id —— 'C' without checking.
1269
1271
 1272
1274
1275
                    register int
register int
                                            try = 5;
err;
                                                                   /* times to try to read disk
 1277
 1278
                    union REGS
                                            regs:
                    if( (id = toupper(id)) -- 'C' || id -- 'c' )
1280
1281
                                return 1:
                                                           /* /*
/*
/*
                                                                    Service #4 (verify sec) */
                    regs.h.ah = 4:
 1283
                    regs.n.an = 4;
regs.h.al = 1;
regs.x.cx = 0;
regs.h.dh = 0;
regs.h.dl = id - 'A';
                                                                    # of sectors
track # & sector #
head #
 1286
 1287
                                                                    drive #
                                Actually read the disk. Loop until we've gotten a timeout error (0x80) from dos try times.
 1289
 1290
1291
                      */
 1292
                                 err = int86(0x13, &regs, &regs) & 0xff;
 1295
                    } while( (err & 0x80) && (--try >= 0) );
 1296
 1298
 1299
                     if( err )
  1300
                                                                    /* Recalibrate diskette system */
                           regs.h.ah = 0x0;
int86(0x13, &regs, &regs);
  1301
 1302
                           fprintf(stderr, "Cd: can't log on drive %c, ", toupper(id));
 1304
```

(Continued on next page)

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#### C CHEST LISTING

```
(Listing Continued, text begins on page 18)
                                fprintf(stderr, "DOS error code = 0x%02x\n" , err
         1306
         1307
                          return(!err);
         1309 }
         1310
1311 /*---
         1312
                                      cd ( name )
         1314 register char
                                      *name;
         1315 (
                                     Change the current directory to the indicated name. Log in a new disk if necessary. This routine is a bit more sophisticated than DOS itself, in that it checks if a disk exists before trying to log it on. This checking is done using the "Get diskette status" service of BIOS interrupt 0x13.
         1316
1317
         1318
         1319
         1320
         1321
                                      Get to the current directory on another disk by saying "cd x:" Get to another directory on another disk by saying "cd x:/dir/subdir/etc"
         1322
         1324
                            */
         1325
         1326
1327
                          if( *name && name[1] == ':')
         1328
                                      if (!disk_present(*name))
         1330
                                                 return:
         1331
                                      else
         1332
         1333
                                                 bdos ( 0xe, toupper (*name) - 'A', 0);
         1334
                          }
         1336
         1337
1338
                          1339
         1340 }
1341
         1342 /*-
         1344 shift()
         1345 {
         1346
1347
                                      Process the "shift" command (move all the $ args left
                                      one notch) .
                            */
         1348
         1350
                          if ( Numc )
         1351
                                      --Numc;
         1353
                                      ++Numv;
         1354
1355 }
1356
         1357 /*
         1359 doenv()
         1360
         1361
1362
                                      Reads and initializes the various enviornment variables.
                                      This routine should only be called once and it must be called before Shlev is used and before command line processing is done.
         1363
1364
         1365
                            */
         1366
1367
                          static char
                                                sbuf[16];
*p;
         1368
                          register char
         1369
1370
                          1371
         1372
                          1373
1374
                                                                    ( g* 33
                           sprintf(sbuf, "SHLEV=%d", ++Shlev);
         1376
1377
                           setenv (sbuf);
         1379
         1380 /*
         1382 use exit()
         1383 {
                                     We get here on a SIGINT (°C) interrupt
         1385
         1386
                          signal( SIGINT, use exit );
fprintf(stderr,"Use \"exit\" or \"logout\" to leave outer shell.\n");
         1388
         1389 }
1390
         1391 /*----
         1392
                          tokenize ( buf )
         1394 register char
1395 (
                                      *buf:
        1396
1397
                                     This is an extremely primitive token recognizer that will eventually be replaced with something more reasonable.
```

(Continued on page 98)

\*/

if( if( if( if( if( if( if(

if( !strcmp( "alias",
 if( !strcmp( "cd",
 if( !strcmp( "exit",
 if( !strcmp( "history

if (!strcmp( "unset",

1398 1399

1400 1401

buf)) return ALIAS;

buf)) return CD; buf)) return EXIT; buf)) return HISTORY;

buf)) return LOGOUT; buf)) return PWD; buf)) return REM;

buf)) return SETENV; buf)) return SET; buf)) return SHIFT;

buf)) return UNSET;

return UNALIAS;

buf))

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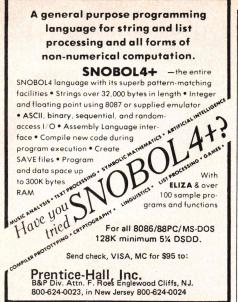
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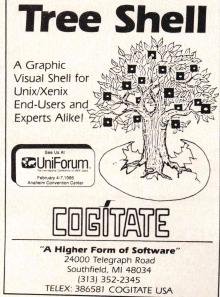
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#### C CHEST LISTING

```
(Listing Continued, text begins on page 18)
```

```
return CMD;
1412
1413 }
1414
1415 /
1416
1418 (
                              Process commands using the current Ifunct. Continue till end of input is reached. If expand vars is false then aliases and $args aren't expanded. This lets us read in the shrc.bat file unmolested
1419
1420
1421
1422
1423
1424
1425
                   char
                                           cmdbuf[9];
                                                                 /* Place to put extracted command */
                                          *p, *start; *cmd;
                   char
1427
                   register char
register int
                                          rval - 0:
1429
1430
                   while ( start = cmd = next cmd() )
1432
1433
                               DIAG("cmds: next cmd returned <%s>\n", cmd );
1435
1436
                                          Strip the command name from the rest of the command
1438
1439
                               p = cmdbuf;
for( i = 8; --i >= 0 && *cmd && *cmd != ' '; *p++ = *cmd++)
1441
1442
1443
1444
                               *p = '\0';
                               DIAG("cmds: partitioned <%s>", cmdbuf);
DIAG("+ <%s>\n", cmd);
1445
1447
                               SKIPWHITE (cmd) :
                               /* If Echo is set then echo the command to standard
 * output. Note that if i == CMD then the command
 * will be echoed in execute() (called by docmd()).
1449
1450
1452
1453
1454
1455
                               if ( (i = (int) tokenize(cmdbuf)) != (int) CMD && Echo )
                                          puts ( start );
1456
1457
                               switch(i)
1458
                                                      alias
1459
                               case ALIAS:
case CD:
                                                               ( cmd );
( cmd );
1460
                                                      cd (cmd);
rval = docmd(start);
                                                                                                     break;
                               case CMD:
                                                                                                    break:
                               case HISTORY:
1462
                                                      print_hist( stdout );
                                                      pwd
                                                                                                     break;
1464
                               case REM:
                                                                                                     break:
1465
                               case SET:
                                                                                                     break;
                                                      set (cmd);
setenv (cmd);
                               case SETENV:
                                                                                                    break:
                               case SHIFT:
case UNALIAS:
                                                                                                     break;
1467
                                                      shift ( );
unalias ( cmd );
                                                                                                     break;
                                                     unsetvar( cmd );
printf ("Illegal token\n");
                               case UNSET:
                                                                                                     break:
1470
                               default:
1471
1472
1473
                               case LOGOUT:
                                          reset fileinput();
if( access(Filename = "/logout.bat", 04) == 0 )
                                                      cmds();
1476
1477
                                           /* Fall through to EXIT case */
1479
                               case EXIT:
1480
                                          goto exit;
1482
1483 exit:
1484
                   return rval;
1485 }
1486
1488
1489 main(argc, argv)
1490 char **argv;
1491
1492
1493
1494
                              Exit status is that of the most
                              recently excuted program. In the case of a batch file
the exit status will be that of the shell processing
the batch file (ie. of the last program executed by
1495
1496
                              that shell).
1497
1498
                   1500
1502
1504
                              Process the two automatic batch file. /shrc.bat is executed every time a shell is created. /login.bat is only executed when a level 0 shell is created.
1506
1507
1508
1509
1510
                   if( access(Filename = "/shrc.bat", 04) == 0 )
1511
                              cmds()
1513
                              reset_fileinput();
1514
1515
1516
                  if ( Shlev == 0 && access(Filename = "/login.bat", 04) == 0 )
1518
                              cmds();
```

```
reset fileinput();
1520
1521
1522
                                        Now process the command line arguments. Doargs will
set the Echo, Verbose, Cmd and Mode variables as
appropriate. If we're in an interactive, level
0 shell, Call signal to prevent °C from working on
the shell itself and print a copyright notice.
1524
1527
1528
                         doargs (argc, argv);
                                                                                  /* Process command line args
1530
1531
                         if ( Shlev -- 0 && Mode -- INTERACTIVE )
1532
1533
1534
                                    signal( SIGINT, use exit );
fprintf(stderr,"SH (ver %s) - Copyright (c) 1985, ", VER);
fprintf(stderr,"Allen I Holub. All rights reserved.\n");
1535
1536
1537
1538
1539
                                  Finally, process commands from the input source determined by the command line.
1540
1541
1542
1543
                         exit( cmds() );
```

**End Listing** 

Shell listings continue next month

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#### PL/68K

#### LISTING ONE (Text begins on page 26)

```
Declare the format of a macro table node.
struct node {
             struct node * next; /* Pointer to next node. */
int nargs; /* Number of args in macro. */
char * name; /* Pointer to name of macro. */
char * text; /* Pointer to replacement text of macro. */
1:
             Define the hash table used to access nodes of the macro table.
#define MAC PRIME 101
struct node * ht [MAC PRIME];
             Look up a symbol in the macro table.
                                         if found, NZ if not found.
             Output: 2
                                                  pointer to text of symbol.

number of arguments (0-n), -1 if no arguments.
#define hash val d2
#define bp0 al
#define bp1 a2
lookup (symbol)
register char * symbol;
              register struct node **bp0, *bp1;
              register int hash_val;
              /* Get the hash value of the symbol into hash_val. */
             hash (symbol);
hash_val = d0;
              /* Point bp0 into the hash table. */
             bp0 = 6ht;
hash val *= sizeof(struct node *);
             adda (hash_val, bp0);
             /* Search down the list of buckets hanging from the hash table. */ for (bpl = *bp0; bp1; bpl = bpl -> next) {    str_eq(symbol, bpl -> name);    if \overline{(2)} {
                                         a0 = bpl -> text;
d0 = bpl -> nargs;
                                         move (Z BIT, ccr);
return;
             /* No match. */
move(NZ_BIT, ccr);
```

**End Listing One** 

#### LISTING TWO

```
ht:
         ds.1
lookup:
         movem.l
                           al/a2/d2, -(sp) ; function entry
         subq
                           #4, sp
         move.1
                           20 (sp), (sp) ; get the hash value into hash val
                           hash
d0, d2
         move.1
                           #ht, al ;point bpO into the hash table
#4, d2
d2, al
         move.1
         adda.1
         search down the list of buckets
                           (a1), a2; for (bp1 = * bp0;...;...)
         move.1
1:
         move.1
                           6(a2), (sp); str_eq(symbol, bpl -> name); 20(sp), -(sp)
         move.1
                           str_eq
                           #4, sp
         addq
                                                       ;if (2)
a0 = bp1 -> text;
d0 = bp1 -> nargs;
move(2_BIT, ccr);
        bnz
move.1
                           2
6(a2), a0;
4(a2), d0;
         move.w
         move
                                                                return;
_2:
                            (a2), a2; for(...;...; bp1 = bp1 -> next)
3:
         cmpa.1
                           #0, a2 ; for (...; bp; ...)
                           #0, ccr; move (NZ_BIT, ccr);
         move
4:
                           #4, sp ; function exit (sp)+, al/a2/d2
         addq
         movem.1
```

**End Listings** 

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#### MULTITASKING OS LISTING

```
(Text begins on page 44)
                                                 Terra Nova Communications multi-tasking kernel
                                                 Initialization and task-switcher
                                                 Note: this is not intended to be a complete listing. It's only a sample of some of the techniques used in our system.
                                                 PSECT
                                                                    Kernel
                                                 External symbols (defined in other code segments)
                                                                                                    ; vector table for hardware vector list
                                                                    VecTable,
                                                 EXTERN
                                                                                                    ; jump table for system calls
                                                                    JMPTable.
                                                                                                    ; length of ju mp table in longwords
                                                                     JMPTabLen.
                                                                     KernEnd,
                                                                                                    ;end of kernel code item in heap
                                                                     IOInit,
                                                                                                    ;our private I/O initialization routine;our private heap initialization;system variable initializer
                                                                    HeapInit,
SysInit,
                                                                     SysConMon,
                                                                                                     ;entry point for system console
                                                                                                    ;monitor task
;entry point for heap munger task
;entry point for disk munger task
                                                                    HeapMunger,
                                                                     DiskMunger
                                                 Entry points in this module (referenced from elsewhere)
                                                                                                    ;primary entry point to boot our OS
;main context switcher
;alternate context switcher (puts
                                                 ENTRY
                                                                    Start.
                                                                     ConSwitch,
                                                                     ConSwSleep
                                                                                                     ; calling task to sleen)
                                                 Include files (mostly equates)
                                                                     SysEqu
                                                                                                     ; contains the low-memory absolute
                                                 INCLUDE
                                                                                                     ;address equates (jump table, etc);defines the heap data structure
                                                 TNCT.IDE.
                                                                     HeapDef
                                                 INCLUDE
                                                                                                     ; contains hardware I/O equates
                                                                     SysIO
                                                 Miscellaneous storage
                       CodeHeap
                                                                                                     ; heap header for kernel heap item
                                                 DS.L
                                                                                                     ;system stack before tasking starts
                       StackEnd
                                                 DS.L
                                                                     40
                                                                                                     ;top of startup stack area
                        StackBegin
                                                 Pre-tasking initialization
                                                 this code works in single-task mode
                                                 prior to the invocation of the context switcher
                                                                     ;Initial entry. Calling operating system is still ;alive and kicking at this point.
                                                                     ReEntry, Al
                                                                                                     ;point to re-entry instruction
                        TakeOver
                                                  LEA
                                                                     A1,$20.W
                                                                                                     ; move short absolute to the vector
                                                  MOVE L
                                                                                                     ; for privilege exceptions
                                                                     USP, AO ;try a privileged instruction. If it; works, then we're in priv. mode. If not, then trap to
                                                 MOVE.
                                                                     ReEntry and be in privileged mode anyway
StackBegin, A7 ;set up initial stack
                        ReEntry
                                                  T.F.A
                                                                     StackBegin, A7
                                                  Turn off all interrupts in the system
                                                 Note: this is device-specific code.
The labels in the operand fields are from our own
                                                  SysIO include file.
                                                  CLR.B
                                                                     FDCIntMask
                                                                                                     ;clear floppy disk & system console ;clear hard disk completion int. mask
                                                  CLR.B
                                                                     HDTnt Mask
                                                                     SerIOlIntMask
                                                                                                     ; clear serial boards
                                                  CLR.B
                                                  CLR.B
                                                                     SerIO2IntMask
                                                  Initialize the vector table
                                                  Copy the vectors from an assembled table (in another module)
                                                  into the actual hardware vector list in low RAM
                                                  LEA
                                                                     VecTable, A0
                                                                                                    ; source (in another code segment)
                                                  LEA
                                                                     $0.W, A1
                                                                                                     ;destination (begins at $00 0000)
                                                                                                     ;192 longwords to move ;move a longword
                                                  MOVE
                                                                     #191,D7
                                                                     (A0) + (A1) +
                        VecMove
                                                  MOVE L
                                                                                                     ;repeat till done (fast loop on 68010)
                                                                     D7, VecMove
                                                 Copy system routine {\tt JMP} table from assembled object code (in another module) to low memory jump table, where everyone
                                                  can get at them.
                                                  T.F.A
                                                                                                     ; source
                                                                     JMPTable, A0
                                                                     ;Solice ;System,W,A1 ;the jump table. "System" is from the SysEqu include ;file. It's the context switcher) #JMPTablen/4,D7 ;number of longwords to move
                                                  T.F.A
                                                  MOVE.
                                                                                                     ;move a longword
;repeat till done (fast loop on 68010)
                        JPTMove
                                                  MOVE.L
                                                                     (A0) +, (A1) +
                                                  DBRA
                                                                     D7. JPTMove
                                                 Clear low memory to zero (between jmp table and kernel)

LEA StackEnd,Al ;point to top of destination

; and bottom of destination (end of the jump table)
                                                  LEA
                                                                     System+JMPTabLen.W, A0
                                                                     A0, A1
                                                  SIIRA
                                                                                                      ; calculate the length
                                                                                                     ;move to D7 for counting ;divide by 16 for 16-byte blocks
                                                  MOVE L
                                                                     A1,D7
#4,D7
                                                  LSR.L
                        LowClr
                                                  CLR. T.
                                                                     (A0) + (A0) +
                                                                                                     ; clear 16 bytes, quickly
                                                  CLR.L
                                                  CLR.L
                                                                      (A0) +
                                                  CLR.T.
                                                                     (A0) +
                                                  DBRA
                                                                     D7. LowClr
                                                                                                     :do it until done.
```

```
Clear high memory to zero (between kernel and end of RAM)
                                 (RAMEnd is first byte beyond RAM, defined in SysEqu)
                                                        RAMEnd,Al ;point to top of destination; and bottom of destination (end of the jump table)
                                LEA
                                                        KernEnd, AO
                                                                                                ;calc the length
;move to D7 for counting
;divide by 16 for 16-byte blocks
;clear 16 bytes, quickly
                                SUBA
                                                        A0, A1
                                                        A1,D7
#4,D7
                                MOVE.L
                                LSR.L
HiClr
                                CLR.L
                                                         (A0) +
                                CLR.L
                                                         (A0) +
                                CLR.L
                                                         (A0) +
                                CLR.L
                                                         (A0) +
                                DBRA
                                                        D7, HiClr
                                                                                                ;do it until done.
                                Initialize all of the primary I/O devices
Note: this is a device specific routine not treated in the article.
JSR IOInit
                                Initialize the heap Note: this is a routine in the heap manager, which creates
                                valid heap headers for the three initial heap items discussed
                                in the text: the deletion below the kernel, the kernel code item, and the deletion above the kernel.
                                                        HeapInit
                                Initialize the system zone of low memory Note: this sets up the TCB and master handle arrays, as discussed in the text, as well as initializing the time of day and the date and the order miscellaneous system values.
                                                        SysInit
                                Spawn off the initial tasks
                                This will create TCBs and TData items for the tasks, but won't
                                invoke them. They're invoked only by the context switcher.
                                                        SysConMon, A0
                                                                                                ;point to system console entry point;
tell it how much RAM for TData
;jump through jump table entry
                                LEA
                                                        #4096,D0
                                MOVE.L
                                                        Spawn ;jump through jump; ("Spawn" is a jump table equate in SysEqu)
HeapMunger, AO; spawn the heap munger
HeapMunger's That.
                                JSR
                                LEA
                                MOVE . L
                                                         #512,D0
                                                                                               ;heap munger's TData size
                                JSR
                                                        Spawn
                                LEA
                                                        DiskMunger, AO
                                                                                                ;Spawn the disk munger
;(TData includes one disk buffer)
                                MOVE L
                                                        #8192,D0
                                JSR
                                                        Spawn
                                LEA
                                                        TCB1.W, A2
                                                                                                ;get address of first TCB in array
;(TCB1 is defined in SysEqu)
                                BRA.S
                                                        ConSw1
                                                                                                 ; now start the context switcher!
                                Context Switcher: primary version Simple task-switch, nothing fancy.
                                SysFlags is a low-RAM system flag byte, defined in SysEqu.
                                Sysings is a low-kaw system flag byte, defined in SysEqu.

The data structure for the TData item is defined in SysEqu.

The data structure for the TCB is defined in SysEqu.

BTST #StopSys, SysFlags.W ;task switching inhibited?

BNE.S ConSwX ;yes, exit back to caller

MOVE.L OurTCB (A5), A0 ;get TCB address from TData

SUBA,L A5,SP ;subtract TData base addr from stack
ConSwitch
                                                        SP, TCBSP (A0)
                                                                                                ; save relative displacement in TCB
                                MOVE . L
                                MOVE.L
                                                        TCBNxt (A0), A2
TCBA5 (A2), A5
TCBSP (A2), SP
                                                                                                ;get address of next TCB
;get new TData base address
;get stack relative displacement
ConSw1
                                MOVE.L
                                MOVE.L
                                ADDA.L
                                                        A5, SP
                                                                                                 ;restore absolute address
ConSwX
                                RTS
                                                                                                 ;return to next task
                                Context Switcher: alternate very the calling task to sleep.

BTST #StopSys, SysFlags.W ;task switching inhibited?

ConSwX ;yes, exit back to caller without going to sleep
ConSwSleep
                                                                                                ;get TCB address from TData
;subtract TData base addr from stack
;save relative displacement in TCB
                                                        OurTCB(A5), A0
A5, SP
SP, TCBSP(A0)
                                MOVE. I
                                SUBA.L
                                MOVE.L
                                MOVE L
                                                        TCBNxt (A0), A2
                                                                                                ;get address of next TCB
                                                        TCBPrev(A0),A1; get addr of previous TCB
A2,TCBNxt(A1); close the pointers around the now-
A1,TCBPrev(A2); sleeping task.
                                MOVE.L
                                MOVE.L
                                MOVE . L
                                MOVE . B
                                                        #Sleep, TCBState (A0) ; mark it as asleep
                                MOVE.L
                                                        TCBA5 (A2), A5
                                                                                                ;get new TData base address
                                MOVE.L
                                                        TCBSP (A2), SP
                                                                                                ;get stack relative displacement
                                ADDA.L
                                                        A5, SP
                                                                                                ;restore absolute address
                                RTS
                                                                                                return to next task
```

**End Listing** 

#### 8080 SIMULATOR

#### LISTING ONE (Text begins on page 76) bne optprnt 8080 Simulator for MC68000 optprnt equ \* ifne trodsk ; If FCB tracing, print header. With CP/M 2.2 call support, optional tracing and lea.l fcbmsg,a0 Morrow HDDMA DMA buffer translating. bsr lpstr endc Version 1.2 1/21/85 JEC Fixed Extent bug in OPEN logic. : Execute simulation mloop: \* Sped up code, sample MAC from 2:13 to 1:40. Now runs at a 1.4 MHz equivalent based on MAC sample. mloop: ifne trace ; Optional trace. tst traceflg Version 1.1 8/29/84 JEC bne dotrace cmpa.1 tracesad, pseudopc Fixed BDOS call #6 bug. bne notrace Version 1.0 05/25/84 by Jim Cathey move.b #1.tracefle dotrace bsr dump This program has been written for speed wherever possible, as such tends to be large because of the separate subroutine for each and every opcode of the target processor. cmpa.l traceead, pseudopc bne notrace move.b #0.traceflg notrace equ \* On an 8MHz 68000 (Compupro) system the simulation speed is a little better than a 1MHz Z-80 when running MAC. The time for a sample assembly was 2:13 for the simulation vs 0:35 moveq #0.d0 ; Execute appropriate simulation subroutine on a 4MHz Z-80, both systems used identical hard disk systems. move.b (pseudopc)+,dØ Grab next opcode. as1 #2,d0 ; (DØ high word is still Ø!) move.1 0(opptr,d0.w),a0 It is not a complete simulation, as some flag handling isn't quite right, but it is enough to run the programs I wrote it for (DDT, LU, MAC, and Morrow's FORMATHW). jmp (aØ) : To the subroutine page text. page Illegal instructions and Dumping. \* This file contains the startup routines, the simulator core, tracing code, and the CP/M 2.2 simulation. ; Illegal opcode, say what & where, illegl move.l fillgmsg,d1 move.w #9,dØ bdos lea.l -1(pseudopc),a0 xdef optabl,flags,mnops globl mloop,illegl,service move.b (a0),d1 suba.l targbase,a0 bsr pbyte move.l #ilgmsg2,d1 Conditional assembly flags. move.w #9.dØ bdos equ Ø ; Non-zero for trace routine inclusion. ; Non-zero for FCB trace routine inclusion. move.l a0,d1 trcdsk equ Ø bsr pword move.l filgmsg3,d1 dmpdsk equ Ø ; Non-zero for register dump in FCB trace. move.w #9,dØ !! diskio is in file COM2.S !! \*diskio equ Ø ; Non-zero for special HDDMA support. bdos move.1 #dumpmsg,d1 move.w 19,dø bdos ; and spill guts. ; Quit simulation. Register definitions for the simulation. bsr dump rts equ @16,r return ; JMP (return) is fast return to MLOOP. ; 8080's PC is register A5. ; Pointer to opcode dispatch table. pseudopc equ @15,r equ @14,r movem.1 dØ-d1/aØ,-(sp) opptr dump ; Pointer to opeced elspatch table. ; 8080's SP is register A3. ; Pointer to 8080's flag lookup table is A2. ; Pointer to 8080's address space is A1. ; Base pointer to 8080's registers is A1. ; Dump all registers, ; used for illegals and tracing. move.1 /dmpmsg2,d1 pseudosp equ @13,r flagptr equ 012,r targbase equ 011,r move.w #9,d0 bdos equ @11,r regs move.b rega,d1 regconde equ 7,r ; Register based constant #\$E (for speed). bsr pbyte regcon01 equ 6,r : Register based constant #\$1. move.b regf,d1 regconof equ 5,r ; Register based constant #\$F. bsr pbyte bsr pspace Register based constant /\$FF. 8080's Flags 8080's Accumulator regconff equ 4,r regf equ 3.r move.w regb(regs),d1 rega equ 2,r bsr pword bsr pspace move.w regd(regs),d1 Note, only leaves DØ-D1/AØ for free use by entire bsr pword program without saving registers for temporary use. bsr pspace move.w regh(regs),d1 bsr pword .opd Ø,\$4e42 .opd Ø,\$4e43 ; BDOS 'macro'.; BIOS 'macro'. bdos bsr pspace bios move.l pseudosp,d1 sub.l targbase,d1 bsr pword bsr pspace move.1 pseudosp,aØ swap d2 : Save REGA Initialization and Main Opcode dispatcher. move.w #3,d2 tosloop move.b 1(a0),d1 \* ror w #8 d1 move.b Ø(aØ),d1 start lea.l target, targbase ; Start of target memory. ifne trace ; Optional trace code. bsr pword ; Enter trace delimiting addresses ; if the code is desired. bsr entrads bsr pspace addq.1 /2,a0 endc dbra d2, tosloop bsr lodfdos ; Load up the fake FDOS in target mem. swap d2 bsr lodregs ; Load the remaining simulation registers; ; Load the .COM program, ; quit if unsuccessful. move.l pseudopc,d1 sub.l targbase,d1 bsr loadcom

bsr pword

```
bsr pspace
       bsr pspace
       move.b (pseudopc).d1
       bsr pbyte
       bsr pspace
                               ; Now show mnemonic
       bsr pspace
       moveq #0,d0
       move.b (pseudopc),dØ
       asl.w #2,d0
       lea.l mnops, a0
       move.l (a0,d0.1),d1
       move.l d1,-(sp)
       inc.1 d1
       move #9,dØ
       move.l (sp)+,aØ cmp.b /" ",(aØ) beq nooprnd
       cmp.b #"C", (a0)
       bne notcons
       move.b 1(pseudopc),d1
       bsr pbyte
bra nooprnd
notcons cmp.b /"A",(a0)
       bne nooprnd
       move.b 2(pseudopc),d1
       move.b 1(pseudopc),d1
        bsr pbyte
nooprnd bsr pspace
                               ; In case of conout calls during trace,
        bsr pspace
                               ; they will be visible at end of line.
        bsr pspace
       movem.l (sp)+,dØ-d1/aØ
Initialization subroutines.
lodfdos lea.l fdos.a6
                               ; Load up the fake FDOS.
```

```
move.1 targbase, pseudosp
         adda.l #$10000, pseudosp
         lea.1 -256(pseudosp),a0
         move.w /fdoslen,d0
lodloop move.b (a6)+,(a0)+
dbra d0,lodloop
         lea.1 -256(pseudosp),a0
         move.l a0,d0
         sub.l targbase,dØ
move.b #$c3,0(targbase) ; Build BIOS & BDOS jumps.
         move.b #$c3,5(targbase)
         move.b d0,6(targbase)
         rol.w #8,d0
         move.b dØ,7(targbase)
         rol.w #8,d0
         add. w #3. dØ
         move.b d0,1(targbase)
         rol.w #8,dØ
         move.b dØ,2(targbase)
         move.w #0,-(pseudosp)
                                    ; Set up a return stack to exit simulation.
lodregs lea.l optabl,opptr
                                    ; Point base register to opcode dispatch
                                      table.
         lea.l mloop, return
         lea.1 flags,flagptr
move.1 targbase,pseudopc
         adda.l /$100,pseudopc
                                    ; Start execution at 0100H in target space
         moveq /$e,regcon@e
moveq /$1,regcon@1
                                    ; Set up quick constants.
         moveq #$f,regcon@f
         move.l /$ff,regconff
         moveq #0,rega
moveq #0,regf
entrads move.1 /tracemsg,d1
                                    ; Enter trace address if necessary.
         move.w #9,dØ
         bdos
                                               (Continued on next page)
```

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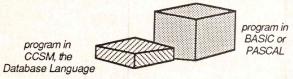
# 8080 SIMULATOR

### LISTING ONE (Listing Continued, text begins on page 76) ; Handle BIOS/BDOS service request ; Get trace start address. service moveq #0,d0 move.b (pseudopc)+,d0 of form HLT DB opcode. and.1 #\$ffff.d1 bne biosfn ; BDOS or BIOS? move.1 d1,a0 adda.l targbase,a0 moveq #0,d1 bdosfn move.b regc(regs),dØ : Get BDOS function number. move. 1 a0. tracesad move.w regd(regs),d1 cmp #31,d0 ; Get argument. move.1 #tracemg2,d1 ; Can't do Disk Parm Hdr function move.w #9,dØ beq badbdos bdos cmp #27.dØ ; or ALLOC vector fn. bsr atol : Get trace end address. bne okbdos and.1 #\$ffff,d1 badbdos move.l #ilgbdos,d1 move.1 d1,a0 move.w #9.dØ adda.l targbase,a0 bdos move.l ag.traceead : CRLF to end line. bsr dump move.w #10.d1 move.w #2,d0 bra quitprg bdos ; Translate target address to real address okbdos cmp #9,d0 move.w #13.d1 blt noconv move.w #2.dØ bdos beq noconv cmp #32,dØ beg noconv cmp #37,dØ OPEN file to be loaded, and load it into target beq noconv space if successful. add.1 targbase.d1 ; Save last known DMA address nocony cmp #26, dØ loadcom link a6,#0 : Mark stack frame. bne notdma ; (in case of OPEN processing). movem.1 d2-d3/a2-a4,-(sp) move.l d1.dmaaddr ; Get the address of the base page. move.1 12(a6),a0 notdma move.b #0, fcbflag ; Separate FCB type requests ; Get FCB address. lea.1 \$5c(a0).a2 cmp #15,dØ ; from the rest of the swine. move.b #'C',9(a2) ; mash filename to .COM blt notfcb : (Assume not, at first). move.b #'0',10(a2) move.b #'M',11(a2) cmp #24.dØ blt fcb move.l a2,d1 cmp #30,d0 move.w #15,dØ beq fcb cmp #33,dØ ; OPEN file. hdos cmpi.w #255,dØ · ERROR? blt notfcb beg openerr cmp #37,dØ ; Start loading at \$0100 in target. hlt. fcb move.1 pseudopc.d2 cmp #40,d0 filelod move.1 d2,d1 ; Set DMA address. beq fcb move.w #26.dØ bra notfcb bdos move.1 a2,d1 move.w #20,d0 : Read file until EOF. page fcb swap d2 bdos : Move the FCB to host working buf, move.w #35,d2 move.l d1,a0 tst dø bne basepg add.1 /128,d2 move.l a1,-(sp) lea.l fcbstor,a1 bra filelod fcb1 move.b (a0)+,(a1)+ dbra d2,fcb1 basepg lea.1 \$80(targbase),a2; Set up the target's base page. move.1 (sp)+,a1 ; Start with default DMA address. move.l a2,d1 ; and swap the random record bytes ; to make them match the 68000's. lea.l fcbstor,a0 move.1 a2, dmaaddr move.b 33(a0),d2 move.w #26,d0 move.b 35(a0),33(a0) bdos move.b d2,35(a0) lea.1 \$38(a0),a2 lea.1 \$5c(targbase),a3 ; Copy host's 2nd FCB to target's 1st FCB. swap d2 move.b #1, fcbflag ; Set flag for proper recovery. move.w #36,d0 move.1 d1,-(sp) ; (Gotta put the pig back in pen!) fcbloop move.b (a2)+,(a3)+ move.l a0,d1 dbra dØ,fcbloop ifne tredsk lea.1 \$80(a0),a2 ifne dmpdsk ; Optional^2 Register dump. lea.1 \$80(targbase),a4 bsr dump lea.l \$81(targbase),a3 endc clr dø endo move.b d0, (a4) ; Grab command tail from host's buffer. cmp.w #15.dØ : OPEN has a problem in that CP/M-68K move.b (a2)+,dØ ; can only open the base extent, unlike : Hack off ?. COM filename. bne notopen tail1 cmp.b #\$20,(a2)+ ; CP/M-80. So we have to check and do ; an OPEN then SEEK (RREAD) if required. dbeq dØ,tail1 bne loaded tst.b 12(a0) ; If there's any tail left, then beg notopen bsr openproc tail2 cmp.b #\$20,(a2)+ ; remove leading whitespace. dbne dØ tail2 beg loaded notopen: dec.1 a2 notopen: subq #2,dØ ifne trcdsk ; Optional FCB trace. tail3 move.b (a2)+.(a3)+; Move the rest of the tail. move.l d2,-(sp) move.b #' ',d2 inc.b (a4) dbra dØ.tail3 bsr fcbtrc1 move. b #0. (a3) move.1 (sp)+,d2 bra loaded endc openerr move.l fopnmsg,d1 : Can't open file. move.w #9,dØ notfcb cmp #6,d0 ; Not an FCB request. bdos clr dØ bne notdcon a direct console I/O function? cmp.b #\$ff.d1 ; Yes, make host's look like target's. bne notdcon loaded movem.1 (sp)+,d2-d3/a2-a4 move.w #\$fe,d1 unlk a6 : Trantor. bdos rts tst dØ beg results move.w #6.d0 move.w #SFF,d1 BIOS and BDOS service request handler. notdcon bdos ; FINALLY! Do the translated function. results move.w dØ, regh(regs) move.b dØ,rega

```
move.b regh(regs), regb(regs)
         tst.b fcbflag
                                     ; Do we need to restore a FCB?
         beq done
         ifne tredsk
         bsr fcbtrc2
         lea.l fcbstor.ag
                                      ; Restore the FCB to target, in proper
         move.b 33(a0),d2
move.b 35(a0),33(a0)
         move.b d2,35(a0)
         move.l (sp)+,a0 move.l a1,-(sp)
         lea.l fcbstor,a1
         move.w #35,d2
move.b (a1)+,(a0)+
fcb2
         dbra d2, fcb2
         swap d2
         move.l (sp)+,a1
move.b rega,d0
done
         and.w regconff,d0
         move.b Ø(flagptr,dØ.w),regf
openproc:
  openproc:
         ifne trodsk
                                      ; Optional FCB trace.
         swap d2
         move.b /' '.d2
         bsr fcbtrc1
         swap d2
         bsr fcbtrc2a
         endc
         move.b 33(a0),-(sp)
                                      : Save away RR fields!
         move.b 34(a0),-(sp)
move.b 35(a0),-(sp)
         movem.1 dØ-d2,-(sp)
         moveq #0,d2
move.b 12(a0),d2
                                      : Save desired extent.
         clr.b 12(a0)
```

```
ber febbdos
                                      ; Do BDOS (with opt. tracing).
         tst.b dØ
                                      ; No seek if not good OPEN.
; Make EXTENT / into record offset.
         bmi badopen
         asl.1 #7,d2
         moveq #0,d0
move.b 32(a0),d0
         bclr #7,dØ
         add.1 d2.d0
                                       ; Add onto CR to make abs record #.
         move.w d0,34(a0)
                                       : Put into FCB.
         swap dØ
          move.b d0,33(a0)
         move.l #junkbuf,d1
move.w #26,d0
                                       : Set DMA addr elsewhere for Rand Seek.
         bdos
         movem.1 (sp)+,dØ-d2
move.w #33.dØ
                                         Random READ (SEEK) desired extent.
                                         Do BDOS (with opt. tracing).
(OPEN) must always be successful because
          bsr fcbbdos
         clr dØ
                                         of the way CP/M-80 & CP/M-68K differ
                                       ; on OPENing non-zero extents.
; Restore the proper DMA address.
          movem.1 dØ-d1,-(sp)
         move.w #26,dØ
          move.l dmaaddr.d1
          bdos
          movem.l (sp)+,dØ-d1
restore move.b (sp)+,35(a0)
move.b (sp)+,34(a0)
                                       ; Restore RR fields.
          move.b (sp)+,33(a0)
badopen movem.1 (sp)+.dØ-d2
          bra restore
fcbbdos:
  fcbbdos:
          ifne tredsk
                                       ; BDOS call with optional FCB trace.
          move.l d2,-(sp)
move.b #'+',d2
          bsr fcbtrc1
                                                    (Continued on next page)
```

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# 8080 SIMULATOR

### LISTING ONE (Listing Continued, text begins on page 76) move.1 (sp)+,d2 movem.1 dØ-d1,-(sp) endc : Line termination if no result bdos bra fcbtr21 ; is to be presented. ifne trodsk bsr fcbtrc2 endc rts Misc. service routines. biosfn cmp #1,dØ : Handle Bios calls. (Inelegant, but rarely used so they stand as is). beq quitprg cmp #\$f,d0 ; List Status is ok. beq gudbios cmp #7,dØ : Don't allow disk functions! bge badbios pbyte move.1 #\$20018.d0 ; 2 nybbles, 24 bit shift first. gudbios clr.w d1 bra pdigits move.1 #\$40010,d0 move.b regc(regs),d1 pword ; 4 nybbles, 16 bit shift first. bra pdigits move.1 /\$60008,d0 movem. 1 d2-d7/a0-a6.-(sp) bios paddr ; 6 nybbles, 8 bit shift first. movem.1 (sp)+,d2-d7/a0-a6 bra pdigits move.1 #\$80000,d0 move.b dØ, rega ; 8 nybbles, no shift first. rts pdigits rol.1 d0,d1 ; Do shift. bra pdigent ; Flag illegal BIOS call badbios move.b dØ,-(sp) pdiglop swap do : Save nybble count. move.l /biosmsg,d1 move.w /9.d0 ; and spill guts. rol.1 #4,d1 ; Print variable in d1. bsr ntoa pdigent swap dØ : Get nybble count. move.b (sp)+,d1 dbra dØ,pdiglop bsr pbyte move.l /biosmg2,d1 rts move.w #9,dØ movem.1 dØ-d1,-(sp) ; Nybble in d1 to ASCII, then output. bdos and #\$f.d1 bsr dump cmp /sa,d1 blt ntoa2 add.b /'A'-'9'-1,d1 ; Trash return address and quitprg move.1 (sp)+,dØ add.b /'0',d1 ntoa2 rts : quit simulation. move.w #2,dØ bdos movem.1 (sp)+.dØ-d1 rts FCB Tracing support routines. pspace move.w #32,d1 ; Print a space. move.w #2,dØ bdos ifne trodsk rts fcbtrc1 movem.l dd-d2/ad,-(sp); Dump to printer each FCB usage move.b #9,d1; in format FN #, Disk, Name (ASCII) bsr lpchar; and the rest, all in hex but the move.w dd,d1; name field. Print the returned bsr lpbyte; value after the FCB. page Line Printer versions of above move.b d2,d1 ; Char in D2 is printed after FN /. bsr lpchar lpbyte move.1 #\$20018,d0 ; 2 nybbles, 24 bit shift first. bsr lpspace bra lpdigts lpword move.l #\$40010,d0 ; 4 nybbles, 16 bit shift first. bsr lpspace bra lpdigts lpaddr move.l #\$60008,d0 move.b (a0)+,d1 bsr lpbyte ; 6 nybbles, 8 bit shift first. bra lpdigts bsr lpspace move.w #10,d2 lplong move. 1 #\$80000, d0 ; 8 nybbles, no shift first. : Print Name field ... fcbtr1 move.b (a0)+,d1 lpdigts rol.1 dØ,d1 ; Do shift. bra lpdgent bsr lochar dbra d2, fcbtr1 lpdiglp swap dø ; Save nybble count. rol.1 /4,d1 bsr Intoa bsr lpspace ; Print variable in d1. move.w #3.d2 move.b (a0)+,d1 lpdgent swap dØ ; Get nybble count. ; Ex .. Rc dbra dØ, lpdiglp bsr lpbyte bsr lpspace dbra d2,fcbtr2 ; Nybble in d1 to ASCII, then output. Intoa movem.1 dØ-d1,-(sp) bsr lpspace and #\$f, d1 bsr lpspace lea.l 16(a0),a0 cmp /sa,d1 ; Skip dØ..dn field. blt Intoa2 move.w #3,d2 add.b /'A'-'9'-1,d1 fcbtr3 move.b (a0)+.d1 ; CR .. R2 add.b #'0',d1 1ntoa2 bsr lpbyte bsr lpspace Intoa3 move.w #5,dØ dbra d2.fcbtr3 bdos movem.1 (sp)+,dØ-d1 bsr lpspace bsr lpspace rts move.l dmaaddr.d1 sub.l targbase,d1 lpchar movem.l dØ-d1,-(sp) ; Print a character. bsr lpword bra Intoa3 bsr lpspace movem.1 (sp)+,dØ-d2/aØ lpspace movem.1 d0-d1,-(sp) move.w #32.d1 ; Print space. bra Intoa3 page fcbtrc2 movem.l dØ-d1,-(sp) ; Line termination of FCB trace. bsr lpspace bsr lpspace move.b d0,d1 Remaining misc. service routines. bsr lpbyte fcbtr21 move.b #10.d1 bsr lpchar move.b #13,d1 lpstr movem.1 dØ-d1,-(sp) ; Print a null-terminated string. lpstr1 move.b (a0)+,d1 bsr lpchar movem.1 (sp)+,dØ-d1 beq lpstr2 bsr lpchar bra lpstr1

```
lpstr2 movem.l (sp)+,dØ-d1
konin
       move.w #1,dØ
                      ; Console input for 'atol'.
atol
        moveq #0,d1
                        ; ASCII to long, stops on invalid hex char.
        clr d2
                        ; Returns long in d1, terminator char in d0,
atol1
       bsr konin
                        ; d2=1 if any chars entered before terminator.
        cmp.b #$40,d0
        blo atol2
       and /$5F,d0
cmpi.b /'0',d0
                        ; Mask to upper case.
                        ; Check range (0..9, A..F).
ato12
        blo atolend
        cmpi.b /'F', dØ
        bhi atolend
        cmpi.b #'9',dØ
        bls atol3
        cmpi.b #'A',dØ
        bhs atol3
        bra atolend
ato13
       moveq #1,d2
sub.b #'0'.d0
                        ; Valid characters entered, flag it.
        cmp.b #$9,d0
        bls atol4
        sub.b /'A'-'9'-1,dØ
       ext dØ
atol4
                       ; To long.
        ext.1 dØ
        asl.1 #4,d1
                        ; Tack it onto D1.
       add.1 dØ,d1
        bra atol1
atolend rts
              ***********
        Target processor's data registers.
        Fake FDOS.
```

7	even	A
regop3		; Operand 1 for DAA storage.
regb		; Offsets from register base pointer for
regc	equ -7	; 8080's pseudo-registers.
regd		; A & F are in Data Registers.
rege		; Pseudo-PC is kept in an Address Register.
regh		
regl	equ -3	
regop1		; Operand 1 for DAA storage.
regop2	equ -1	; " 2 " " "
fcbstor	ds.b 36	; Host works FCB's out of here.
fcbflag	ds.b 1	; Flag says we used the FCB buffer.
t =======	even d ds.l 1	
	d ds.1 1	; Trace start address.
		; Trace end address.
ti acei i	g ds.w 1	; Tracing enabled flag.
dmaaddr	ds.1 1	; DMA address storage.
	page	
fdos	dc.b \$76,0,\$C9	; Fake BDOS for target system.
*		; BIOS Jump Table
	dc.b \$C3,\$33,\$F	
		r : WDOOL
	dc.b \$C3,\$36,\$F	
	dc.b \$C3,\$36,\$F dc.b \$C3,\$39,\$F	F ; Const
		F ; Const F ; Conin
	dc.b \$C3,\$39,\$F	FF ; Const FF ; Conin FF ; Conout
	dc.b \$C3,\$39,\$F dc.b \$C3,\$3C,\$F	F ; Const F ; Conin F ; Conout F ; List
	dc.b \$C3,\$39,\$F dc.b \$C3,\$3C,\$F dc.b \$C3,\$3F,\$F	F ; Const F ; Conin F ; Conout F ; List F ; Punch
	dc.b \$C3,\$39,\$F dc.b \$C3,\$3C,\$F dc.b \$C3,\$3F,\$F dc.b \$C3,\$42,\$F	F ; Const F ; Conin F ; Conout F ; List F ; Punch F ; Reader
	dc.b \$C3,\$39,\$F dc.b \$C3,\$3C,\$F dc.b \$C3,\$3F,\$F dc.b \$C3,\$42,\$F dc.b \$C3,\$45,\$F	F ; Const F ; Conin F ; Conout F ; List F ; Punch F ; Reader F ; Home
	dc.b \$C3,\$39,\$F dc.b \$C3,\$3C,\$F dc.b \$C3,\$3F,\$F dc.b \$C3,\$42,\$F dc.b \$C3,\$45,\$F dc.b \$C3,\$48,\$F	F ; Const F ; Conin F ; Conout F ; List F ; Punch F ; Reader F ; Home F ; Seldsk
	dc.b \$C3,\$39,\$F dc.b \$C3,\$37,\$F dc.b \$C3,\$37,\$F dc.b \$C3,\$42,\$F dc.b \$C3,\$45,\$F dc.b \$C3,\$48,\$F dc.b \$C3,\$48,\$F	F ; Const F ; Conin F ; Conout F ; List F ; Punch F ; Reader F ; Home F ; Seldsk F ; Settrk
	dc.b \$C3,\$39,\$F dc.b \$C3,\$3C,\$F dc.b \$C3,\$3F,\$F dc.b \$C3,\$42,\$F dc.b \$C3,\$45,\$F dc.b \$C3,\$48,\$F dc.b \$C3,\$48,\$F dc.b \$C3,\$48,\$F	F ; Const F ; Conin F ; Conout F ; List F ; Punch F ; Reader F ; Home F ; Seldsk F ; Settrk F ; Setsec
	dc.b \$C3,\$39,\$F dc.b \$C3,\$3C,\$F dc.b \$C3,\$3F,\$F dc.b \$C3,\$42,\$F dc.b \$C3,\$48,\$F dc.b \$C3,\$48,\$F dc.b \$C3,\$48,\$F dc.b \$C3,\$4E,\$F dc.b \$C3,\$4E,\$F	F ; Const F ; Conin F ; Conout F ; List F ; Punch F ; Reader F ; Seldsk F ; Settrk F ; Setsec F ; Setdma
	dc.b \$C3, \$39, \$F dc.b \$C3, \$3C, \$F dc.b \$C3, \$3F, \$F dc.b \$C3, \$42, \$F dc.b \$C3, \$48, \$F dc.b \$C3, \$48, \$F dc.b \$C3, \$48, \$F dc.b \$C3, \$48, \$F dc.b \$C3, \$45, \$F dc.b \$C3, \$51, \$F	Const
	dc.b \$C3, \$39,\$F dc.b \$C3,\$57,\$F dc.b \$C3,\$57,\$F dc.b \$C3,\$42,\$F dc.b \$C3,\$45,\$F dc.b \$C3,\$48,\$F dc.b \$C3,\$48,\$F dc.b \$C3,\$54,\$F dc.b \$C3,\$51,\$F dc.b \$C3,\$51,\$F	Const

(Continued on next page)

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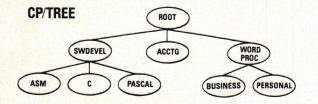
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# 8080 SIMULATOR

```
LISTING ONE (Listing Continued, text begins on page 76)
                                                                                regconff equ 4,r
                                                                                                           Register based constant #$FF.
       dc.b $76,1,$C9 ; Fake BIOS for target system
                                                                                                            Often used constants 10 & 18 are predominantly
       dc.b $76,2,$C9
                                  Const
                                                                                                            used by instructions that have 'quick' modes which encompass these values -- no register
                                   Conin
       dc.b $76,3,$C9
       dc.b $76,4,$C9
dc.b $76,5,$C9
                                  Conout
                                                                                                            needed (or available, either).
                                  List
                                                                                                           8080's Flags
                                                                                         equ 3.r
                                                                               regf
       dc.b $76,6,$C9
                                                                                         equ 2,r
                                                                                                           8080's Accumulator
                                                                                rega
        dc.b $76.7.$C9
                                   Reader
       dc. b $76.8.$C9
                                   Home *
                                                                                        egu -9
                                                                                                         : Operand 3 for DAA storage.
                                                                                regop3
        dc.b $76,9,$C9
                                   Seldsk *
                                                                                        equ -8
                                                                                                           Offsets from register base pointer for
        dc.b $76,10,$C9
                                   Settrk *
                                                                                regb
                                                                                rego
                                                                                        equ -7
                                                                                                           8080's pseudo-registers.
        dc.b $76.11.$C9
                                   Setsec *
                                                                                        equ -6
                                                                                regd
                                                                                                           A & F are in Data Registers
        dc.b $76.12.$C9
                                   Setdma *
                                                                                                           Pseudo-PC is kept in an Address Register.
                                                                                        equ -5
                                                                                rege
        dc.b $76,13,$C9
                                                                                        equ -4
                                                                                regh
        dc.b $76,14,$C9
                                   Write *
                                                                                        equ -3
                                                                                regl
        dc.b $76,15,$C9
                                   Listst
                                                                                        equ -2
                                                                                                         ; Operand 1 for DAA storage.
                                                                                regop1
        dc. b $76.16.$C9
                                 : Sectran *
                                                                                regop2
                                                                                        equ -1
                                                                                        data
                                                                                        page
page
                                                                                        dc.l nop00,lxib,staxb,inxb,inrb,dcrb,mvib,rlca
dc.l nop08,dadb,ldaxb,dcxb,inrc,dcrc,mvic,rrca
        Messages.
                                                                                         dc.l nop10,lxid,staxd,inxd,inrd,dcrd,mvid,ral
                                                                                        dc.l nop18,dadd,ldaxd,dcxd,inre,dcre,mvie,rar
dc.l nop20,lxih,shld,inxh,inrh,dcrh,mvih,daa
dc.l nop28,dadh,lhld,dcxh,inrl,dcrl,mvil,cma
illgmsg dc.b $d,$a,'Illegal instruction $'
ilgmsg2 dc.b ' at $'
ilgmsg3 dc.b '.$'
                                                                                         dc.1 nop30, lxis, sta, inxs, inrm, dcrm, mvim, stc
                                                                                         dc.1 nop38,dads,lda,dexs,inra,dera,mvia,cmc
                                                                                         dc.1 movebb, movebc, movebd, movebe, movebh, movebl, movebm, moveba
dumpmsg dc.b $d,$a,'Register contents:$'
                                                                                         dc.l movecb, movecc, movecd, movece, movech, movecl, movecm, moveca
dmpmsg2 dc.b $d,$a
                                                                                         dc.l movedb, movedc, movedd, movede, movedh, movedl, movedm, moveda
        dc.b '-AF- -BC- -DE- -HL- -SP- -SØ- -S1- -S2- -S3- -PC- -op-',$d,
                                                                                         dc.1 moveeb, moveec, moveed, moveee, moveeh, moveel, moveem, moveea
             $a,'$'
                                                                                         dc.1 movehb, movehc, movehd, movehe, movehh, movehl, movehm, moveha
biosmsg dc.b $d,$a,'Illegal BIOS call $'biosmg2 dc.b '.$'
                                                                                         dc.1 movelb, movelc, moveld, movele, movelh, movell, movelm, movela
                                                                                         dc.1 movemb, movemc, movemd, moveme, movemh, moveml, halt, movema dc.1 moveab, moveac, movead, moveae, moveah, moveal, moveam, moveae
blosmg2 dc.b 13,10,'Start trace at >$'
tracemg2 dc.b 13,10,'End trace at >$'
opmmsg dc.b 'Cannot open .COM file.$'
ilgbdos dc.b 'Unsupported BDOS call.$'
fcbmsg dc.b 9,'Fn# Dr NAME EX
                                                                                         dc.1 addb,addc,addd,adde,addh,addl,addm,addaa
                                                                                         dc.l adcb,adcc,adcd,adce,adch,adcl,adcm,adca
                                                                                         dc.1 subb, subc, subd, sube, subh, subl, subm, subaa
                                      EX S1 S2 RC CR RØ R1 R2 Addr Rslt
                                                                                         dc.1 sbbb, sbbc, sbbd, sbbe, sbbh, sbbl, sbbm, sbba
                                                                                         dc.l andb,andc,andd,ande,andh,andl,andm,anda
dc.l xrab,xrac,xrad,xrae,xrah,xral,xram,xraa
             ,10,13
         dc.b 9,'-
                                                                                         dc.l orab, orac, orad, orae, orah, oral, oram, oraa
             ,10,13,0
                                                                                         dc.l cmpb,cmpc,cmpd,cmpe,cmph,cmpl,cmpam,cmpaa
dc.l rnz,popb,jnz,jmpa,cnz,pushb,adi,rstØ
dc.l rz,ret,jz,nopCB,cz,call,aci,rst8
         bss
                                                                                         dc.l rnc.popd.jnc.out.cnc.pushd.sui.rst10
dc.l rc.nopD9,jc.in.cc.nopDD.sbi.rst18
dc.l rpo.poph.jpo.xthl.cpo.pushh.ani.rst20
dc.l rpe.pchl.jpe.xchg.cpe.preED.xri.rst28
*************************
        Target processor's address space.
                                                                                         dc.l rp,popp,jp,di,cp,pushp,oria,rst30
dc.l rm,sphl,jm,ei,cm,nopFD,cpi,rst38
; Actual storage for 8080's other registers.
registers ds.b 10 target ds.b $10000
                                ; 8080's universe.
junkbuf ds.b $80
                                 ; For BDOS' OPEN faking (RREAD buffer).
                                                                                         Flag register lookup tables.
         .end
                                                                                 End Listing One
                                                                                 flags dc.b $00,$01,$04,$05,$40,$41,$44,$45,$80,$81,$84,$85,$C0,$C1,$C4,$C5
                                                                                       LISTING TWO
                             *************
                                                                                       This file contains the target processor (8080) simulation
         routines.
                                                                                       Opcode dispatch table. One longword entry per opcode of the target (8080) processor, including illegals.
                                                                                        dc.b $80,$84,$84,$80,$84,$80,$80,$80,$84,$84,$80,$80,$84,$80,$84,$80
                                                                                        dc.b $80,$84,$84,$80,$84,$80,$80,$84,$84,$80,$80,$84,$80,$84,$80,$84,$80
                                                                                        dc.b $84,$80,$80,$84,$80,$84,$84,$80,$80,$84,$84,$80,$84,$80,$84,$80,$84
                                                                                     globl optabl, flags, nopØØ
          xdef mloop,illegl,service,preED,outspec
                                                                                          Opcode simulation subroutines.
                          ; Non-zero for special HDC/DMA support.
 diskio
                                                                                                 I/O instructions are based at 68000 address $FF0000
          equ @16,r
                          ; JMP (return) is fast return to MLOOP.
                                                                                                  as is appropriate for the CompuPro CPU-68K card.
 pseudopc equ @15,r
                            8080's PC is register A5.
                            Pointer to opcode dispatch table.
8080's SP is register AJ.
Pointer to 8080's flag lookup table is A2.
Pointer to 8080's address space is A1.
 opptr
          equ @14.r
                                                                                          Also, all routines assume that the high word of DØ = Ø!
 pseudosp equ @13,r
 flagptr equ 012,r
targbase equ 011,r
                                                                                             *************
          equ @11,r
                          ; Base pointer to 8080's registers is A1.
                                                                                          text
                                                                                          even
 regconde equ 7,r
regcond1 equ 6,r
regcondf equ 5,r
                          ; Register based constant #$E (for speed).
                          ; Register based constant /$1. Register based constant /$F.
                                                                                 nopøø jmp (return)
                                                                                                                                    ; 00 Nop
```

lxib	<pre>move.b (pseudopc)+,regc(regs) move.b (pseudopc)+,regb(regs) jmp (return)</pre>	; Ø1 Lxi BC,nnnn	dcxb	dec.w regb(regs) jmp (return)	; ØB Dcx B
			inrc	inc.b regc(regs)	; ØC Inr C
staxb	move.w regb(regs),dØ	; Ø2 Stax B		move sr,dØ	
	move.b rega,Ø(targbase,dØ.1) jmp (return)			and.w regconde,dd and.w regcond1,regf	
	Jmp (return)			or.b Ø(flagptr,dØ.w),regf	
inxb	inc.w regb(regs)	: 03 Inx B		jmp (return)	
LIIAU	jmp (return)	, by Thix B		Jmb (recorn)	
	Jup (1 codi ii)		dere	dec.b regc(regs)	: ØD Der C
inrb	inc.b regb(regs)	: Ø4 Inr B	dere	move sr.dØ	, bb bci o
	move sr.dØ			and.w regconde,d0	
	and.w regconde,d0			and.w regcon01,regf	
	and.w regcon01,regf			or.b Ø(flagptr.dØ.w),regf	
	or.b Ø(flagptr,dØ.w),regf			jmp (return)	
	jmp (return)				
			mvic	move.b (pseudopc)+,regc(regs)	; ØE Mvi C
dcrb	dec.b regb(regs)	; Ø5 Dcr B		jmp (return)	
	move sr,dØ				
	and.w regconde,d0		rrca	ror.b /1,rega	; ØF Rrc
	and.w regcon01,regf			bra docyf	
	or.b Ø(flagptr,dØ.w),regf				
	Jmp (recurit)		nop1Ø	bra illegl	; 10 Illegal for 8080
mvib	move.b (pseudopc)+,regb(regs)	; Ø6 Mvi b.nn			
	jmp (return)		lxid	move.b (pseudopc)+,rege(regs)	; 11 Lxi DE,nnnn
				move.b (pseudopc)+,regd(regs)	, II BAL DE, IIIIIII
rlca	rol.b /1,rega	; Ø7 R1c		jmp (return)	
docyf	bcs rlc1				
	bolr #0,regf				
	jmp (return)		staxd	move.w regd(regs),d0	; 12 Stax D
rlc1	bset #0,regf			move.b rega, Ø(targbase, dØ.1)	
	jmp (return)			jmp (return)	
nopØ8	bra illegl	; Ø8 Illegal for 8Ø8Ø			
nopue	bia illegi	; be illegal for enen	inxd	inc.w regd(regs)	
dadb	move.w regb(regs),dØ	: Ø9 Dad B	Ilixu	jmp (return)	; 13 Inx D
4445	add.w dØ,regh(regs)	, , , , , , , , , , , , , , , , , , , ,		Jmp (recurr)	
	bra docyf				
ldaxb	move.w regb(regs),d0	; ØA Ldax B			
	move.b Ø(targbase,dØ.1),rega				
	jmp (return)				(Continued on next page



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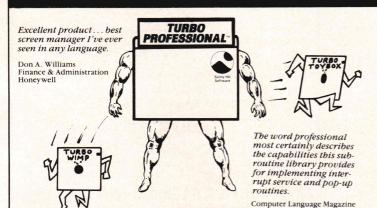
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# 8080 SIMULATOR

# LISTING TWO (Listing Continued, text begins on page 76)

nrd	inc.b regd(regs) move sr,d0	; 14 Inr D		move.b (pseudopc)+,regh(regs) jmp (return)	
	and.w regconde,d0				
	and.w regcon01,regf		shld	move.b 1(pseudopc),dØ	; 22 Shld addr
			Silla	rol.w #8,dØ	, as silla addi
	or.b Ø(flagptr,dØ.w),regf				
	jmp (return)			move.b (pseudopc),dØ	
				addq.1 /2,pseudopc	
crd	dec.b regd(regs)	: 15 Dcr D		move.l dØ,aØ	
	move sr, dØ			adda.l targbase,a0	
	and.w regcon@e,d@			move.b regl(regs),(a0)+	
	and.w regcon01,regf			move.b regh(regs),(a0)	
	or.b Ø(flagptr,d0.w),regf			jmp (return)	
	jmp (return)				
			inxh	inc.w regh(regs)	; 23 Inx H
vid		; 16 Mvi D,nn		jmp (return)	
	jmp (return)				
			inrh	inc.b regh(regs)	; 24 Inr H
al	roxr.b #1,regf	; 17 Ral		move sr.dØ	
	[HT NEW HOLD NEW HOL				
	roxl.b /1,rega			and.w regconde, do	
	roxl.b /1,regf			and.w regcon01,regf	
	jmp (return)			or.b Ø(flagptr,dØ.w),regf	
				jmp (return)	
op18	bra illegl	; 18 Illegal for 8080			
	•		dcrh	dec.b regh(regs)	: 25 Dcr H
add	move.w regd(regs),dØ	; 19 Dad D		move sr,dØ	
auu		, I) Dau D			
	add.w dØ,regh(regs)			and.w regconde,d0	
	bra docyf			and.w regcon01, regf	
				or.b Ø(flagptr,dØ.w),regf	
daxd	move.w regd(regs),d0	; 1A Ldax D		jmp (return)	
	move.b Ø(targbase,dØ.1),rega				
	jmp (return)		mvih	move.b (pseudopc)+,regh(regs)	; 26 Mvi H,nn
	Jap (1 cour 11)			jmp (return)	.,
land	dea u read/reas)	; 1B Dcx D		J-P (1 0001 11)	
lcxd	dec.w regd(regs)	, IB DCX D	daa	mous b regen 3/regs \ da	. 27 Das
	jmp (return)		daa	move.b regop3(regs),d0	; 27 Daa
				roxr.b dØ	
nre	inc.b rege(regs)	; 1C Inr E	A Part of the Part	move.b regop2(regs),d0	
	move sr,dØ			move.b regop1(regs),d1	
	and.w regconde,d0			swap regconde	
				move.b rega, regconde	
	and.w regcon01, regf			and.b regcon0f, regcon0e	
	or.b Ø(flagptr,dØ.w),regf				
	jmp (return)			cmp.b /9,regconde	
				bhi halfcy	
dere	dec.b rege(regs)	; 1D Der E	C Plants in	and.b regcon0f,d0	
	move sr.dØ		A CALL	and.b regconØf,d1	
				ori.b /\$fØ.d1	
	and.w regconde,d0			addx.b dØ.d1	
	and.w regcon01,regf				
	or.b Ø(flagptr,dØ.w),regf			bcc nohalf	
	jmp (return)		halfcy	add.b /6, rega	
				bcs fullcy	
nvie	move.b (pseudopc)+,rege(regs)	: 1E Mvi E.nn	nohalf	btst /0,regf	
The Party	jmp (return)			bnz fullcy	
	J. 1. 2001 11/		STATE OF THE STATE OF	move.b rega, regconde	
	name h 41 name	. IF Dan		and.b /\$fØ,regconØe	
rar	roxr.b /1, regf	; 1F Rar	A Share		
	roxr.b /1,rega			cmp.b /\$90,regconde	
	roxl.b /1,regf			bls nofull	
	jmp (return)				
nop2Ø	bra illegl	; 20 Illegal for 8080			
lxih				(7 1-41-4-4-1 5	ontinued next mont
	move.b (pseudopc)+,regl(regs)	; 21 Lxi H,nnnn		(Listings to be (	ontinued next mont

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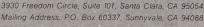
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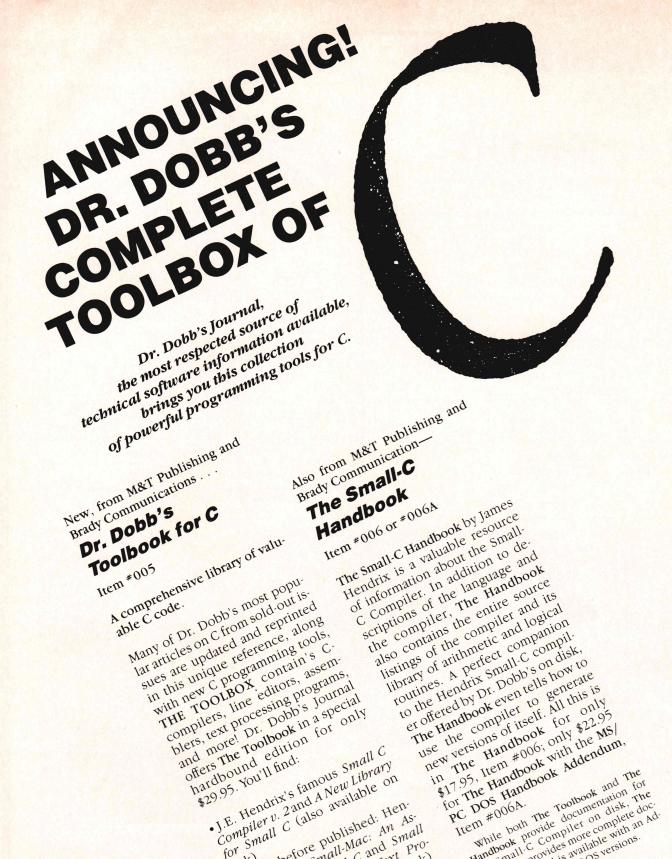
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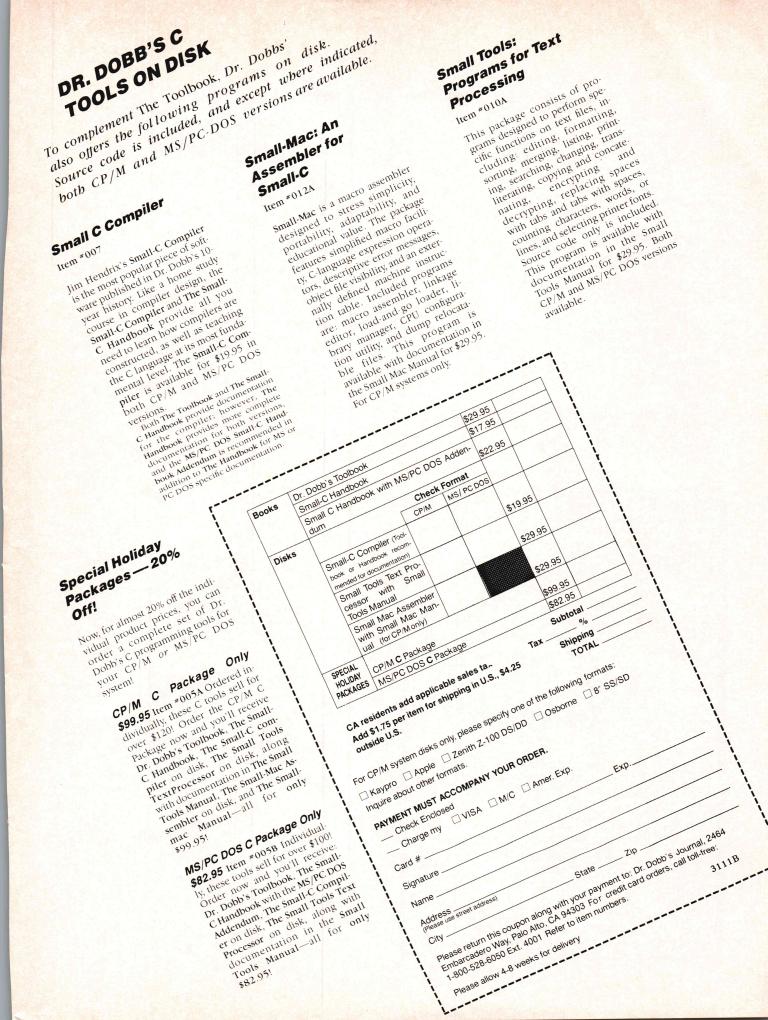
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# 16-BIT

# LISTING ONE (Text begins on page 118)

```
DUMMY SEGMENT TO BE GROUPED WITH HIMEM, WHICH
  IS AN EMPTY SEGMENT AT THE END OF THE USER'S
  PROGRAM
HIMEM GROUP HIDATA
HIDATA SEGMENT COMMON 'HIMEM'
FLARB DW 0
HIDATA ENDS
  'MAIN' IS THE PROGRAM'S ENTRY POINT,
  JUST ABOVE THE PSP
EXTRN MAIN: FAR
ARCODE SEGMENT PUBLIC 'CODE'
      ASSUME CS: ARCODE
      PUBLIC FREMEM
FREMEM PROC FAR
      PUSH BP
                                 STANDARD ENTRY STUFF
      MOV BP, SP
  GET THE ADDRESS OF HIMEM, THE TOP OF THE
  USER'S PROGRAM, STICK IT IN USER'S VARIABLE
                              ; GET TOP SEGMENT
; SEND IT TO THE USER
      MOV AX. SEG HIDATA
       LES
            BX, 18[BP]
            ES:[BX],AX
  FIND OUT HOW MUCH MEMORY IS AVAILABLE ABOVE
  HIMEM BY REQUESTING A LOT OF MEMORY.
THE CALL TO FUNCTION 4AH NEEDS THE SEGMENT
```

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```
ADDRESS OF THE PROGRAM'S PSP. GET IT BY
 SUBTRACTING 10 PARAGRAPHS FROM THE SEGMENT
 OF THE PROGRAM'S ENTRY POINT, WHICH HAS THE
 LABEL 'MAIN'.
           AX, SEG MAIN
                              ; GET ENTRY SEGMENT
                              ; ADJUST TO GET PSP
      SUB
           AX, 10H
                              ; PUT VALUE INTO ES
      MOV
           ES, AX
      MOV BX, -1D
MOV AX, 4A00H
                                 GET ALL OF MEMORY
                                 DOS FUNCTION CODE
                               ; DO IT
      TNT 21H
 THE NUMBER OF AVAILABLE PARAGRAPHS IS IN BX,
 RETURN IT TO USER.
          DI,14[BP]
                               ; STORE IT IN
      MOV ES:[DI], BX
                              ; USER'S VARIABLE
 TRY TO ALLOCATE THE PARAGRAPHS REQUESTED.
 A VALUE OF -1 MEANS DON'T ALLOCATE ANY.
                                 GET THE USER'S
      LES
           DI.10[BP]
           BX, ES: [DI]
      MOV
                                 VARIABLE
            BX,-1D
                                 SHOULD WE ALLOCATE?
      CMP
                                NO, BAIL OUT
GET ENTRY POINT
      JE
            OK
           AX, SEG MAIN
      MOV
                                 ADJUST TO GET PSP
PUT IT IN ES
DOS FUNCTION CODE
      SUB
            AX, 10H
      MOV
            ES, AX
      MOV
            AX, 4A00H
      INT
            21H
                                 DO IT
            ERR
                                 CHECK FOR ERRORS
       JNA
                                NONE, CLEAR FLAG
STORE ERROR CODE
OK:
      XOR
            AX, AX
            BX, 6[BP]
      LES
ERR:
                                 IN USER'S VARIABLE
            ES:[BX],AX
  THAT'S IT
                               ; STANDARD EXIT STUFF
       VOM
            SP, BP
       POP
       RET
           16
FREMEM ENDP
ARCODE ENDS
        END
                                                 End Listing One
```

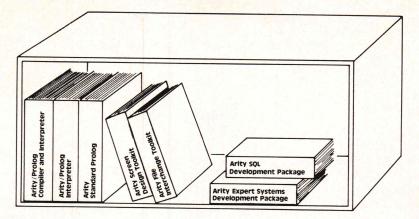
# LISTING TWO

```
; Michael Barr's 32-bit Square Root Routine
  Call with CX:BX = argument
            BX
                    = result
  Returns
                                         ;CX:BX = argument
sart
           proc
                       near
           push
                       ax
                                         ; save other registers
           push
                       dx
            push
                       di
            jcxz
                       sqrt3
                                         ; prepare first try
                       dx, cx
                       di,-1 dx,1
           mov
                                         ;estimating size of arg.
sart1:
            shl
                       sqrt2
                                         ;to guess initial try
            shl
                       dx, 1
            jc
                       sgrt2
            shr
                       di.1
                       sqrt1
            ami
sart2:
                        dx, cx
                                         ; restore argument
            mov
                       ax, bx
                       dx, di
                                         :prevent overflow
            amp
                        sart4
            div
                       di
                       ax, di
                                         ; comp quotient and divisor
            cmp
                        sqrt4
                       di,ax
di,1
                                         :average them
            add
            rcr
                                         ;and do it again ;prepare first try
                        sqrt2
            jmp
sgrt3:
            mov
                       dx, bx
di, Offh
            mov
                                         ; lower half zero?
                        bx, bx
            jnz
                        sart1
                                         ; no, jump
                       di, bx
            mov
sart4:
                        bx, di
                                         ;return result in BX
                        di
            pop
            pop
                        dx
            pop
sart
            endo
```

**End Listings** 

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# Trojan Horse Programs

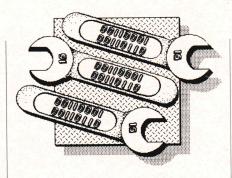
ne of the more ominous developments in the last few months is the appearance of so-called Trojan horse programs on some public bulletin board systems. These have appeared under such names as EGABTR, VDIR, and SYSUTIL. They are usually accompanied by scanty documentation that bills them as super disk directories or something similar, but their actual effect is to trash your system by formatting the hard disk, erasing the file allocation table, or writing random garbage into files. I put the people who create and upload such programs to public BBSs in the same category as terrorists.

Along similar lines, some villains have taken advantage of their knowledge of the public-domain PC remote BBS to upload programs that purport to paint a pretty picture on the screen or do something else cute but actually copy the password file or other vital system files to (unprotected) files under another name. The villain calls back later and downloads the unprotected files, thereby gaining access to privileged files and messages.

As the computer terrorists become more clever, we can expect the appearance of subtle sabotage programs that copy themselves to hidden files on the hard disk or attach themselves to the bootable operating system and don't do their damage until much later. For example, I can envision a hard disk destruction program that would wait until it saw that you had not run Backup for a week! The nature of the damage could also be so subtle that it would drive you crazy, such as simply changing a bit in a random sector of

# by Ray Duncan

The programs published in this month's column are available for downloading from the Laboratory Microsystems RBBS at (213) 306-3530 (300 baud or 1,200 bps).



the hard disk every few days.

These developments have the potential to damage or destroy the proliferation of public-domain software on bulletin boards, which would be very sad. BBS operators will not want to take the risk of being held liable for disasters due to Trojan horse programs. On the LMI RBBS, we are immediately adopting a policy of deleting all programs that are not uploaded in the form of, or at least accompanied by, source code as a clear-text ASCII file. I'd like to hear comments from readers on this subject, especially those who have been victims or who can provide actual samples of these Trojan horse programs.

## **EXEC Calls and FORTRAN**

Robert Sypek of a firm called Argis in Hudson, Massachusetts, writes: "While attempting to write a program that would execute a user or system task from a user program, I ran into many of the problems described in past issues about using the DOS 2.x EXEC function call. My problems were compounded by the fact that Microsoft's languages (FORTRAN, Pascal, C) set up their own data and stack segments, and the code segment is that of the user program or subroutine, not the segment of the root program. This means that the program segment prefix (PSP) is inaccessible from any routine called by a user program. All is not lost, however, as I believe I have found a method of retrieving the necessary information to allow an EXEC function to be called.

"The method depends on two

pieces of information that may be gleaned from the Microsoft user manuals: the compiler defines a null segment called *HIMEM* that is placed at the end of all other segments when the user program is loaded, and the compiler defines a symbol called *MAIN* that is located directly after the PSP. Both the segment and the symbol have the *PUBLIC* attribute, making them accessible to the user.

"The segment of the PSP may be found by subtracting 10H from the segment of MAIN, yielding the value of the ES register needed for the call to EXEC. The value of HIMEM can then be used to calculate the size of the user program and the segment of the next available paragraph of memory."

Mr. Sypek enclosed an assemblylanguage subroutine that we are printing as Listing One (page 116). The routine is invoked in the form:

CALL FREMEM(MEMPTR,MEMAVL, MEMALL,ERR)

where *MEMPTR* is the segment value of the next available paragraph, *ME-MAVL* is the number of paragraphs available, and *MEMALL* is a user-defined variable specifying the amount of memory above the program to leave allocated to the user. (A value of —1 leaves all memory allocated to the program.) An error code of 0 indicates that the function was successful.

### Lightweight Reading

Those who think the mainframe mentality is gone forever should read Martin Healy's article "Toward a Viable OS for the PC" (Datamation, September 1985). At first I took this article for a practical joke because it contains so many distortions of the history and current state of the art in microcomputers, sideways slams at a variety of targets, and outrageous gobbledygook (example: "PC Network could be the answer, but it won't share data, only the file").

The author asserts that "there are in fact two leading real operating systems for micros, Unix and Concurrent DOS from Digital Research." He goes on to say, "The new Version 4.1 [of Concurrent DOS] is the idealized multitasking PC DOS, which due to its maturity should eliminate any Microsoft version (PC DOS 4.0?) thereof. That leaves Microsoft to concentrate on its Unix-like system, Xenix." I am sure the folks at Microsoft will be very relieved to learn from Mr. Healy that they no longer have to waste all that effort on maintaining MS DOS, can take MS DOS 4.0 out of testing and toss it in the trash can, and turn their attention to other matters.

### **Square** Roots

Michael Barr of the Department of Mathematics at McGill University sent us his 8086 assembly-language subroutine for square roots (Listing Two, page 116). He writes: "This routine gives the correct floored square root for any 32-bit number (considered as unsigned). It is also faster than the bit-at-a-time algorithms you have put into DDJ.

"Apropos that last statement, there seems to be a discussion between people who believe that Newton's method is always the best way to do a square root and others who believe equally fervently that the bit-at-atime method is always faster. Common sense would dictate that they are both wrong. I strongly suspect that Newton's method is faster if and only if you have an on-chip (or coprocessor) division of the relevant size. In particular, to do a 32-bit square root, you need a 32-bit by 16bit division. This much I have tested; Newton's method is just about twice as fast (about 330 msec, compared to about 650 msec for the bit-at-a-time). What I haven't tested (I can't face the thought of programming them) are 64-bit square roots. But there is every reason to believe that it will be faster to do the bit-at-a-time square root than to code division and then use that for Newton's method."

# **Big DOS**

The new wave of PCs, based on the 80286 microprocessor, are still limited to 1 megabyte of direct memory addressing because they run MS DOS or its clones in 8086 emulation mode

(called Real Mode by Intel). The ability of the 80286 to address 16 megabytes of RAM in "Protected Virtual Memory Mode" is currently supported only by the 80286 Xenix and iRMX-286 operating systems, neither of which is likely to become very popular due to their incompatibility with WordStar, dBASE III, and Lotus 1-2-3. The average user's exploitation of the full capabilities of the PC/AT and other such machines awaits the arrival of an operating system that runs in Protected Mode and can execute the popular

MS DOS applications unaltered.

Such an operating system has been the subject of much industry rumor and speculation. Digital Research has been talking about "Concurrent DOS-286" for some time now, even announcing the operating system's "immediate availability" at a press conference in New York about six months ago. (See "Concurrent DOS-286: Available Now," Intel Speak Softly Quarterly, vol. 2, no. 2, Second Quarter 1985.) But lately DRI has been backpedaling a bit and now admits

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(Continued from page 119)

that 80286 Concurrent DOS will probably never exist in the form originally advertised. Of course, the company is blaming its problems on "defects" in the 80286 design. (When in doubt, fall back on the classic programmer's defense: "It must be a hardware problem.")

Microsoft Corp. is also known to be working on a Protected Mode, upward-compatible version of PC DOS (already dubbed Big DOS by the trade rags and assumed to be Version 5.0). But Microsoft has been very close-mouthed about these efforts, presumably wanting to make sure it can pull it off before committing itself to such a product in public. Too bad DRI wasn't that careful!

Ross Nelson, who previously worked on the 80286 team at Intel, took the time to write to us with some musings on the future of DOS and Protected Mode. He says, "With regard to the 286 in the marketplace, it seems to me that unless a comparatively low-cost system (PC/II?) is introduced, there won't be a lot of work done that will take advantage of Protected Mode (PM). The installed base of AT users vs. the number of PC-compatible users will make it economically unfeasible. Once people do start working with PM, they will encounter some interesting problems. The great virtue of PM is that no task in the system should be able to corrupt another task (assuming the operating system is stable). As a software engineer, I applaud this philosophy, and I believe that this use of the 286 should be encouraged. Realistically, however, it is clear that there will be a transition period in which PM will only be used to gain access to the larger memory space.

Nelson wrote: "As far as I can tell, there are only two ways of switching back to Real Mode when you are in Protected Mode, and only one of them is feasible with the standard 80286 part. This is essentially the method IBM has chosen [in the VDISK driver supplied with PC DOS 3.x...RD], which is to place enough state information to restart your process in a 'safe' location and RESET the processor. The other method requires a special 'bond-out' part (which Intel uses in

its I2-ICE). By activating a special pin on the bond-out chip, you can issue special instructions that dump and restore the internal state of the machine, including the Machine Status Word. Systems built with the bond-out chip could easily be toggled between Real Mode and Protected Mode.

"Whether or not a DOS 5.0 or Big DOS can be successful in emulating the current PC system architecture on a 286 will depend on how freely the implementors translate the word compatible. I do not believe that 100

percent compatibility can be accomplished without unreasonable overhead. I suspect that even partial compatibility will have large memory requirements. It would not surprise me to see a 512K operating system with an additional 32K required on a per-task basis. Here is how I believe some of the problems that come up can be handled:

"Video—This has quite a few simple solutions: (a) trap each initial write to the display segment, and replace the offending segment register with

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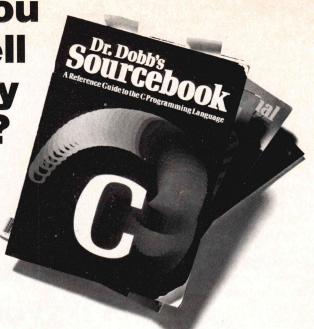
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an OS-created descriptor; (b) use the TopView solution, that is, require applications to issue a software interrupt to get the address of their own 'local' display buffer; (c) trap every display write and deal with it on a case-by-case basis. Solution (c) is the most compatible but the poorest performer; (a) is almost as compatible and would substantially increase performance. Solution (b) is the best and would not cause a great deal of incompatibility, especially with programs that use installable device drivers.

"Communications—Here we have no simple solutions. I would expect that all but the most primitive communications applications would have to be rewritten. Anything more complicated than Int 14H calls should be declared incompatible and rewritten to fit the new OS.

"EXE files—Here, the OS must limit the 'free-memory' size to 64K of code and 64K of data. Programs that need more memory should be forced to [dynamically] allocate it. Because of the large number of programs that indiscriminately load segment registers, however, an OS might want to trap segment load faults and attempt to map them into the 8086-style physical addressing scheme.

"COM files—These programs are the most likely to be incompatible, because they often load and stay resident and have only one segment (CS). Because executable segments are never writable, an OS would incur severe performance penalties trying to simulate the standard DOS mode of operation here. Some heuristics would work for some programs, such as loading DS, SS, and ES with a writable alias of the CS descriptor, but any program that did segment register arithmetic would then yield incorrect results.

"Obviously, there are hundreds of similar problems.... The question is, what solutions will the marketplace accept? It seems that a Protected Mode 286 operating system will have to contain a large measure of DOS/IBM PC compatibility to prevent the 'software gap' problem faced by the Macintosh, Amiga, etc. Unfortunately, this means a lower-performance,

aesthetically unpleasing solution. I would welcome a radically different, optimized system, but only IBM has the clout to pull it off, and there would still be two important factors: IBM would have to want to do it, and it would have to do it right."

### Nothing but the Best for My Little Girl

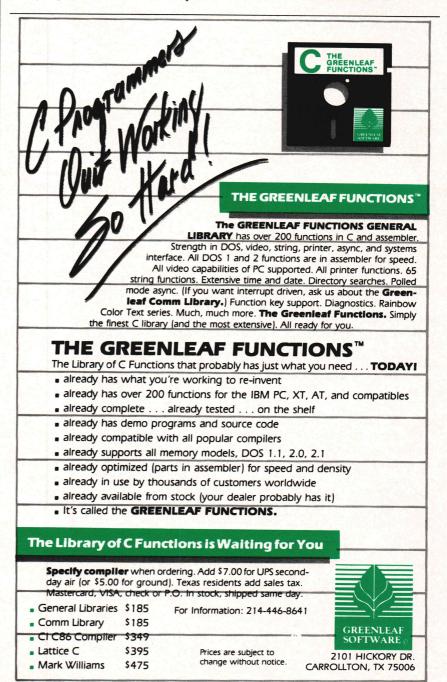
Richard Gaulden, vice president of sales and marketing for EDCOM, Inc., wrote me a letter describing his company as a "national supplier of enrichment materials to school and library systems." His sales pitch wound up with the statements: "ED-COM can free your distribution channels by marketing software that becomes outdated by new releases. The educational market does not have the need for constant upgrades, especially when new versions are more costly." Kind of gives you a warm, secure feeling about your kids' education in computer literacy, doesn't it?

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# (Listing begins on page 116)

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# PROFESSIONAL PROGRAMMER

# Kitchen-Table Entrepreneur

The purpose of this department is to provide information of use to professional programmers. This first installment addresses itself to the independent software developer, romantically and perhaps not altogether inaccurately envisioned as working on a VAX by day and by night designing the compiler of tomorrow on an XT on the kitchen table. We point here to a number of books that might be of use to the kitchen-table entrepreneur.

# Copyright

Read books on copyright with some caution. Both copyright law itself and the interpretations of it in court have changed in their application to software publishing. Don't assume that you need not copyright your software if you favor a Freeware/Shareware approach to distribution.

Remer, Daniel. Legal Care for Your Software. Berkeley, Calif.: Nolo Press, 1984.

Read this for discussion of all the legal issues: workfor-hire agreements, nondisclosure agreements, contracts, license agreements. The author covers the differences among and nuances of copyright, trade-secret protection, patent, and trademark registration. The book explains your legal liabilities and how international copyright works. It has a number of sample forms, such as a beta test site agreement form.

Salone, M.J. How to Copyright Software. Berkeley, Calif.: Nolo Press, 1984.

Read it for deeper discussion of copyright issues. What happens when the first 500 disks go out with no copyright notice on the disk or in the code? What can you do if your copyright is infringed? Hundreds of examples are provided.

### Documentation

Read something on documentation and realize how important it is. A good manual is, after all, the best copy protection available. Many programmers cut their own throats by releasing good software with illiterate, poorly designed documentation. Besides creating an aura of amateurishness for your product, such a manual will hide the capabilities of your product from usersand however experienced a programmer you may be, the odds are you're an amateur in producing documentation.

Houghton-Alico, Doann. Creating Computer Software User Guides. New York: McGraw-Hill, 1985.

Read it to jog your memory about some of the techniques available for communicating about your software—maybe what your compiler documentation needs is a good pie chart. Maybe not. Use the book to get an idea of what professional documentation writers have to do, though their tasks are somewhat different from yours, of course.

Stephan, Peter M. Writing User-usable Manuals. Salt Lake City: Wredco Press, 1984.

Read it as an example of decent low-cost documentation. The author has won awards for his documentation, and although the book may not tell you anything you don't already know, it shows that a manual need not be typeset and perfect bound.

Strunk, William, and White, E.B., *The Elements of Style*. New York: Macmillan, 1972.

Read it. Whether you write the documentation or hire writers, your words are likely to see print somewhere. There is no more concise guide to keeping your foot out of your semiliterate mouth than this book.

### Markets

Read these if you're really an independent software developer and don't want to distribute your product(s) yourself. These books mainly list software distributors.

Amato, Francis. *Guide to Computer Magazines*. Dallas: Steve Davis, 1985.

Read it for a quick idea of the editorial focus, circulation, and audience of selected computer magazines. This can be useful for promotional and advertising purposes, and some of the magazines are software markets in themselves, publishing programs in their pages and/or on disks.

Hoffman, Roger. The Complete Software Marketplace. New York: Warner Books. 1984.

Read it for many reasons. One reason is to get an overview of the things you need to know in running a small software business. The book has lists of distributors, including mail-order companies, electronic distributors, and magazines. There are also lists of disk duplication services, PR firms, market research firms, trade associations and shows, seminars, and conferences. There are case studies and a useful section on freebies for the independent author.

McGehee, Brad M., ed. 1986 Programmer's Market. Cincinnati: Writer's Digest Books, 1985.

Read it for publishers who use "free-lance" material—even the terminology reflects this book's heritage. It was put together by people who view the independent software developer as another kind of author. If that description fits you, this book may also.

### Software Design

Read this to remind yourself of the basic principles and also to learn a vocabulary for communicating the principles to others.

General Electric. Software Engineering Handbook. New York: McGraw-Hill, 1986.

Read it for one approach to the management of large software projects when your design team actually becomes a team. Also read its bibliography for sources on design methods: Nassi/Sneiderman, Warnier, Yourdon, and Jackson methodologies; data-flow driven design.

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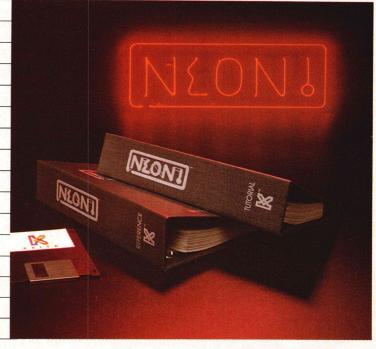
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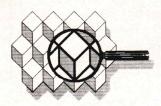
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# OF INTEREST



Peak Electronics has announced a 68K coprocessor board that, it claims, you can plug into any (IEEE-696-1983) S-100 system running CP/M-2.2 and have CP/M-68K running in minutes with no hardware or software modifications. The 68K8-CP is a 10-MHz processor card with an MC68008 (8-bit version of the 68000), up to 512K RAM, up to 128K EPROM, two 38.4 kilobaud serial ports, and an 8-bit parallel printer port. The 68K8-CP doesn't replace your current CPU card but runs in conjunction with it, so you can jump back and forth between operating systems. Peak says that the board will support Concurrent DOS 68K when it's released.

Speaking of Concurrent DOS 68K, maybe DRI is finally going to exploit the head start it had in 68K development software: The company was expected to introduce its 68K developer's kit at Comdex in November. With Concurrent, DRI has two operating systems for the 68000; three if you count the GEM windowing environment. Concurrent DOS 68K has a CP/M front end that will run most CP/ M-68K programs, and DRI is promising GEM support for subsequent versions of Concurrent DOS 68K.

Speaking of GEM, last month we discussed it and other windowing environments and, because of space limitations, mentioned only some of the most visible. We don't feel right about giving all that attention to the big boys and never mentioning two windowing products we have had good experience with: Desqview and dWindow. It was, of course, not only Digital Research, IBM, and Microsoft that went windowing. Quarterdeck's Desgview is an impressive competitor to TopView, which bundled with AST's Rampage expanded memory board, lets users run up to nine programs concurrently in memory. Also, the success of the concept of windows is exemplified by its employment in a nonoperating system application in the product dWindow. For vears, Ashton-Tate's dBASE II set the standard for austerity in user interfaces: You can't get much simpler than a single-period prompt. Liberty Bell Publishing's dWindow does a dazzling cathedral window treatment on dBase II (and now dBase III) that makes it look like an entirely different product.

Speaking of entirely different products, Answer Software has announced an 80286 emulator for developers of 80286-based software. The ICD286 consists of a card for the host system (which must be in the IBM PC/XT/AT family or a compatible), a Buffer Box that plugs into the 80286 slot in the target system, and a symbolic debugger. It allows uploading and downloading of code and data, hardware and software breakpoints, singlestepping, and full-speed emulation up to 10-MHz clock rates.

And speaking of 80286 development, American

ADO has introduced an 80286 board for Multibus systems. The SOL C286-01 (no relation to the Processor Technology Sol computer of story and song) is being manufactured in 6-, 8-, and 10-MHz versions and has up to 512K RAM, a Centronics printer interface, and two 8- or 16-bit SBX bus I/O connectors. Then there's the ET-286Plus, a 10-MHz AT-compatible single-board computer that uses the new 1-Mb dynamic RAMs and allows 4 megabytes onboard. ATS International was expected to show it at Comdex.

As long as we're speaking of the 286 and operating systems, we should mention Locus Computing's Multisystem Merge. This product allows simultaneous, transparent execution of Unix and MS DOS on the same machine according to the company. You can set several Unix tasks to work in the background while you run a DOS application in the foreground. This is the system AT&T is using on its 6300+ Unix/DOS computer. Locus developed some of the technology in the system while writing PC-Interface, a product that links pos computers to a host Unix machine.

DRI, which we were speaking of a few paragraphs back, is of course not the only developer of operating system software for the 68K, as two recent announcements prove. U S Software has announced a real-time multitasking system for embedded applications using the 68K. It's ROMable, requires 3K of code space, and is called USX68K. Integrated Busi-

ness Computers has ported TheOS-16 (formerly Oasis) to its line of 68010 computers and was expected to be showing a beta version at Comdex.

Speaking of beta-test versions, Tall Tree Systems has begun shipping beta versions of its Jlaserprinter interface to software companies it deems closest to achieving compatible products. The interface is a PC/XT/AT card that works with the JRAM 2-megabyte memory board; it's designed to spend memory to buy print speed and typeface flexibility for laser printers. It transfers bitmapped images directly from RAM to the print mechanism and is supposed to provide unlimited type fonts with full graphics capabilities at 300 dots per inch in eight seconds.

Speaking of *mucho* megabytes, Reference Technology has announced a device that lets you attach up to eight of its optical disk drives, for more than 4 gigabytes of storage on a (sturdy, large) desktop. The device is PC/XT/AT or compatible compatible.

Speaking of product an-Speech nouncements, Technology is now selling the cross-assemblers it developed in the process of designing and manufacturing electronic devices to aid the blind (its real business). The MS DOS cross-assemblers for the 8048 and 6502 were written in C and support a subset of C preprocessor commands, macros, three object file formats, and features to support PROM programming. They sell for \$30 each or \$75 for both with source code and are distributed with no restriction on noncommercial copying.

Speaking of copying, we are watching with interest SoftKlone's fortunes with a product it cheerfully describes as a clone of Microstuf's Crosstalk XVI data-communications package. SoftKlone's Mirror was designed to the precise visual specs of Crosstalk, and SoftKlone presents itself as introducing a new idea-mirrorimage software at a lower price than the mirrored product and perhaps with added capabilities. Rather than competing by producing a better or cheaper product, the notion here is to produce the same product better or cheaper.

### Reference Map

Peak Electronics, P.O. Box 700112, San Jose, CA 95170; (408) 253-5108. Reader Service Number 16.

Digital Research, P.O. Box DRI, Monterey, CA 93942; (408) 649-3896. Reader Service Number 17.

Liberty Bell Publishing, 618 N.W. Glisan, Ste. 203, Portland, OR 97209; (503) 221-1806. Reader Service Number 18.

Quarterdeck Office Systems, 1918 Main St., Ste. 240, Santa Monica, CA 90405; (213) 392-9851. Reader Service Number 19.

Answer Software, 20863 Stevens Creek Blvd., Cupertino, CA 95014; (408) 253-7515. Reader Service Number 20.

American ADO, 1840 West 186th St., Ste. 200, Torrance, CA 90504; (213) 532-5010. Reader Service Number 21.

ATS International, 2105 Luna Rd., Ste. 330, Carrollton, TX 75006; (214) 247-5151. Reader Service Number 22.

Locus Computing Corp., 3330 Ocean Park Blvd., Santa Monica, CA 90405; (213) 452-2435. Reader Service Number 23.

U S Software, 5470 N.W. Innisbrook Pl., Portland, OR 97229; (503) 645-5043. Reader Service Number 24.

Integrated Business Computers, 21621 Nordhoff St., Chatsworth, CA 91311; (818) 882-9007. Reader Service Number 25.

Tall Tree Systems, 1120 San Antonio Rd., Palo Alto, CA 94303; (415) 964-1980. Reader Service Number 26. Reference Technology, 1832 North 55th St., Boulder, CO 80301; (303) 449-4157. Reader Service Number 27.

Speech Technology Inc., 16321 176th Ave. N.E., Woodinville, WA 98072; (206) 483-5150. Reader Service Number 28.

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Increase your productivity and avoid agonizing waits. Get instant feedback of your C programs for debugging and rapid prototyping. Then use your compiler for what it does best...compiling efficient code ...slowly.

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- Full K&R C (no compromises)
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- Convenient-- Compiling and running are only a key-stroke or two away. Errors direct you back to the editor with the cursor set to the trouble spot.
- Object Module Support

   Access functions
   and externals in object modules produced
   by your compiler. New: We are now
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   & Aztec C in addition to C.I. C86 & Lattice.
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Price of demo includes documentation & shipping within U.S. PA residents add 6% sales tax. Specify compiler.

 C-terp runs on the IBM PC (or compatible) under DOS 2.x with a suggested minimum of 256Kb of memory.
 It can use all the memory available.

# GIMPEL SOFTWARE

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Mouse	Yes	No	Yes	Yes
Screen Resolution (Non-Interlaced Mode) Color Monochrome	640 x 200 640 x 400	640 x 200 720 x 350**	None 512 x 342	640 x 200*** 640 x 200 ***
Color Output	Yes	Optional	None	Yes
Number of Colors	512	16	None	4096
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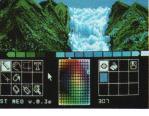
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